

Center for
Quantum Networks
NSF Engineering Research Center



Principles of Quantum Networks

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Funded by National Science Foundation Grant #1941583



Outline

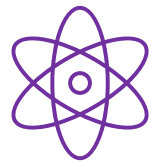
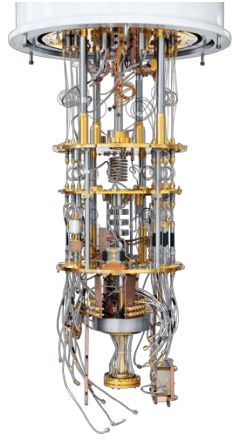


- Introduction
- Fundamentals of Quantum Communications
- Classical vs. Quantum Networks
- Scheduling in Quantum Repeater Chains
- Quantum Network Routing
- Connectionless Quantum Networks
- Quantum Network Management and Tomography
- Summary

Introduction

Purpose

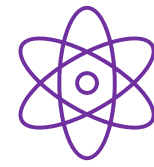
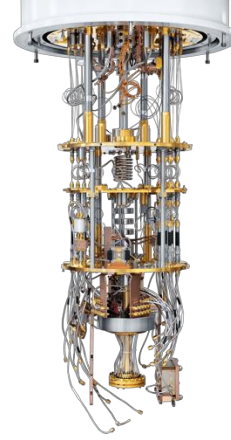
Alice



$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$$



Bob



$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$$

The Quantum Internet

Vision: Quantum network enabling full quantum connectivity between multiple user groups.



Secure Communications



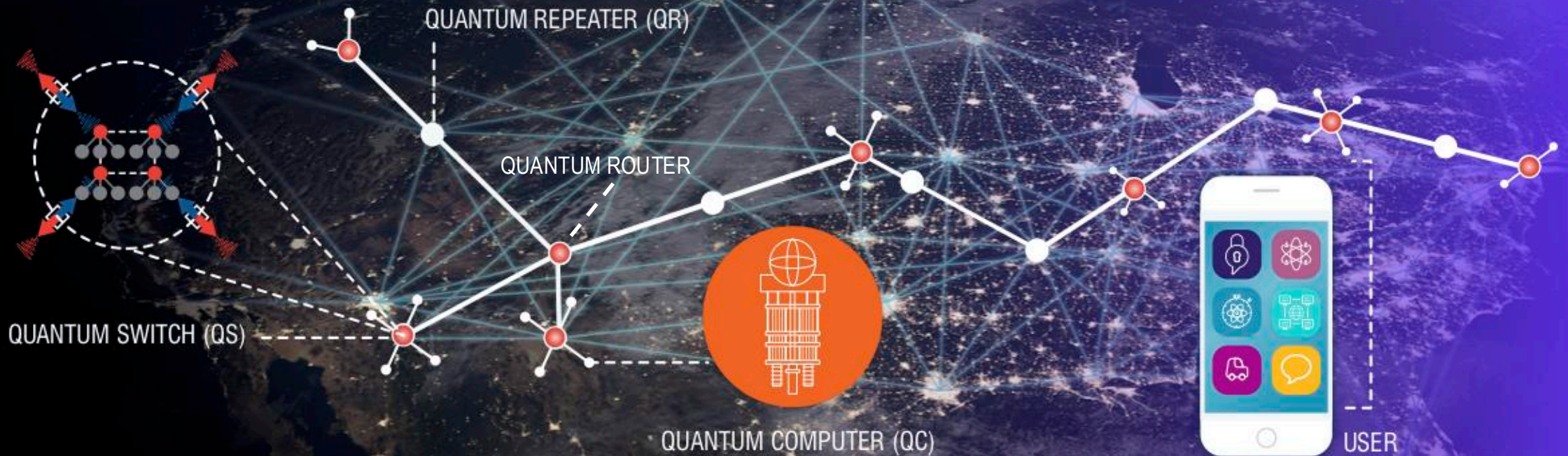
Quantum Multi-User Applications



Sensing, Timing, GPS

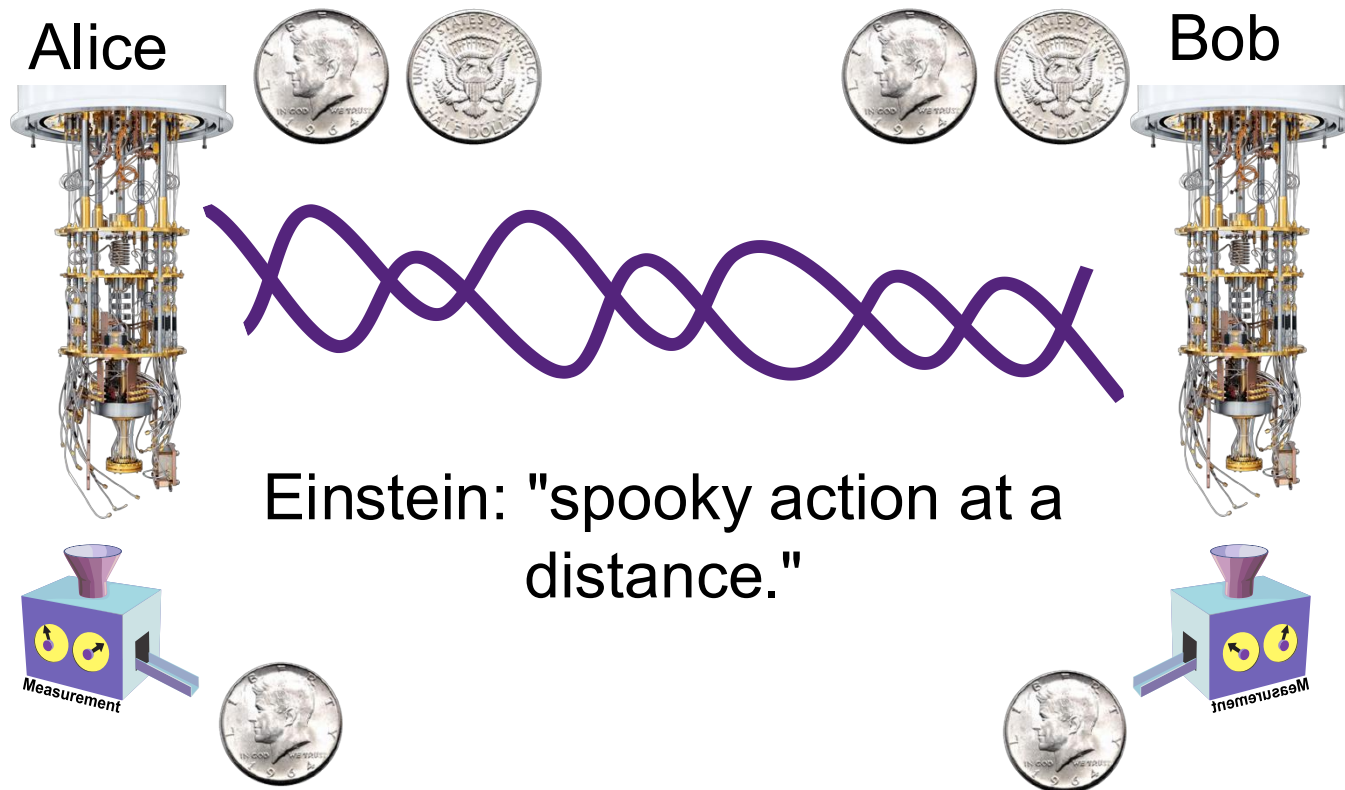


Networked Quantum Computing



Key ingredient

Quantum entanglement, *aka Bell state*, between pair of remote quantum processors



$$\text{Bell state: } \frac{|0_A 0_B\rangle + |1_A 1_B\rangle}{\sqrt{2}}$$

Nobel prize, Physics, 2022: A. Aspect, F. Clauser, A. Zeilinger



Why Quantum Internet?



Cryptography, security – quantum key distribution (QKD)

Distributed quantum computing – breaking web security, solving hard problems

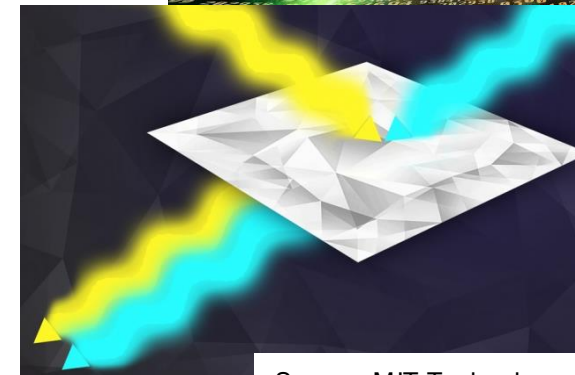
High resolution sensing – exploring the universe

Quantum Teleportation – transmission of quantum information

Source: Physics World



Source: IQOQI, H. Ritsch



Source: MIT Technology

Bell state

- Bell state

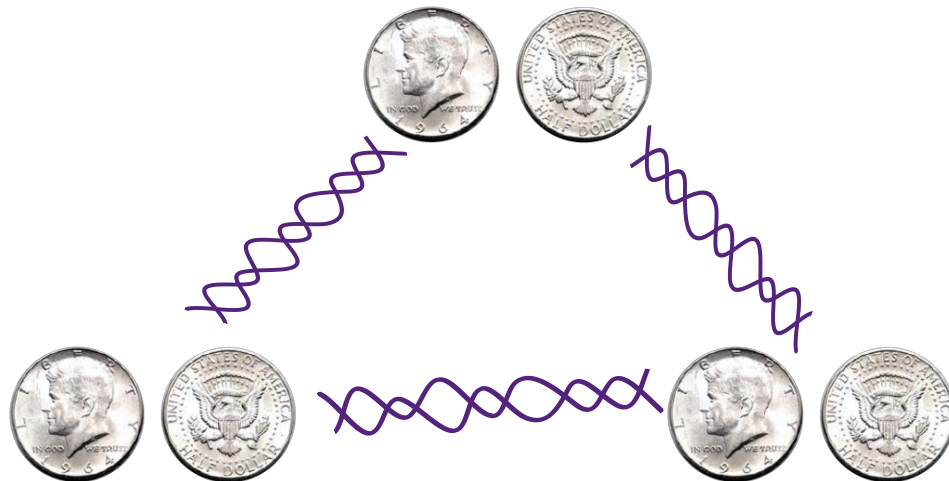
$$\frac{|0_A 0_B\rangle + |1_A 1_B\rangle}{\sqrt{2}}$$

- Measuring Alice's qubit yields 0,1
 - if 0, measuring Bob's qubit yields 0
 - if 1, measuring Bob's qubit yields 1
 - can generate shared randomness across distances
- Key ingredient of quantum teleportation, QKD, and many other applications



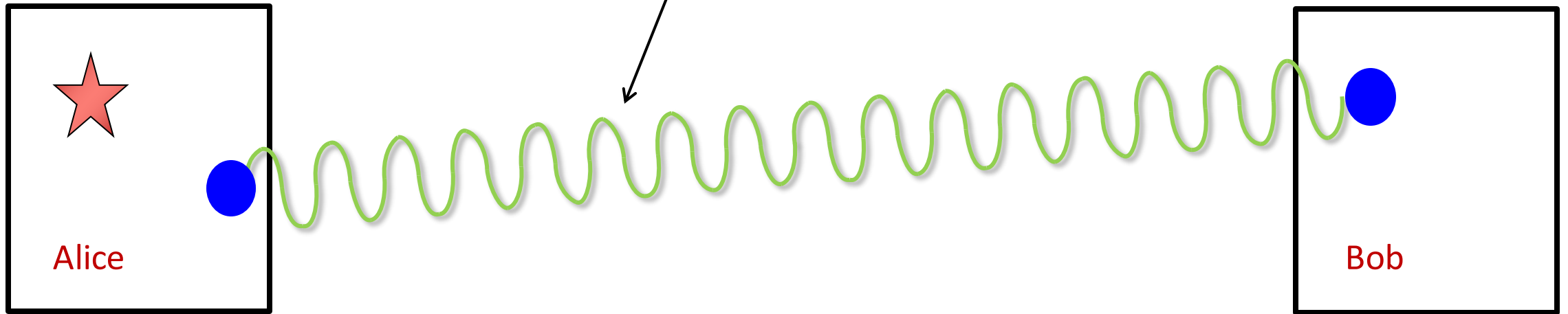
Greenberger–Horne–Zeilinger (GHZ) state

- n -partite GHZ state $|GHZ\rangle = \frac{|00\dots 0\rangle + |11\dots 1\rangle}{\sqrt{2}}$
- used in multiparty QKD, secret sharing, quantum sensing, ...

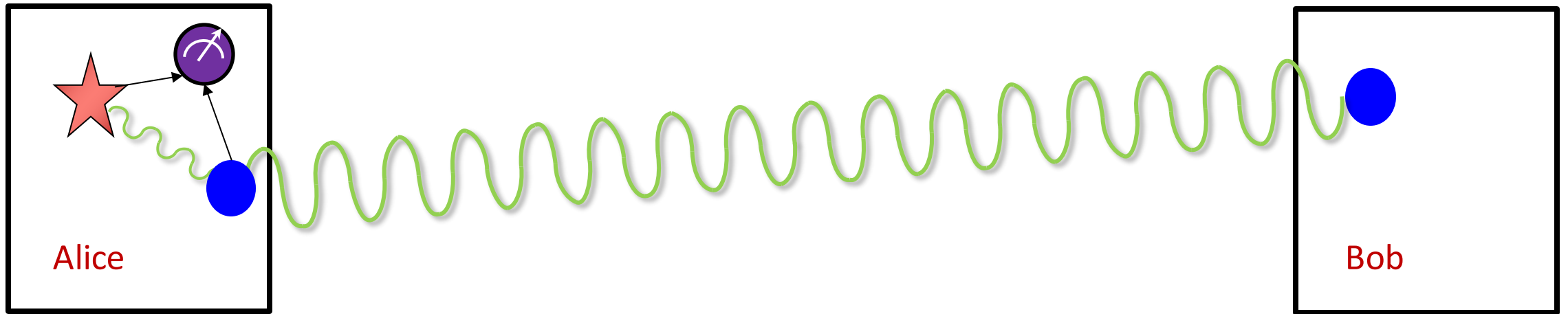


Quantum Teleportation

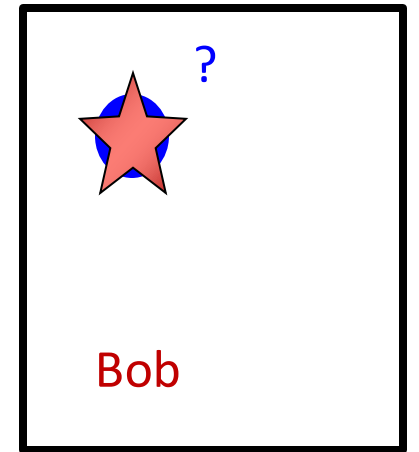
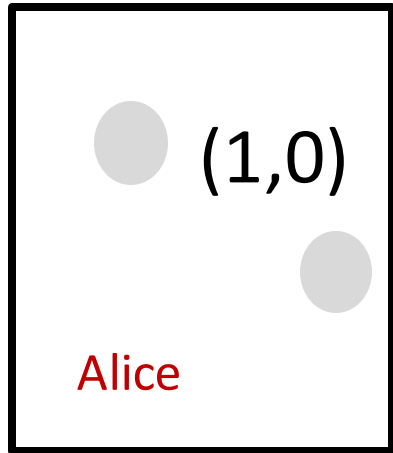
end-to-end entanglement $\frac{|0_A 0_B\rangle + |1_A 1_B\rangle}{\sqrt{2}}$



Teleportation

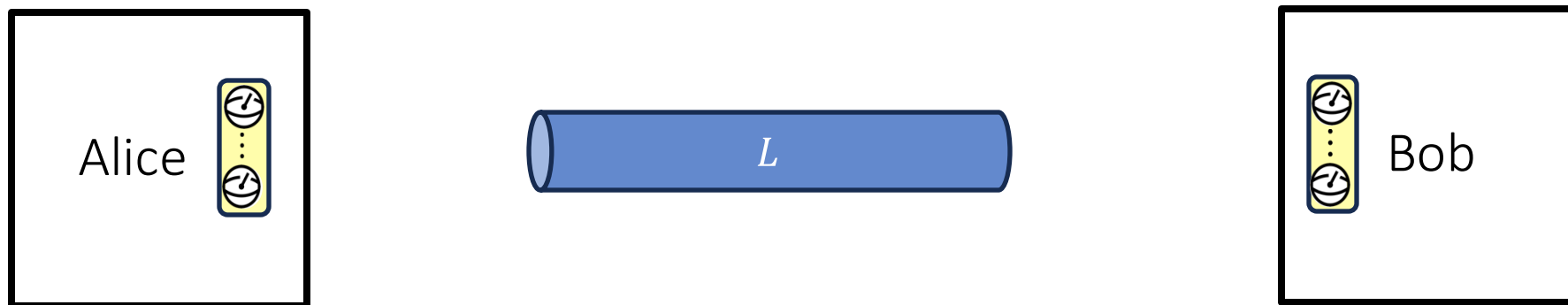


Teleportation



Fundamentals of Quantum Communications

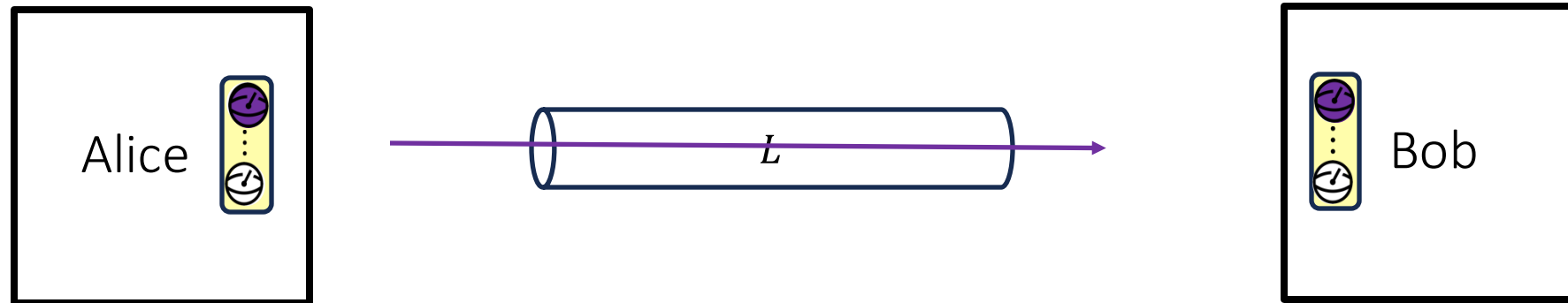
Quantum comms over link



- Nodes have qubits (memories)
- Photons are quantum information carriers
- Fiber/free space link connects nodes

Quantum comms over link

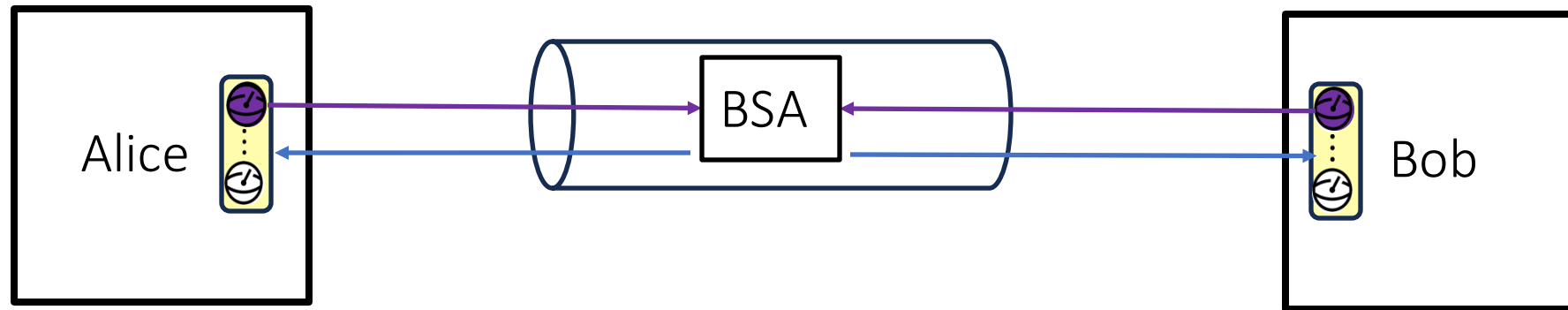
One-way communication



- Alice prepare qubit in state $|\psi\rangle$
- State transduced from memory to photon
- Photon sent to Bob through channel
- State $|\psi\rangle$ loaded in Bob's memory qubit

Quantum comms over link

Two-way communication

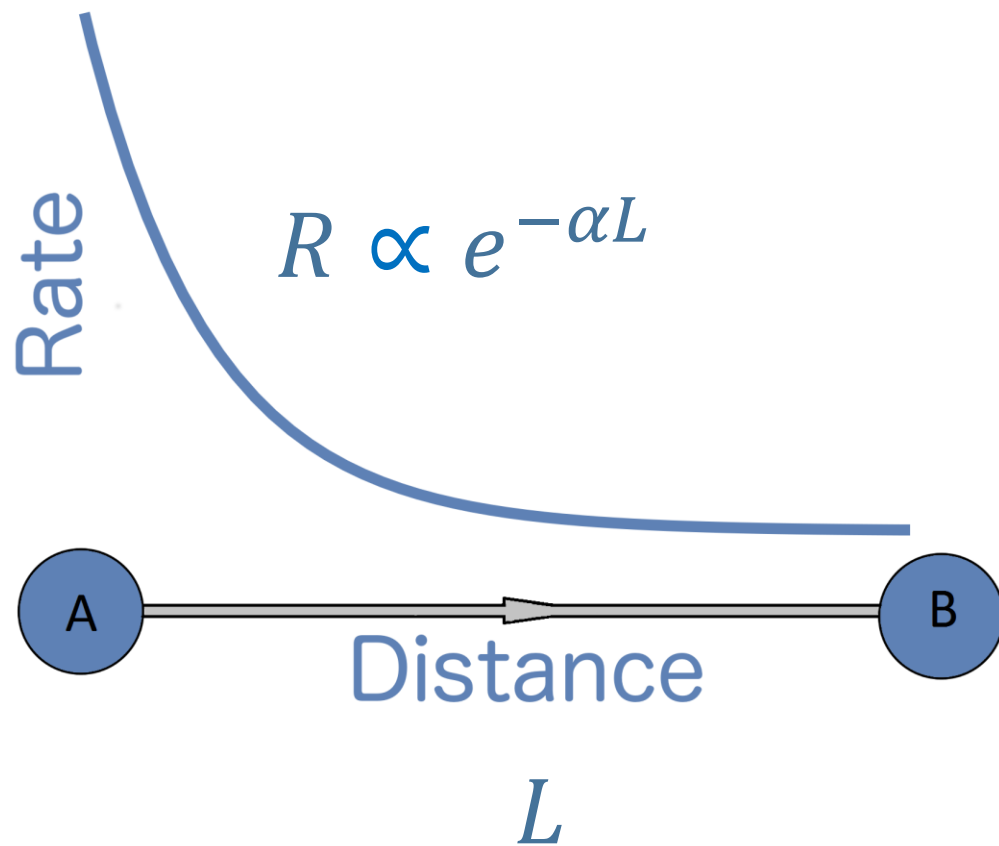


- Alice/Bob generate photon-memory Bell pair
- Photons traverse and meet at Bell state analyzer
- Measurement results sent back to Alice and Bob
- Bell state shared between Alice and Bob
- Additional operations herald state $\frac{|00\rangle + |11\rangle}{\sqrt{2}}$

Why are quantum communications so
hard?

Loss and noise!!!

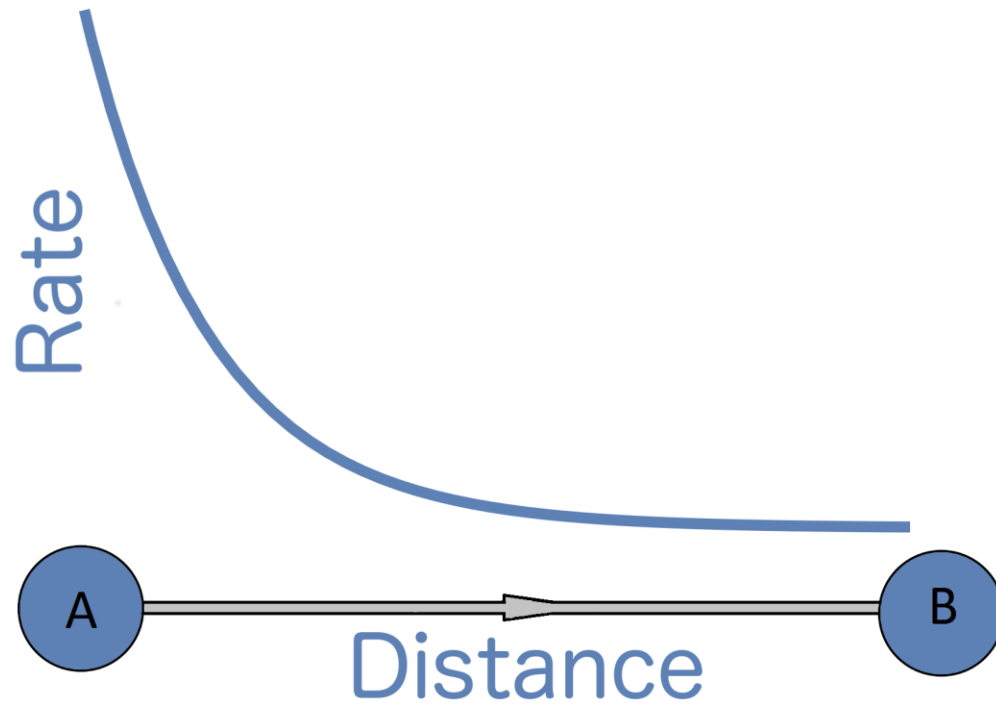
Why are quantum communications so hard?



Can we amplify signal?

Rate decays exponentially
with distance

Why is it so hard?



No cloning theorem!
Quantum signals
cannot be copied

Rate decays exponentially
with distance

Poll Question

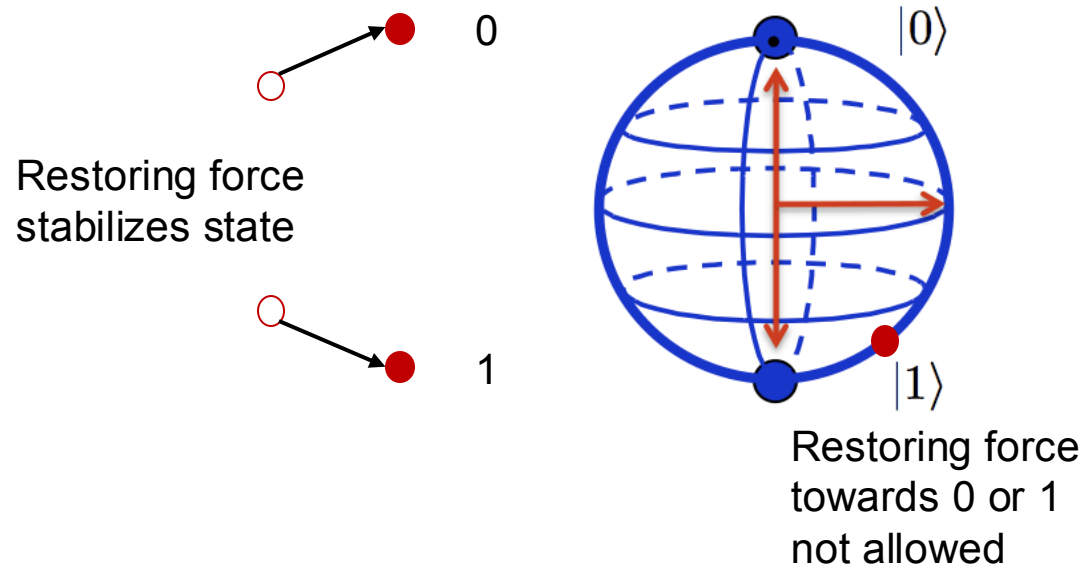


How does the generation rate of EPR pairs decay with distance in direct transmission through fiber?

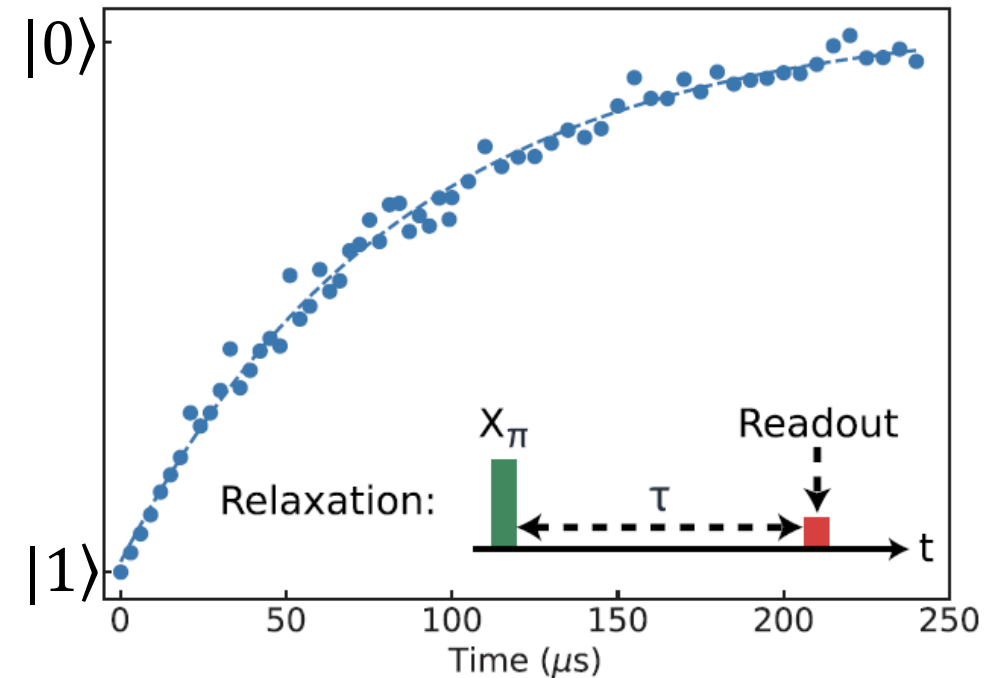
- A. Polynomial decrease
- B. Constant decrease
- C. Exponential decrease
- D. Depends on fiber technology used
- E. I don't know

Why are quantum communications so hard?

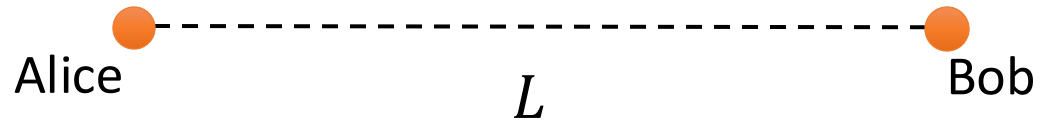
- Qubits not self protected against smallest perturbation



- Qubits have limited coherence times



Noise in quantum states



- Noise introduced thru comms
- States decohere, Bell states decohere twice!
- Noisy gates and memory operations

	One-way	Two-way
Expected	$ \psi\rangle$	$\frac{ 00\rangle + 11\rangle}{\sqrt{2}}$
Obtained	1-qubit state ρ	2-qubit state ρ

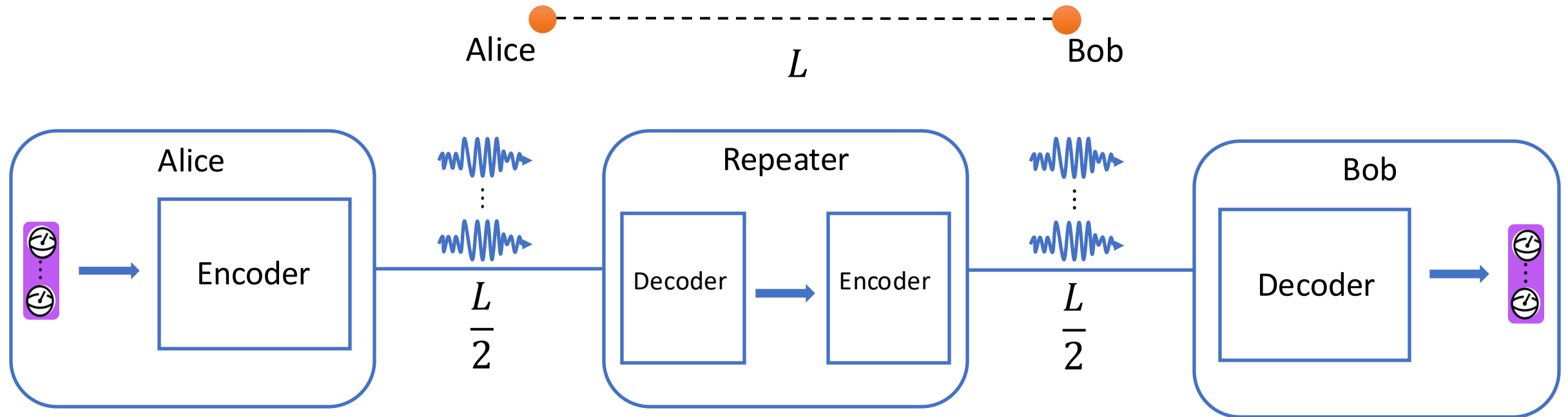
Fidelity: measure of closeness between two quantum states



Fidelity of direct transmission is between ρ and $|\psi\rangle$

Entanglement fidelity is fidelity between ρ and $\frac{|00\rangle + |11\rangle}{\sqrt{2}}$

One-way quantum repeaters



Protect information with Quantum Error Correction (QEC)

State recovered if sufficient photons survive

High-probability of delivering logical state, p depends on encoding

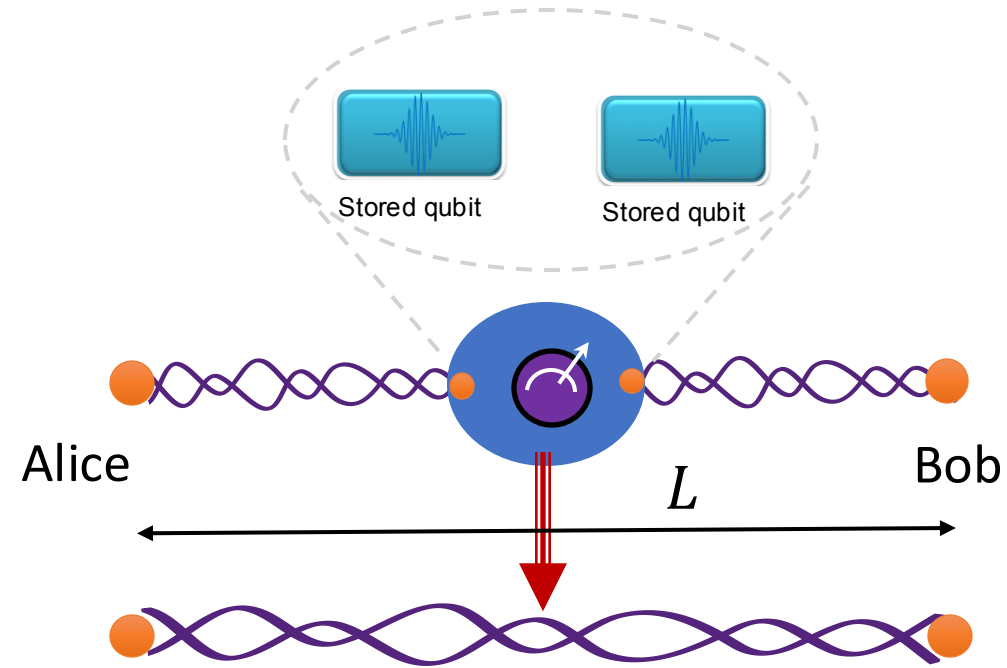
Repeater may not need quantum memories

Two-way quantum repeaters

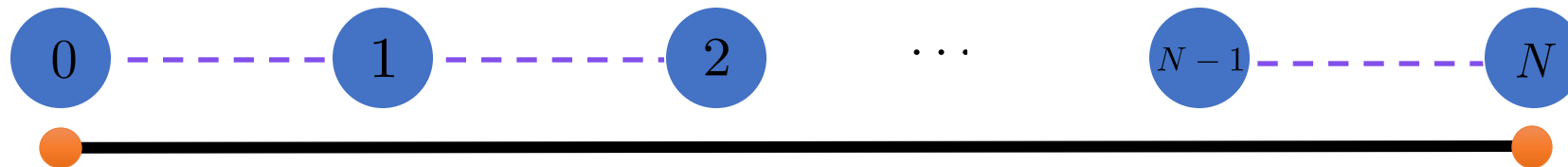
Quantum memories to store entanglement

Phase I: generate link level entanglement (Bell states)

Phase II: measurement propagates entanglements to ends



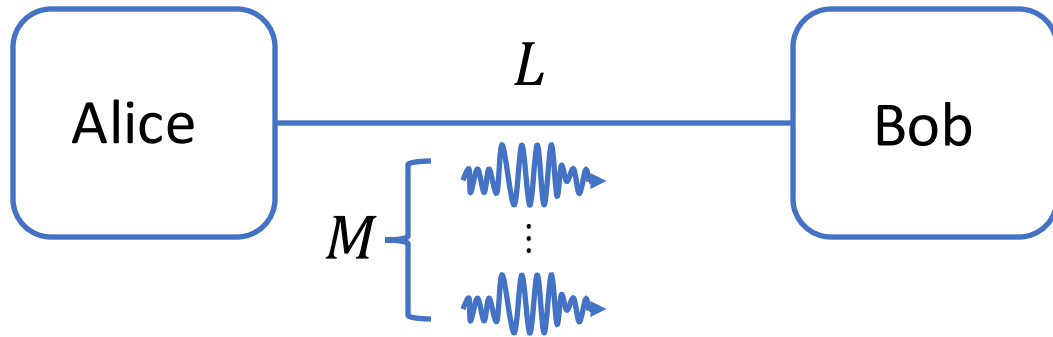
$$R = e^{-\alpha L/2}$$



$$R \propto e^{-\alpha L/N}$$

Multiplexing

What if we do parallel attempts?

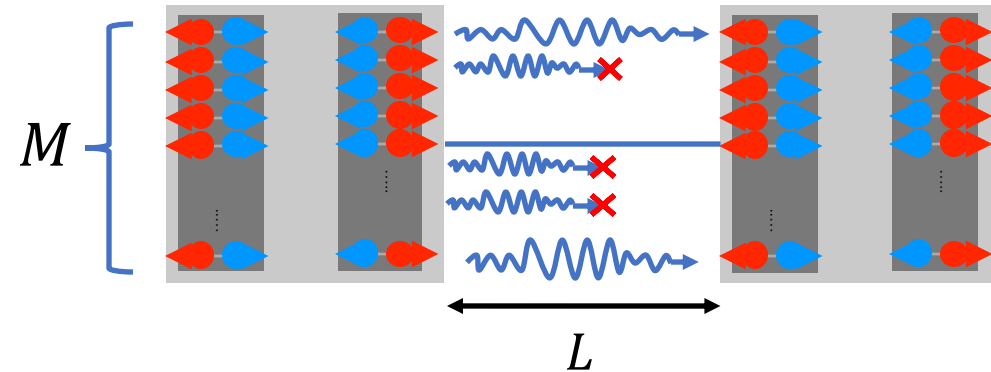


One-way repeaters utilize multiple photons!

QEC is like “quantum multiplexing!”

Use multiple entangled qubits (photons)!

In two-way repeaters



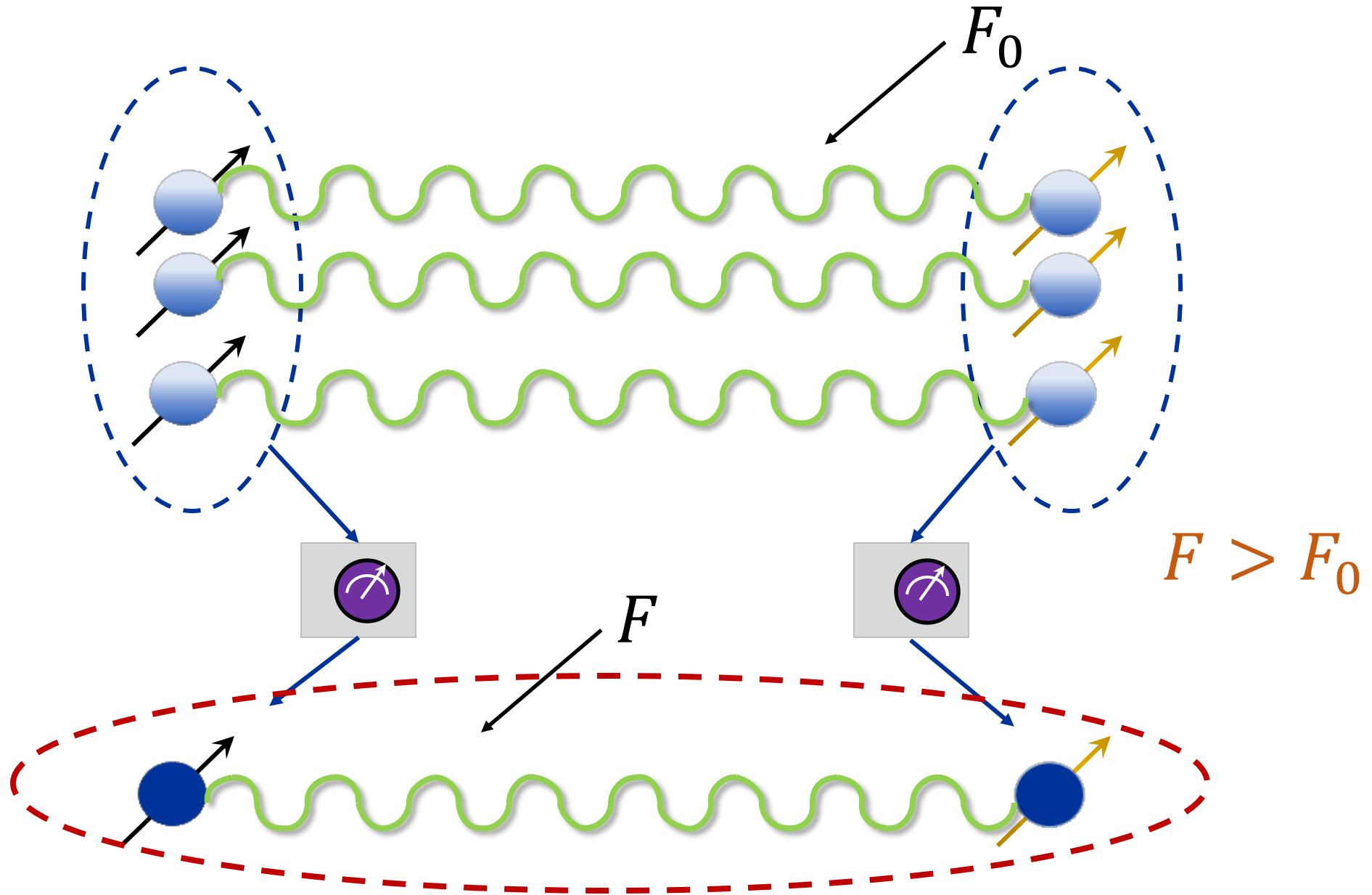
Probability that at least one survive?

$$\Pr[K = 1] = 1 - (1 - p)^M$$

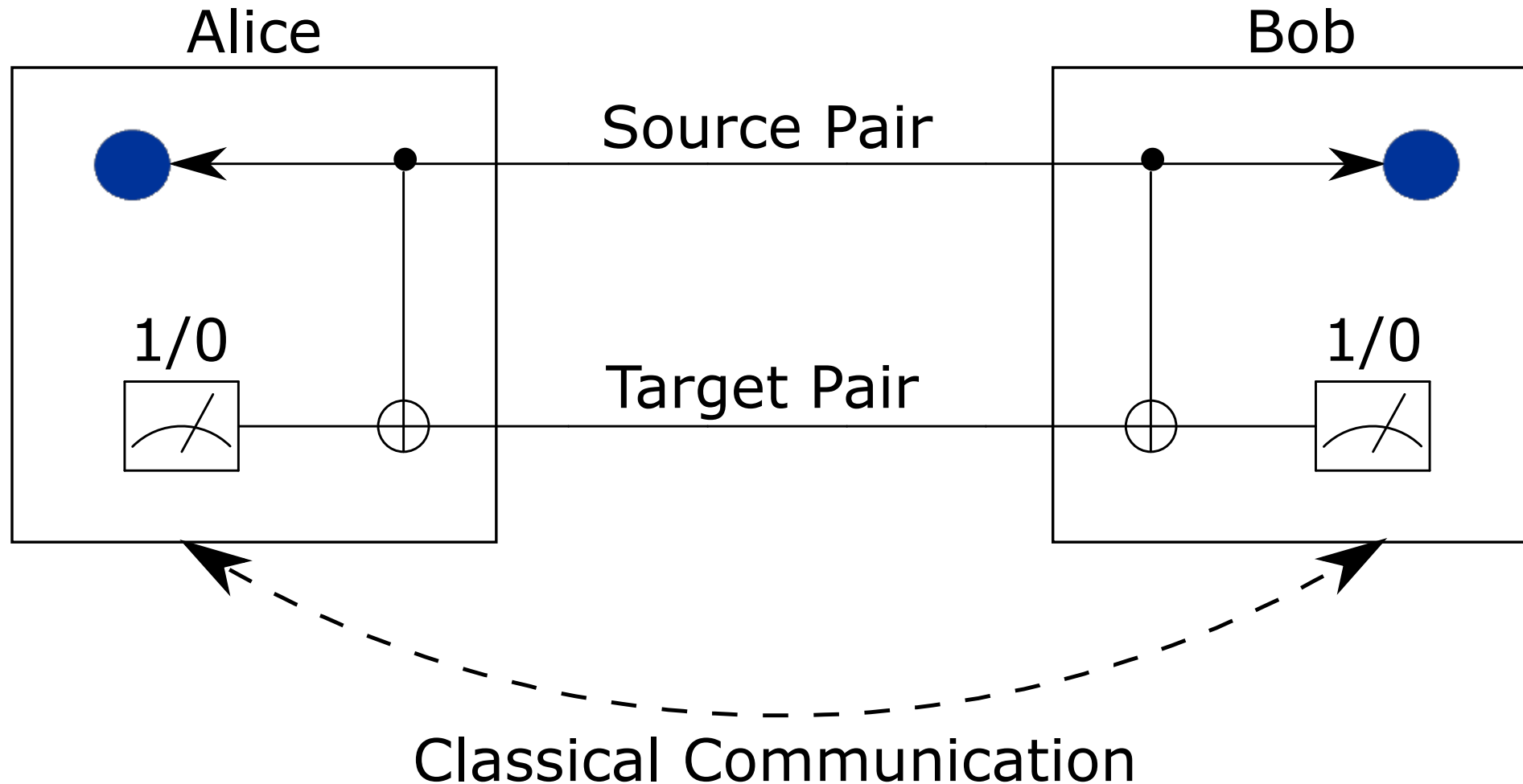
Probability that k survive?

$$\Pr[K = k] = \binom{M}{k} p^k (1 - p)^{M-k}$$

Entanglement distillation



Entanglement distillation



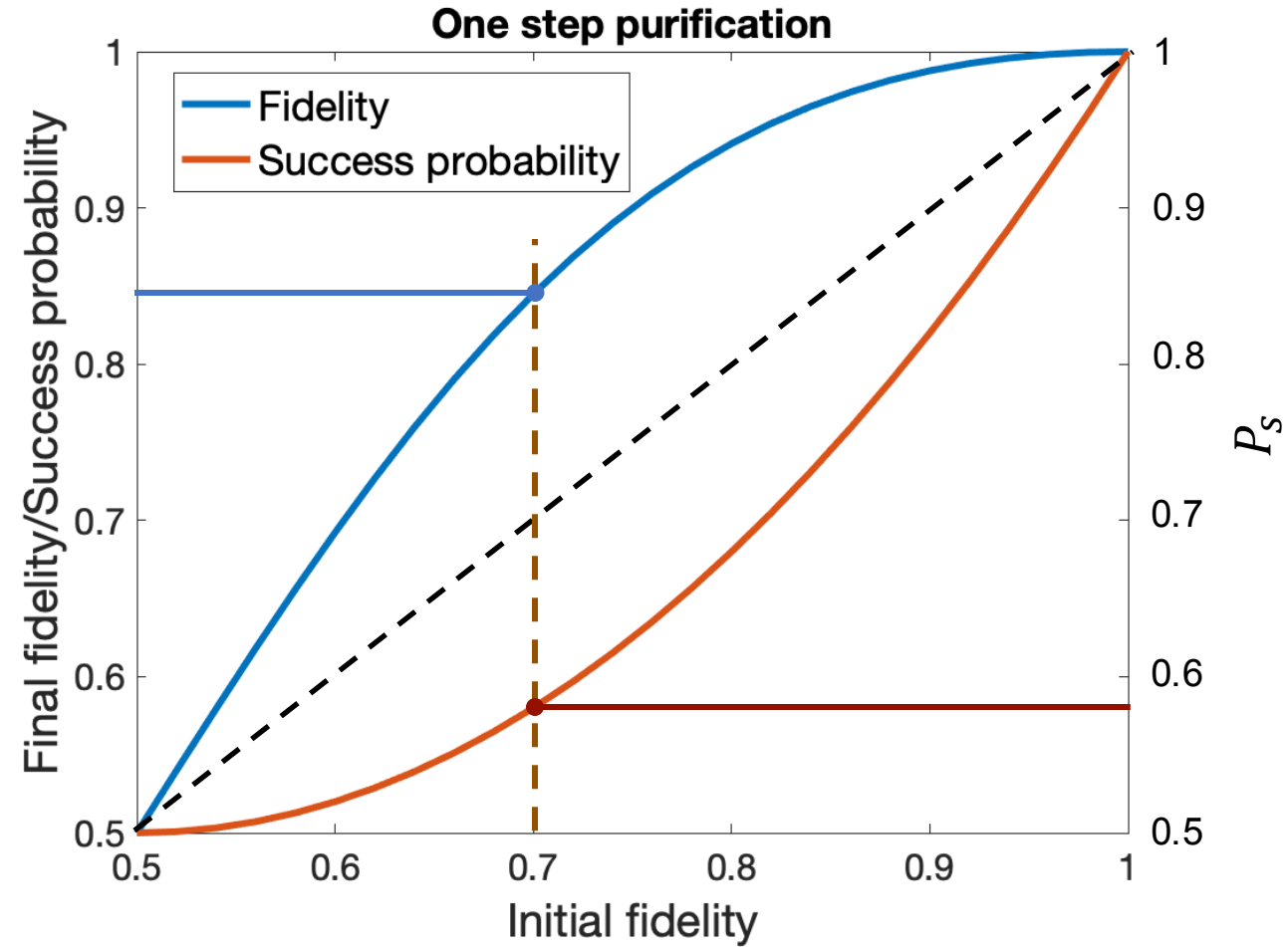
Probabilistically convert multiple noisy entangled pairs into single strongly entangled pair!

QoS metric

Fidelity: measure of closeness of entanglement to perfection

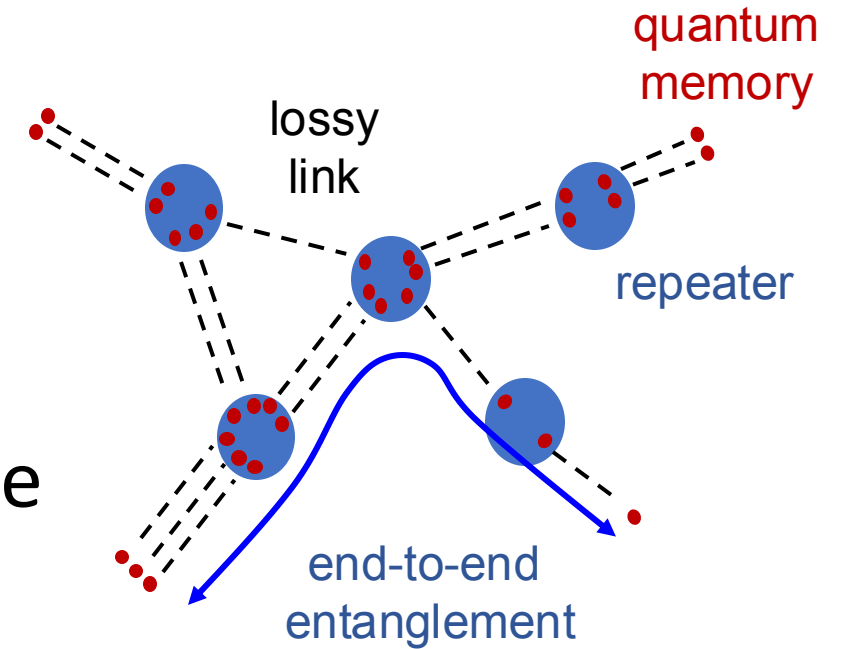


Distillation step succeeds with probability P_s



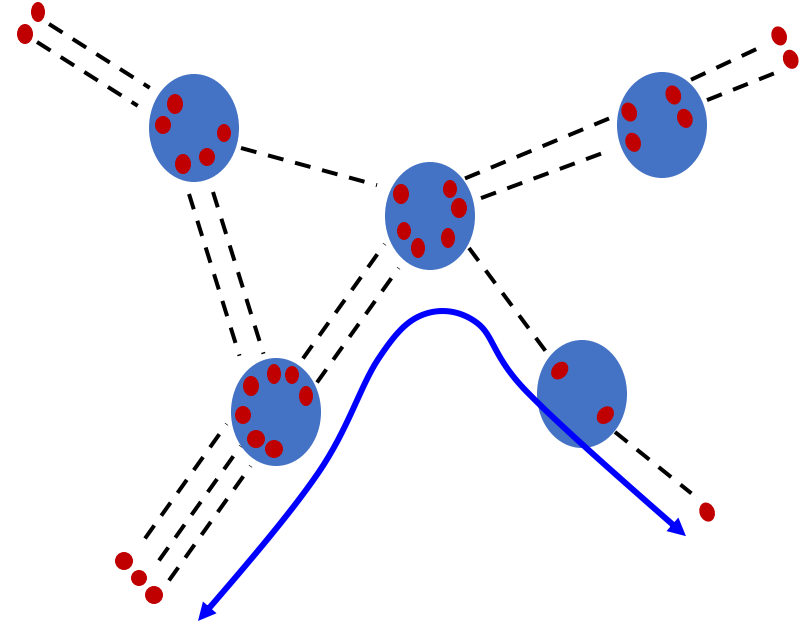
Quantum Networks

- Quantum switches with memories connected via lossy links
- Links generate entanglement / used to transmit quantum info. directly
- Switches concatenate (measure) to realize end-to-end entanglement between end nodes / decode-encode and forward quantum info



Quantum networking challenges

- Service to provide
 - entanglement distribution
 - direct quantum information transfer
- Noise!
- Who to serve
 - performance & resource allocation



- Network management
 - measurement & tomography
- Data, control plane design

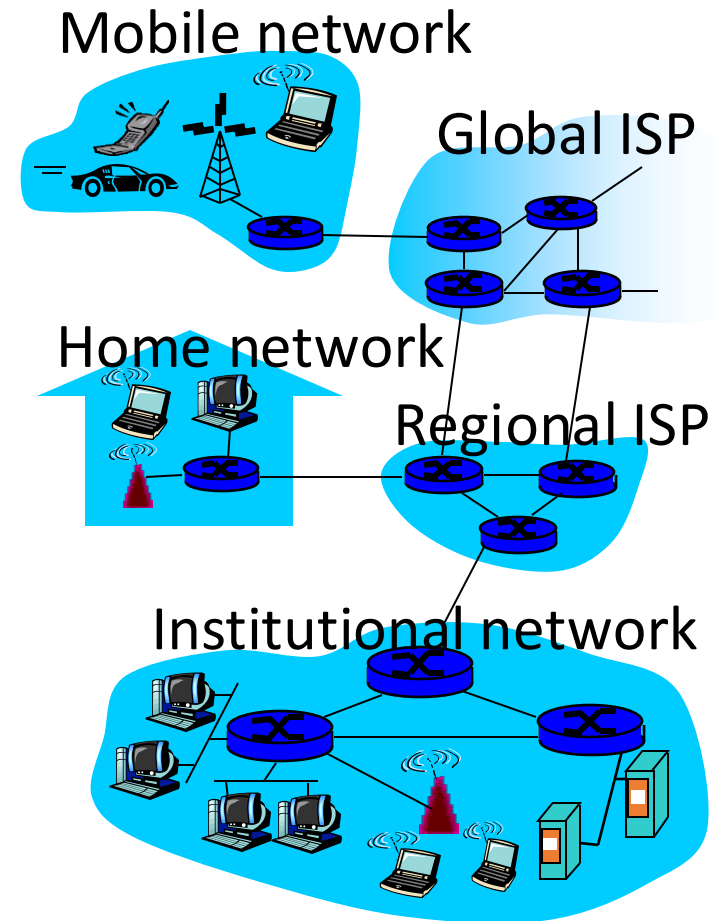
Classical vs. Quantum Networks

- Internet overview
- Network services, routing
- Switch/router design

What's the Internet: "nuts and bolts" view

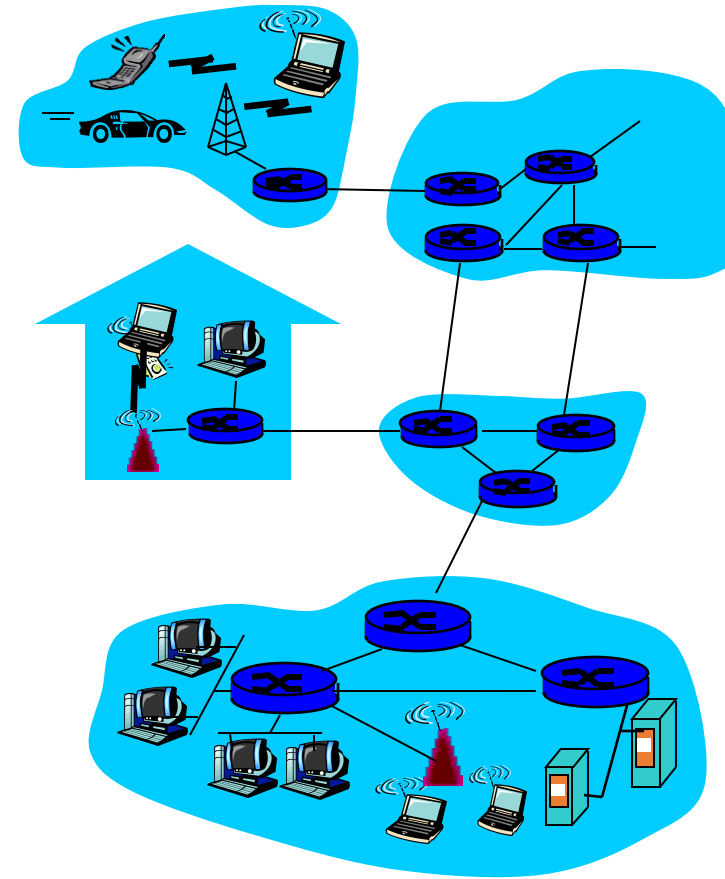


- **Internet: "network of networks"**
 - loosely hierarchical
 - public Internet versus private intranet
- **Protocols:** control sending, receiving of messages
 - e.g., TCP, IP, HTTP, RTMP, Ethernet, WiFi
- Internet standards
 - RFC: Request for comments
 - IETF: Internet Engineering Task Force
 - IRTF: Internet Research Task Force
 - QIRG: Quantum Internet Research Group



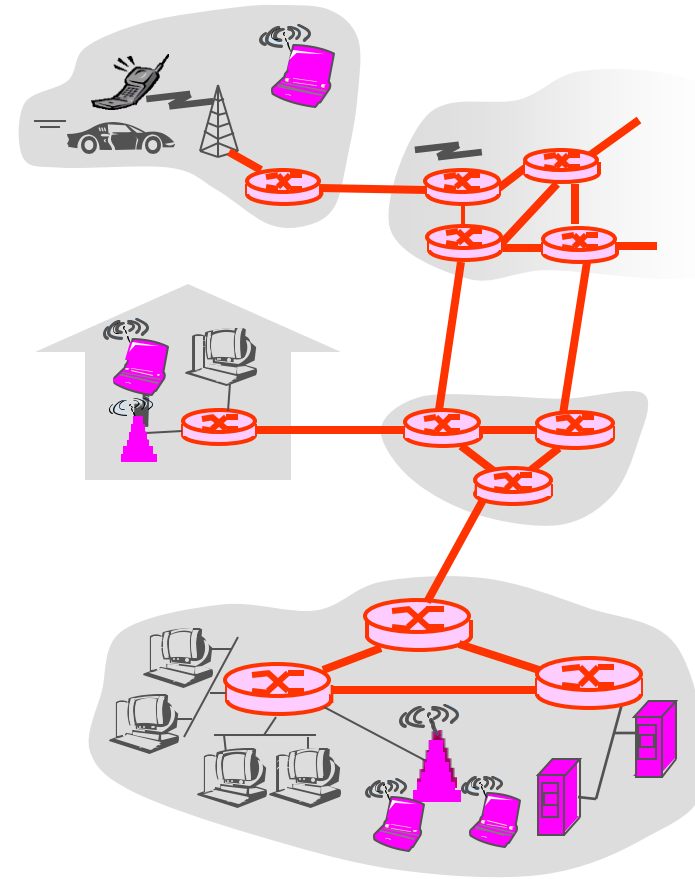
A closer look at network structure

- **Network edge:** applications and hosts
- **Network core:**
 - routers
 - network of networks
- **Access networks**
 - wired
 - wireless



The network core

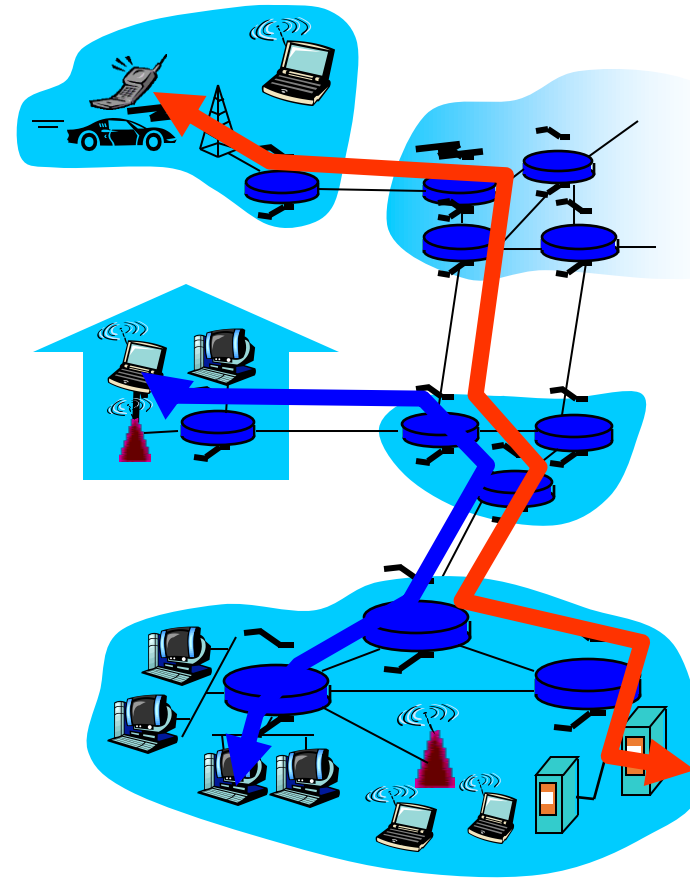
- Mesh of interconnected routers
- **Fundamental question:** how is data transferred through net?
- **circuit switching:** dedicated circuit per call: telephone net
- **packet-switching:** data sent thru net in discrete “chunks”



Network core: Circuit switching

End-end resources reserved for “call”

- Link bandwidth, switch capacity
- Dedicated resources: no sharing
- Circuit-like (guaranteed) performance
- Call setup required



Network core: Packet switching



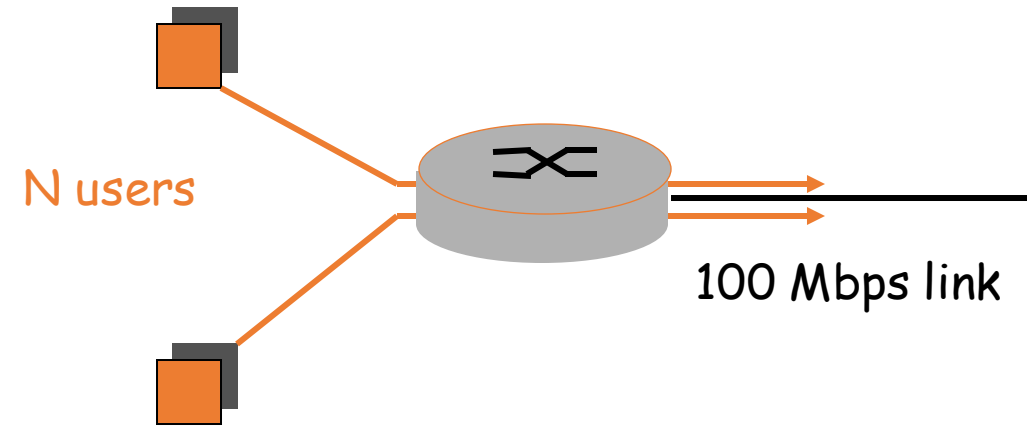
Each end-end data stream
divided into *packets*

- User A, B packets *share* network resources
- Each packet uses full link bandwidth
- Resources used *as needed*
- Resource contention
- Aggregate resource demand can exceed amount available
- Congestion: packets queue, wait to use link
- Store and forward: packets move one hop at a time
 - transmit over link
 - wait turn at next link

Packet switching versus circuit switching

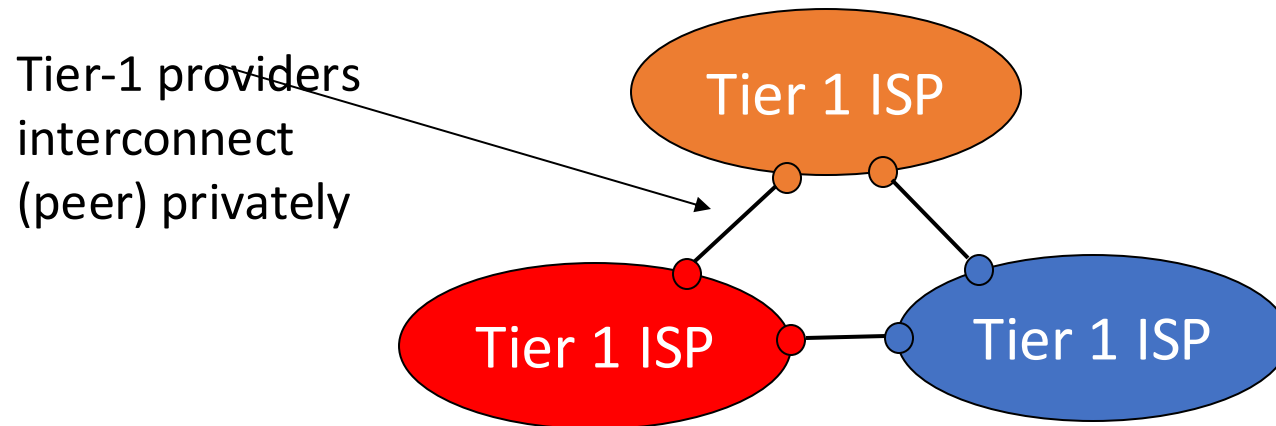


- 100 Mb/s link
- each user:
 - 10 Mb/s when “active”
 - active 10% of time
- Circuit-switching:
 - 10 users
- Packet switching:
 - with 35 users, probability > 10 active less than .0004



Packet switching allows more users to use network!

- Roughly hierarchical
- **At center: “tier-1” ISPs** (e.g., Verizon, Sprint, AT&T, Level 3), national/international coverage
- treat each other as equals

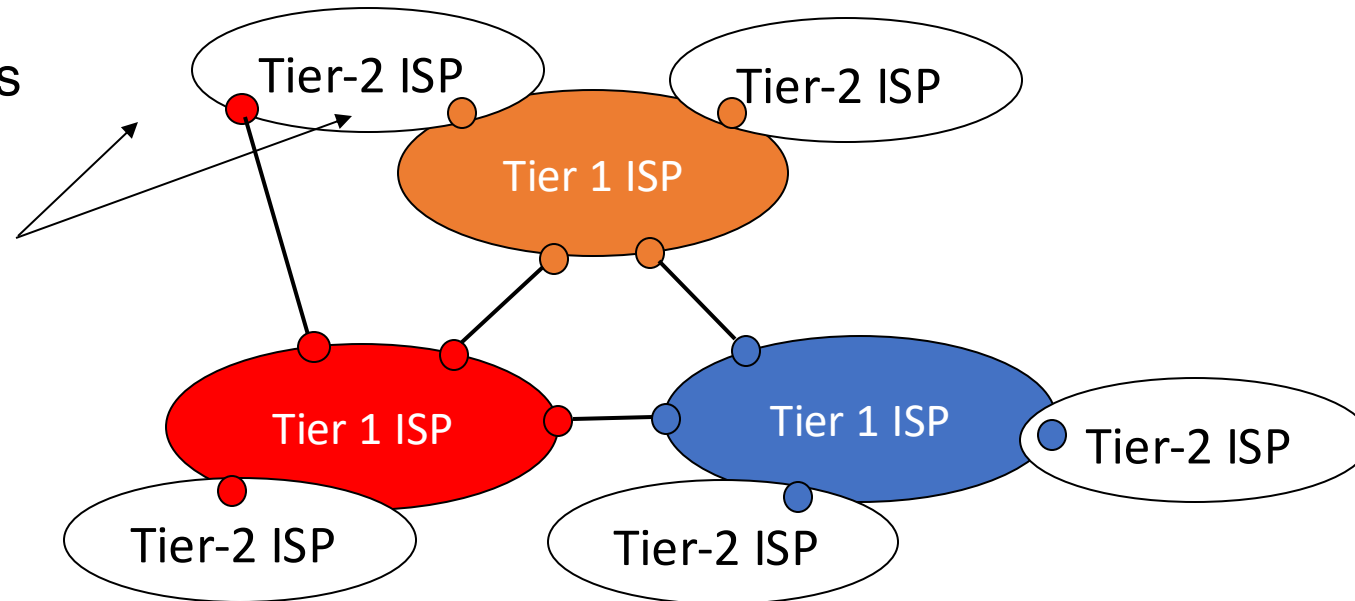


Internet structure: network of networks



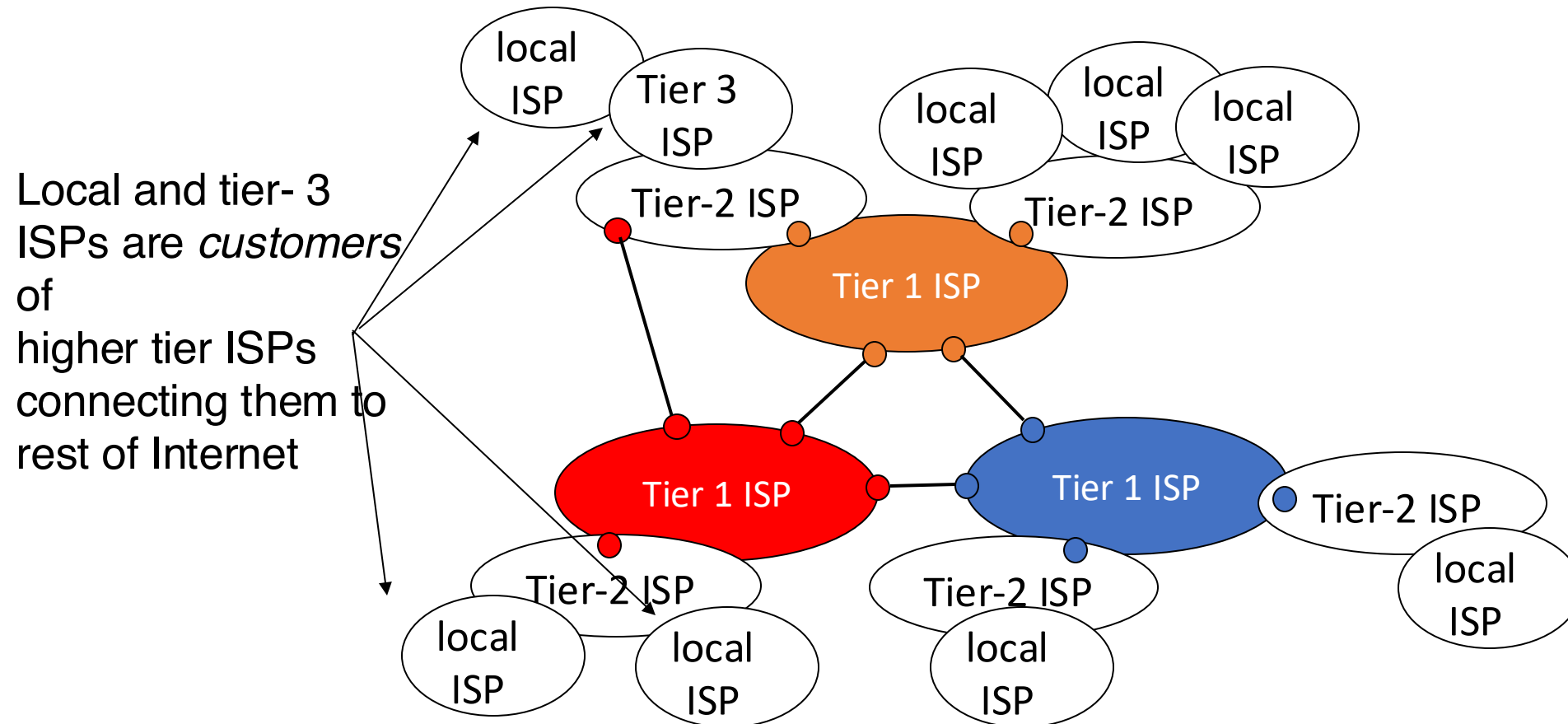
- “Tier-2” ISPs: smaller (often regional) ISPs
 - connect to one or more tier-1 ISPs, possibly other tier-2 ISPs

- tier-2 ISP pays tier-1 ISP for connectivity to rest of Internet
- tier-2 ISP is *customer* of tier-1 provider



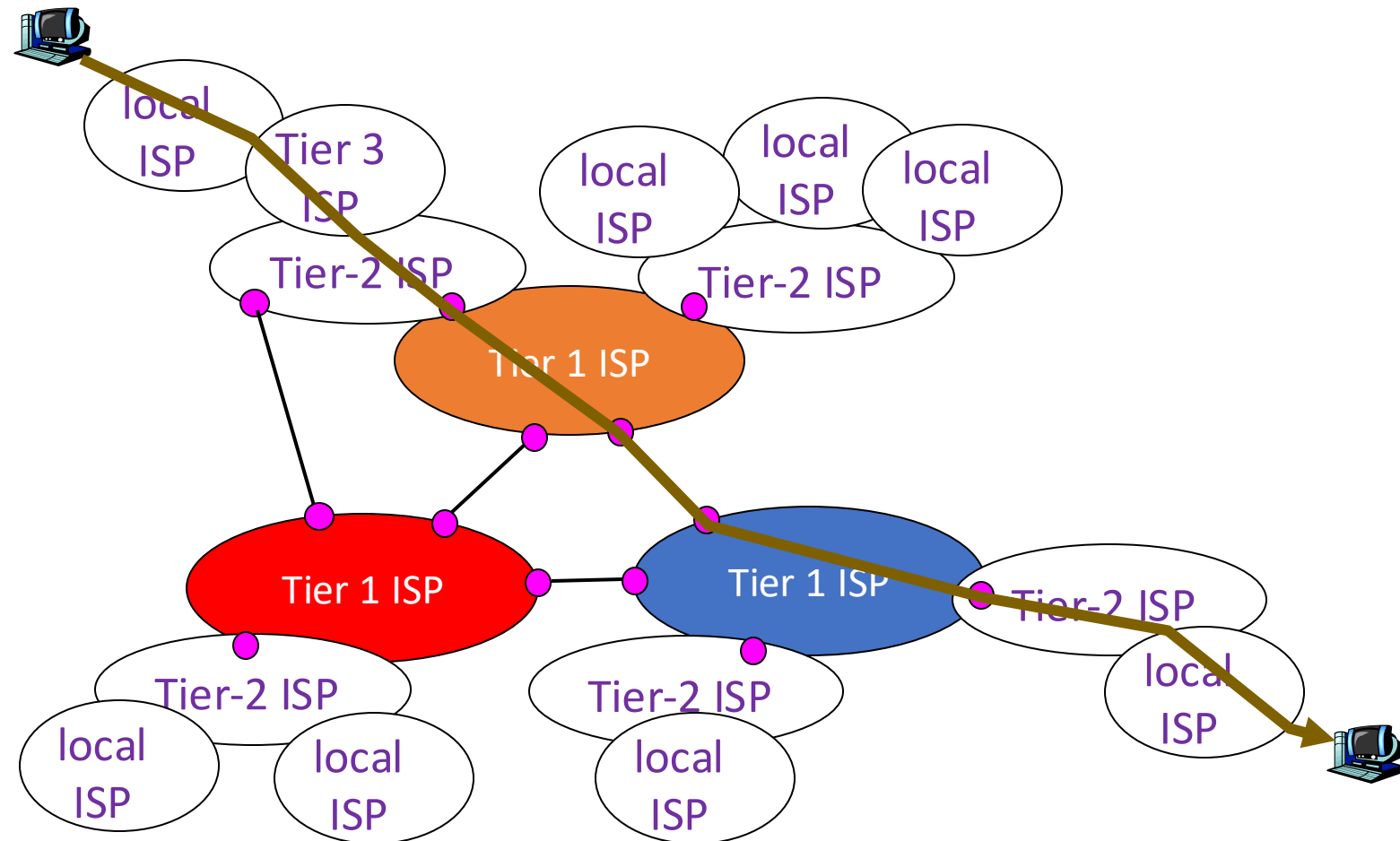
Internet structure: network of networks

- “Tier-3” ISPs and local ISPs
 - last hop (“access”) network (closest to end systems)



Internet structure: network of networks

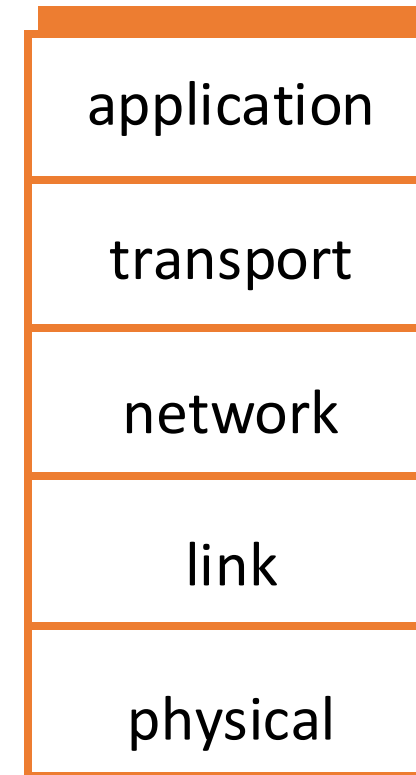
- a packet passes through many networks!



Internet protocol stack

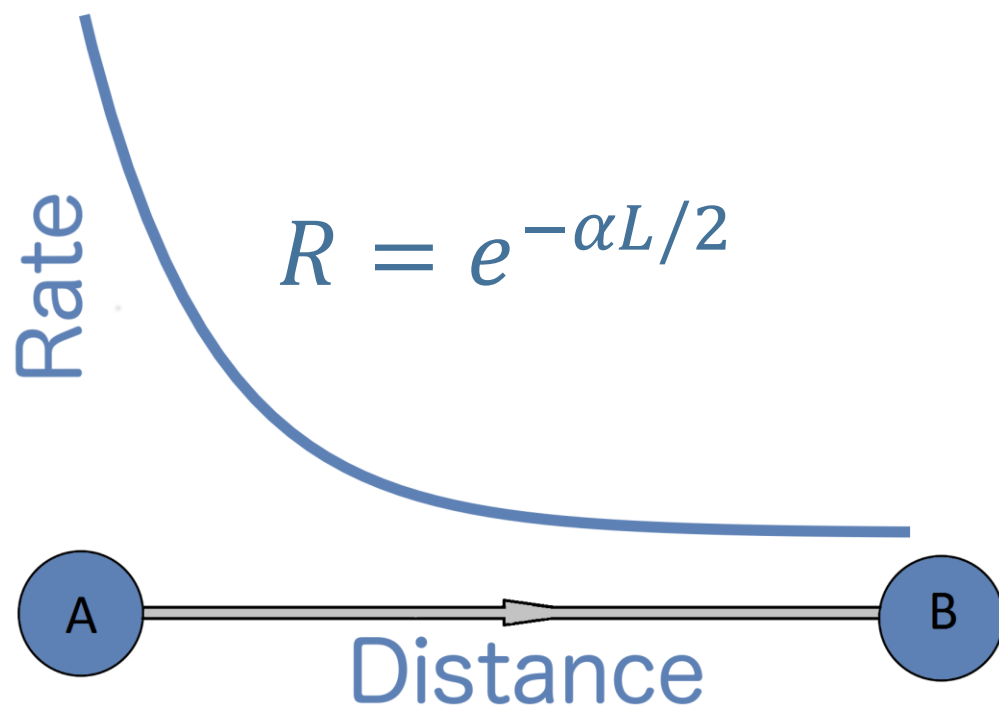


- **Application:** supporting network applications
 - scp, smtp, https
- **Transport:** host-host data transfer
 - tcp, udp
- **Network:** routing of packets from source to destination
 - ip, routing protocols
- **Link:** data transfer between neighboring network elements
 - ppp, ethernet
- **Physical:** bits “on the wire”



Quantum Networks

Why is quantum communications so hard?



**No cloning theorem
precludes copy and
amplification**

Rate decays exponentially
with distance

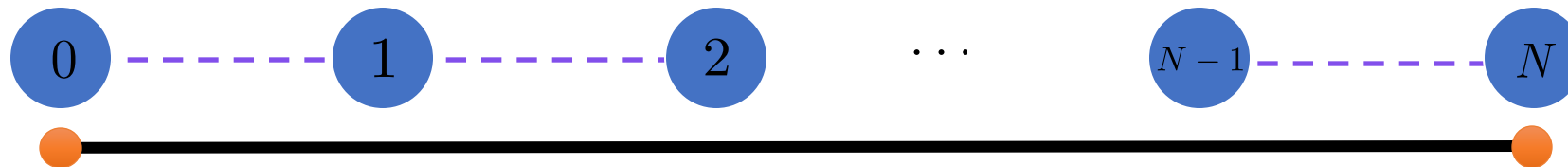
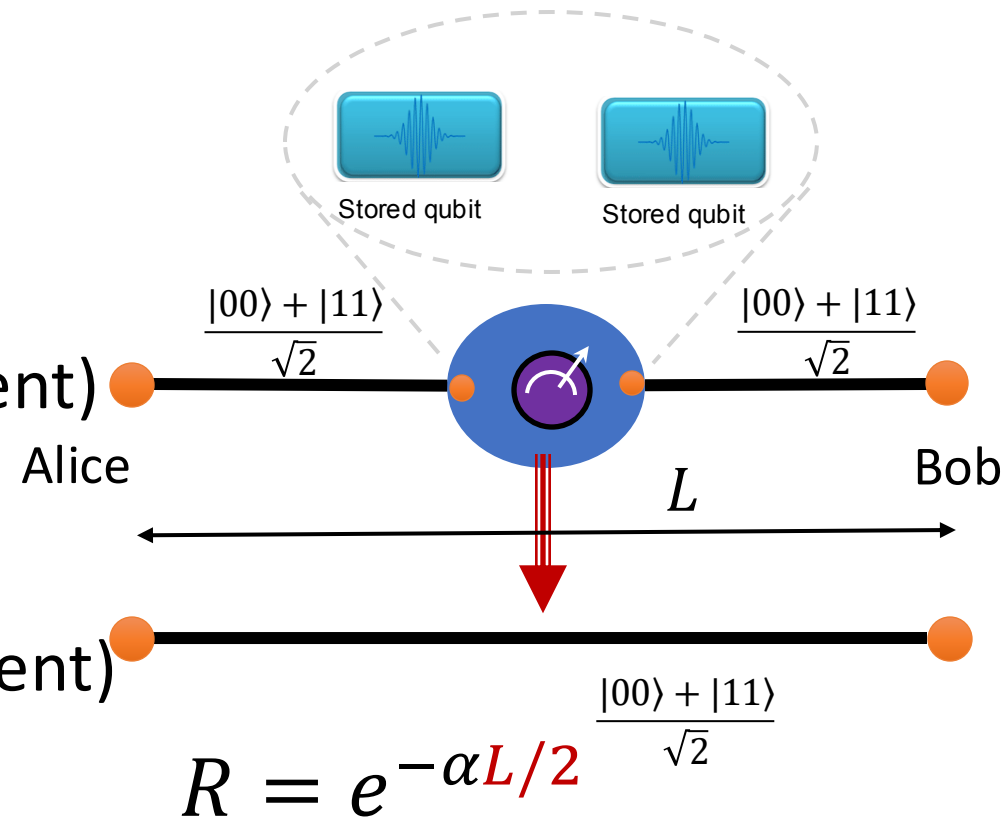
Quantum repeaters

Quantum memories to store qubits

Phase I: generate link Bell states (entanglement)

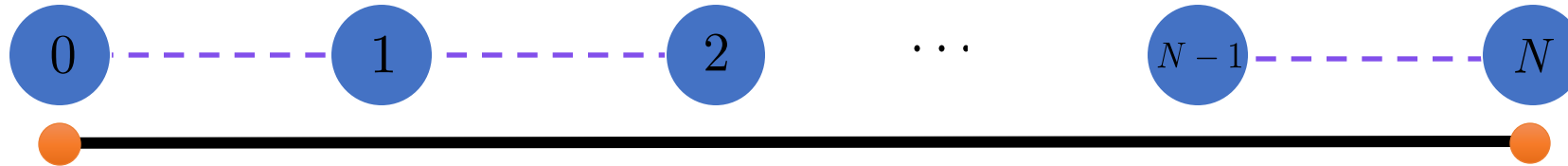
Phase II: propagate entanglements

entanglement swap (Bell state measurement)



$$R \propto e^{-\alpha L/N}$$

Repeater chain



- Infinite memory \Rightarrow distance independent entanglement rate

$$R \propto e^{-\alpha L/N}$$

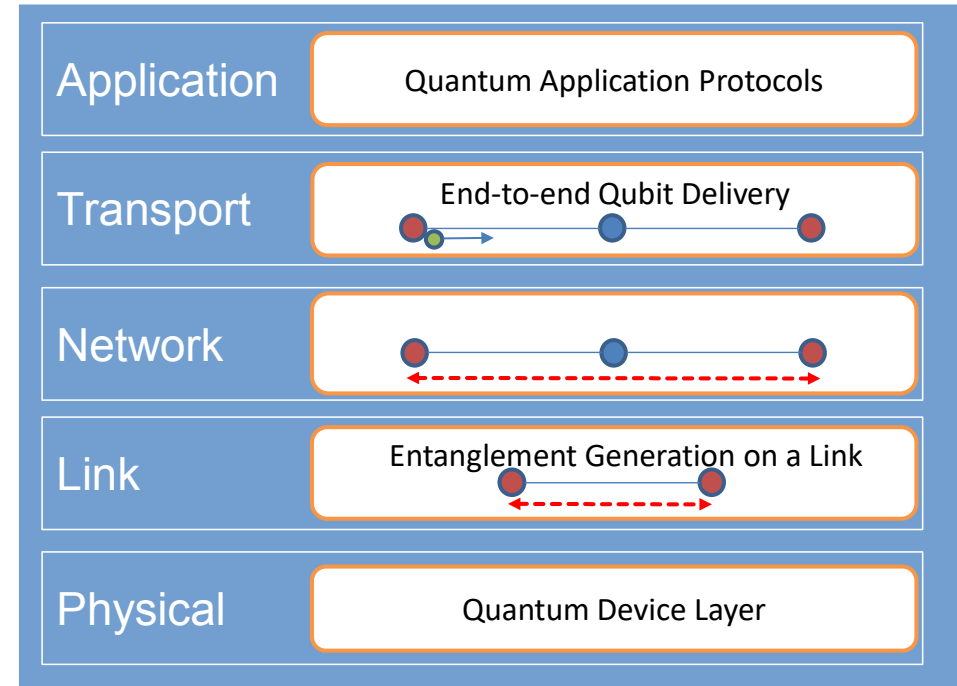
- Finite (one) memory \Rightarrow exponential decay in entanglement rate as function of L

$$R \propto e^{-\alpha L}$$

Quantum Internet



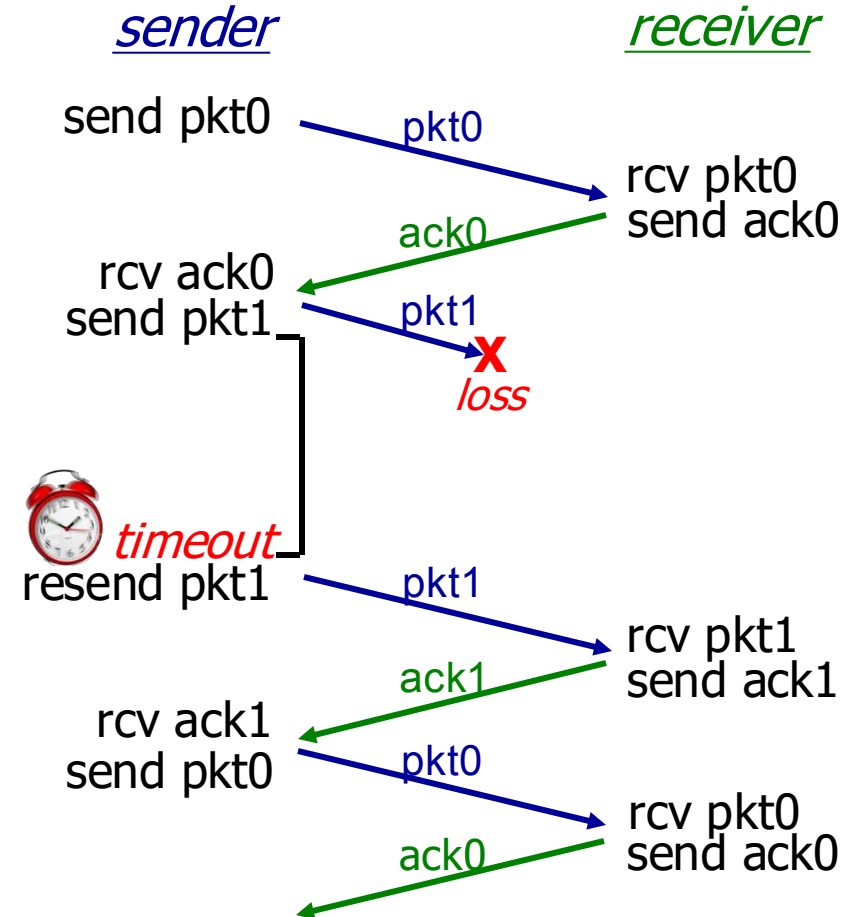
- **Application:** supporting network applications
- **Transport:** host-host quantum data transfer
 - qtcp, qudp
- **Network:** entanglement generation between end nodes
 - qip, path selection protocols
- **Link:** link-level entanglement generation
- **Physical:** photons “on the wire”



Stephanie Wehner et al.

Reliable communications (classical)

- Error models:
 - bit flips, erasures
 - *dropped packets*
- Recovery schemes
 - error detection/correction codes
 - packet retransmission
 - *relies on cloning!*



Reliable communications (quantum)

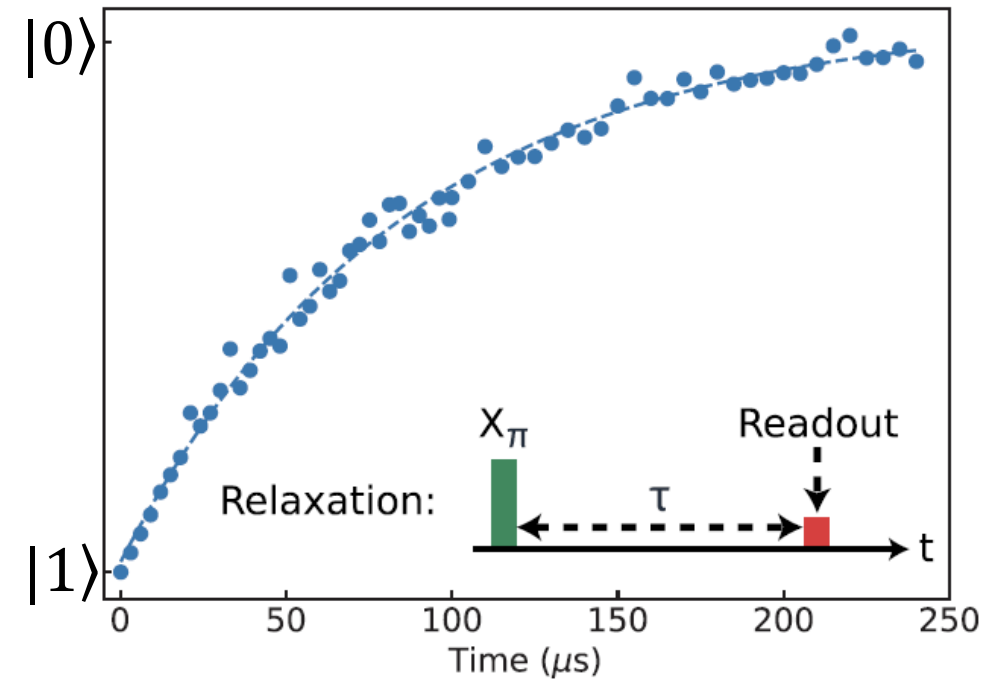
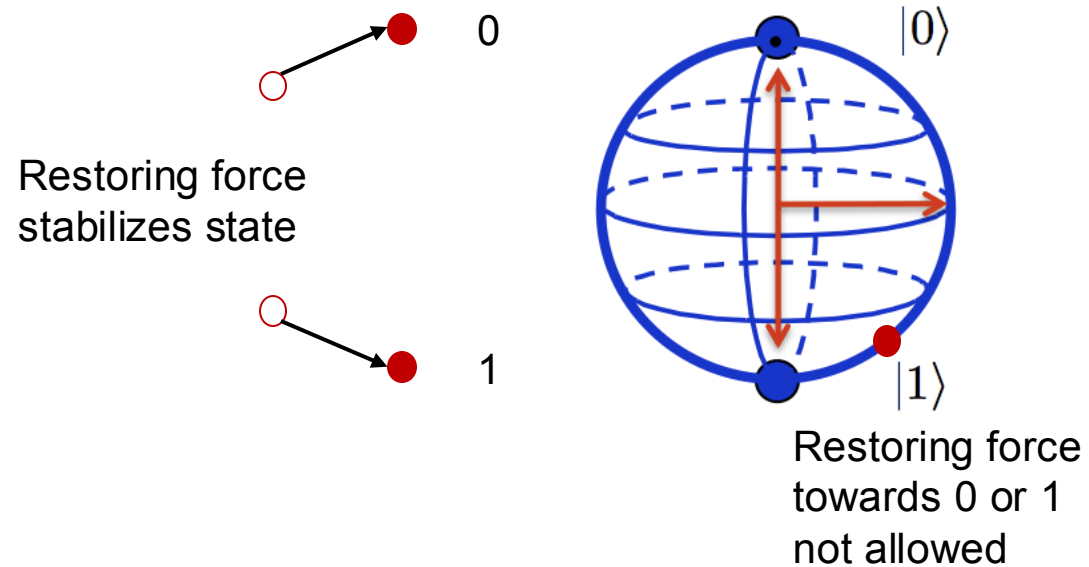


- Errors (non-exhaustive):
 - Models: Pauli channels, erasures, amplitude damping,...
 - Gate noise
 - Memory decoherence
 - Transduction
 - Loss in optical components
 - Collection efficiency
 - Probabilistic operations
- Recovery schemes
 - Quantum error correction (one-way)
 - Distillation
 - Data retransmission (one-way)
 - *Source must regenerate!*
 - Attempt until succeed (two-way)
 - *Two-way retransmission*

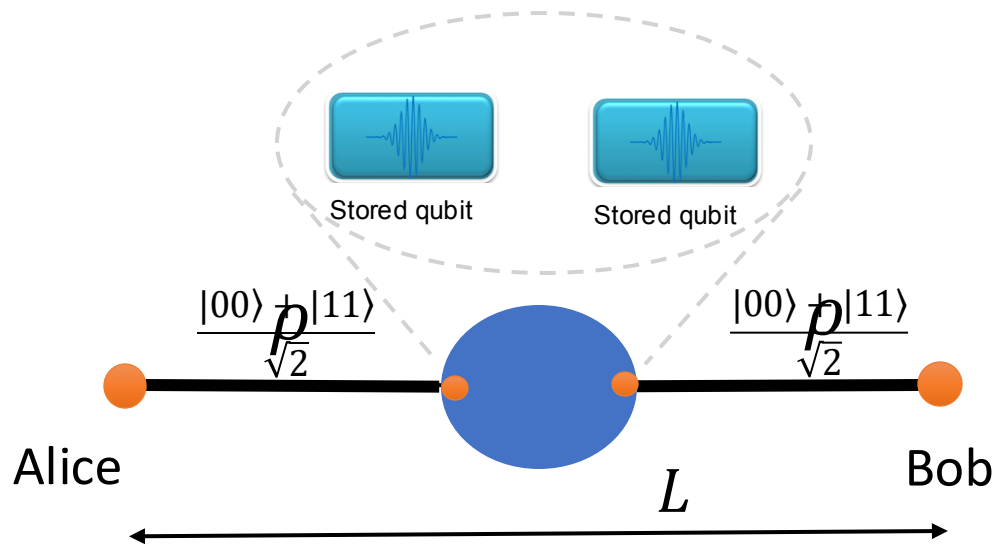
Quantum challenge

- Qubits not self protected against smallest perturbation

- Qubits have limited coherence times



Imperfect Entangled States



- Noise introduced thru comms
- Bell states decohere

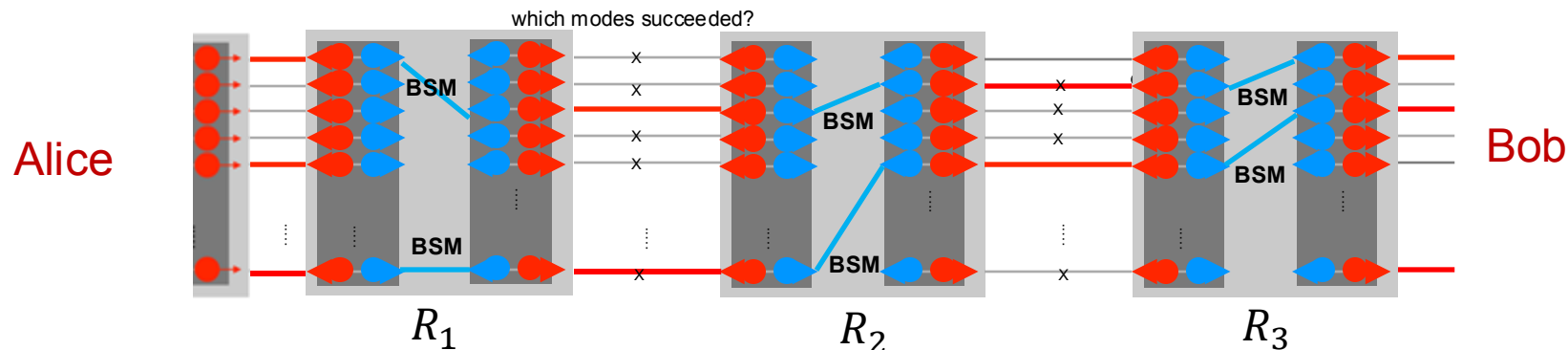
Fidelity: measure of closeness between two quantum states



Entanglement fidelity is fidelity between ρ and $\frac{|00\rangle + |11\rangle}{\sqrt{2}}$

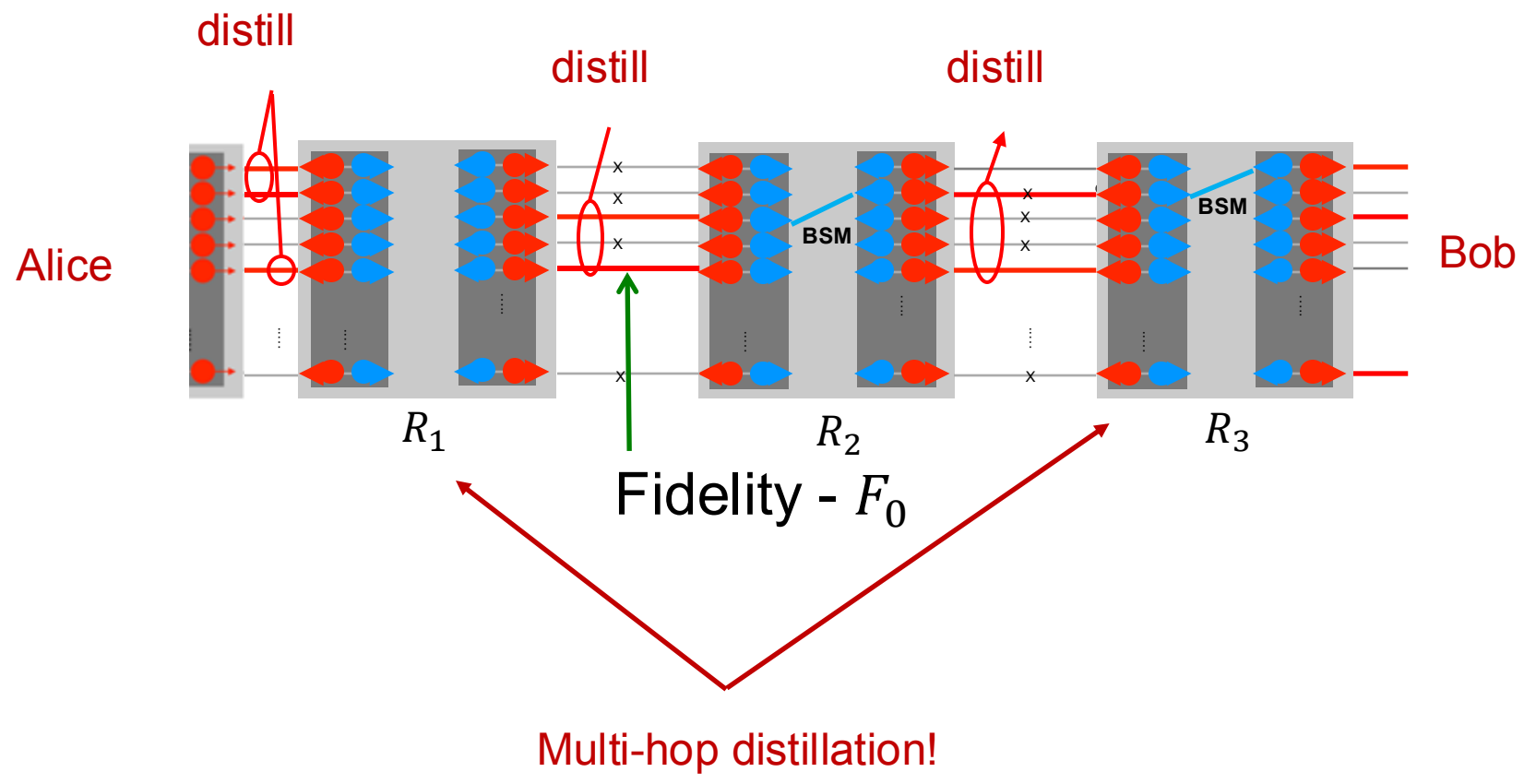
Multiplexed linear repeater network

- Parallel attempts performed at each link

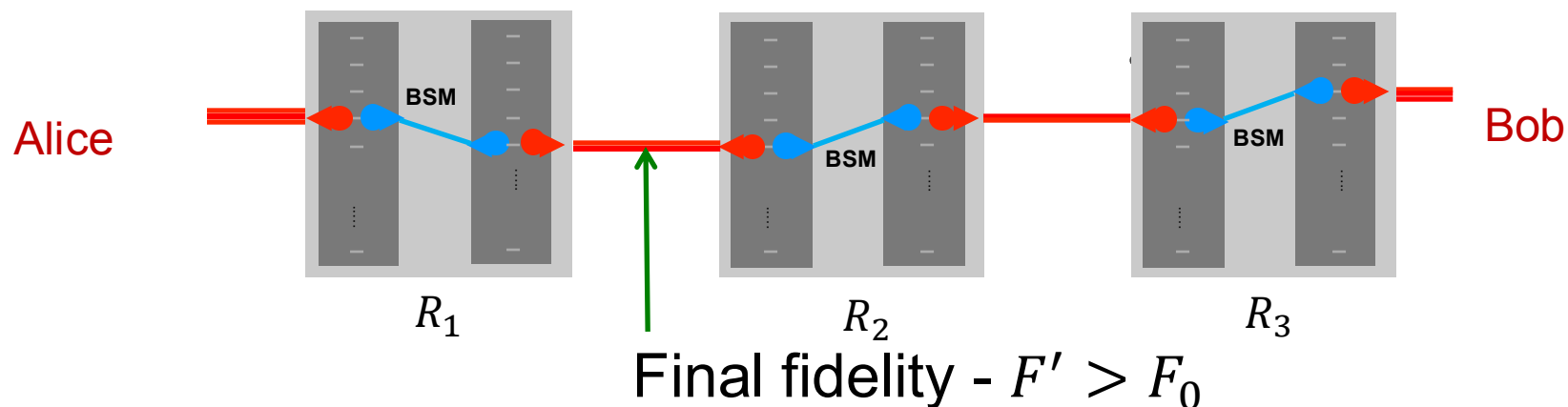


- Multiplexing increases rate
- Provides opportunity for distillation

Distillation



Distillation



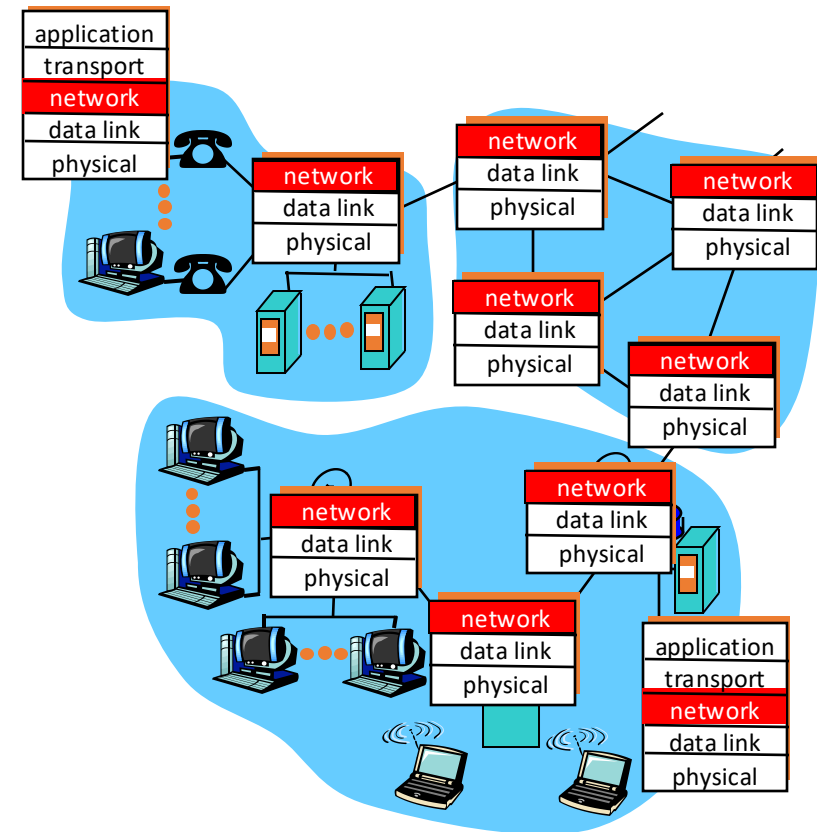
- Determine when and how much to distill
- Whether to distill across single or multiple links
- Possibly with minimum e2e fidelity constraint

Network layer functions

- Transport packet from sending to receiving hosts
- Network layer protocols in *every* host, router

Three important functions:

- *Path selection*: route taken by packets from source to destination (routing algorithms)
- *Switching*: move packets from router's input to appropriate router output
- *Call setup*: some network architectures require router call setup along path before data flows



Network service model



CRUCIAL
question!

Q: What *service model* for
“channel” transporting packets
from sender to receiver?

- guaranteed bandwidth?
- preservation of inter-packet timing (no jitter)?
- loss-free delivery?
- in-order delivery?
- congestion feedback to sender?

service abstraction

The most important
abstraction provided
by network layer:

virtual circuit
or
datagram?

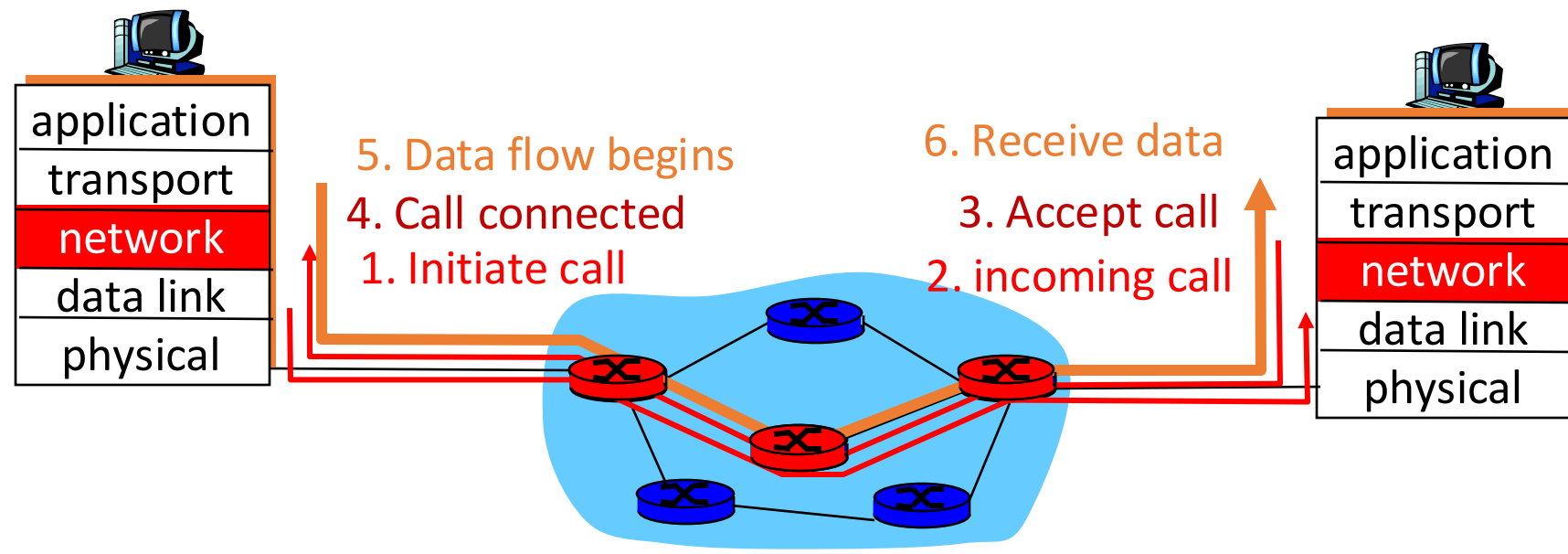
Virtual circuits



“source-to-dest path behaves like telephone circuit”

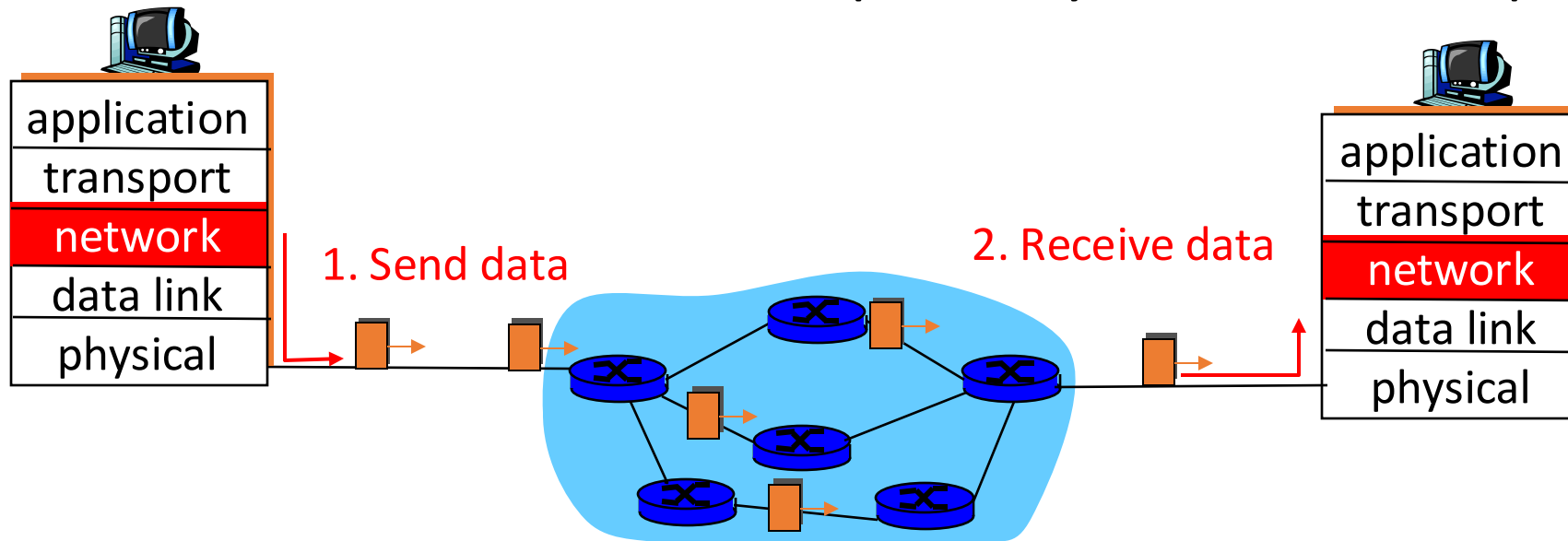
- performance-wise
 - network actions along source-to-dest path
-
- Call setup, teardown for each call *before* data can flow
 - Each packet carries VC identifier (not destination host ID)
 - *Every* router on source-dest path maintains “state” for each passing connection
 - transport-layer connection only involved two end systems
 - Link, router resources (bandwidth, buffers) may be *allocated* to VC
 - to get circuit-like performance

Virtual circuits in practice



Datagram network: The Internet model

- No call setup at network layer
- Routers: no state about end-to-end connections
 - no network-level concept of “connection”
- Packets typically routed using destination host ID
 - packets between same source-dest pair may take different paths



Quantum network service model



Q: What *service model* for
“quantum channel” between
end nodes?

service abstraction

- guaranteed rate?
- latency guarantee?
- minimum fidelity guarantee?

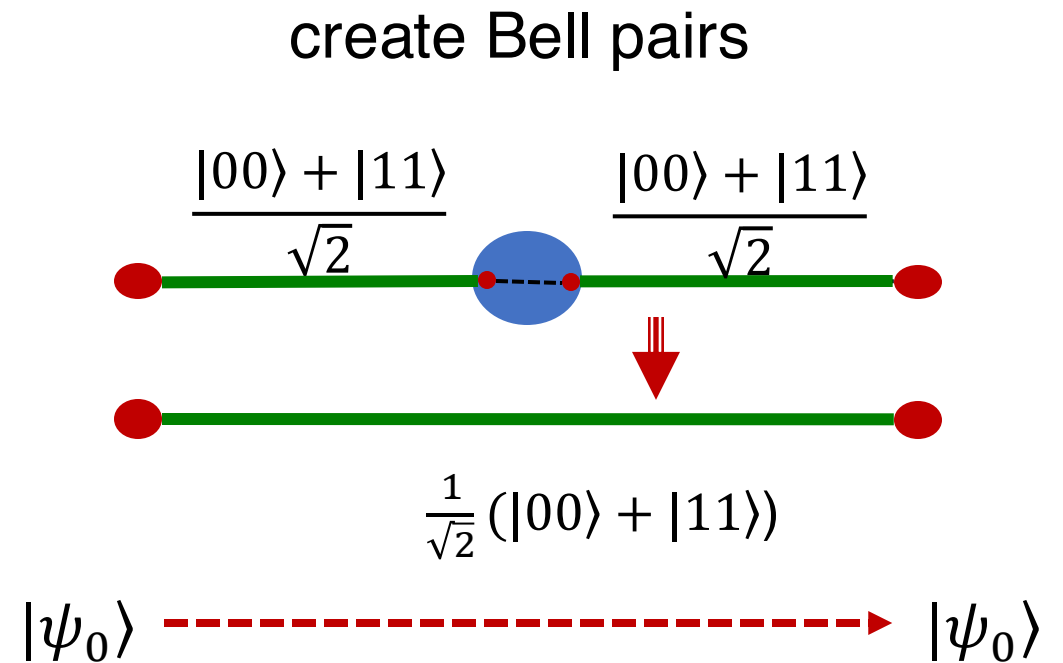
CRUCIAL
question!

The most important
abstraction provided
by network layer:

entanglement generation
or
quantum information transmission

Entanglement distribution (Two-way network architecture)

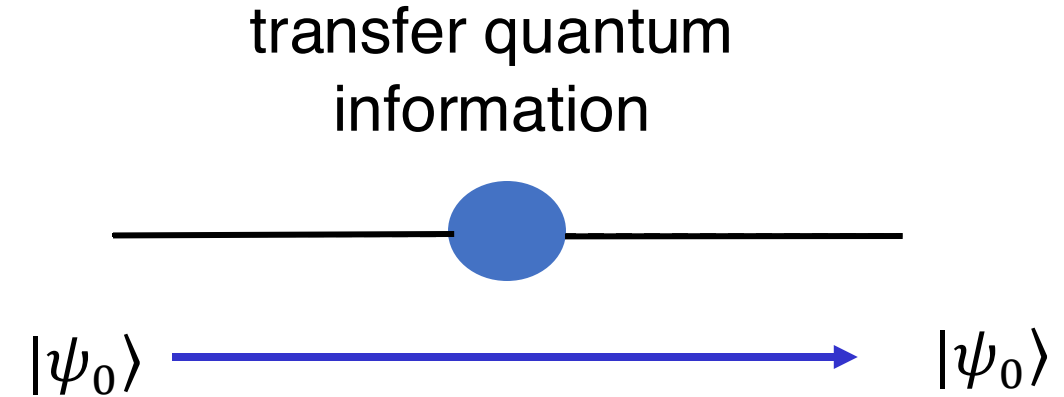
- Creation/distribution of Bell pairs (entanglement)
- Use **teleportation** to transfer quantum information
- Relies heavily on distillation to handle noise
- Requires exchange of classical information for correction



Quantum information transfer

(One-way network architecture)

- Transfer quantum information directly
- Note resemblance to classical network
- Relies heavily on Quantum Error Correction (QEC)
- Does not require exchange of classical info



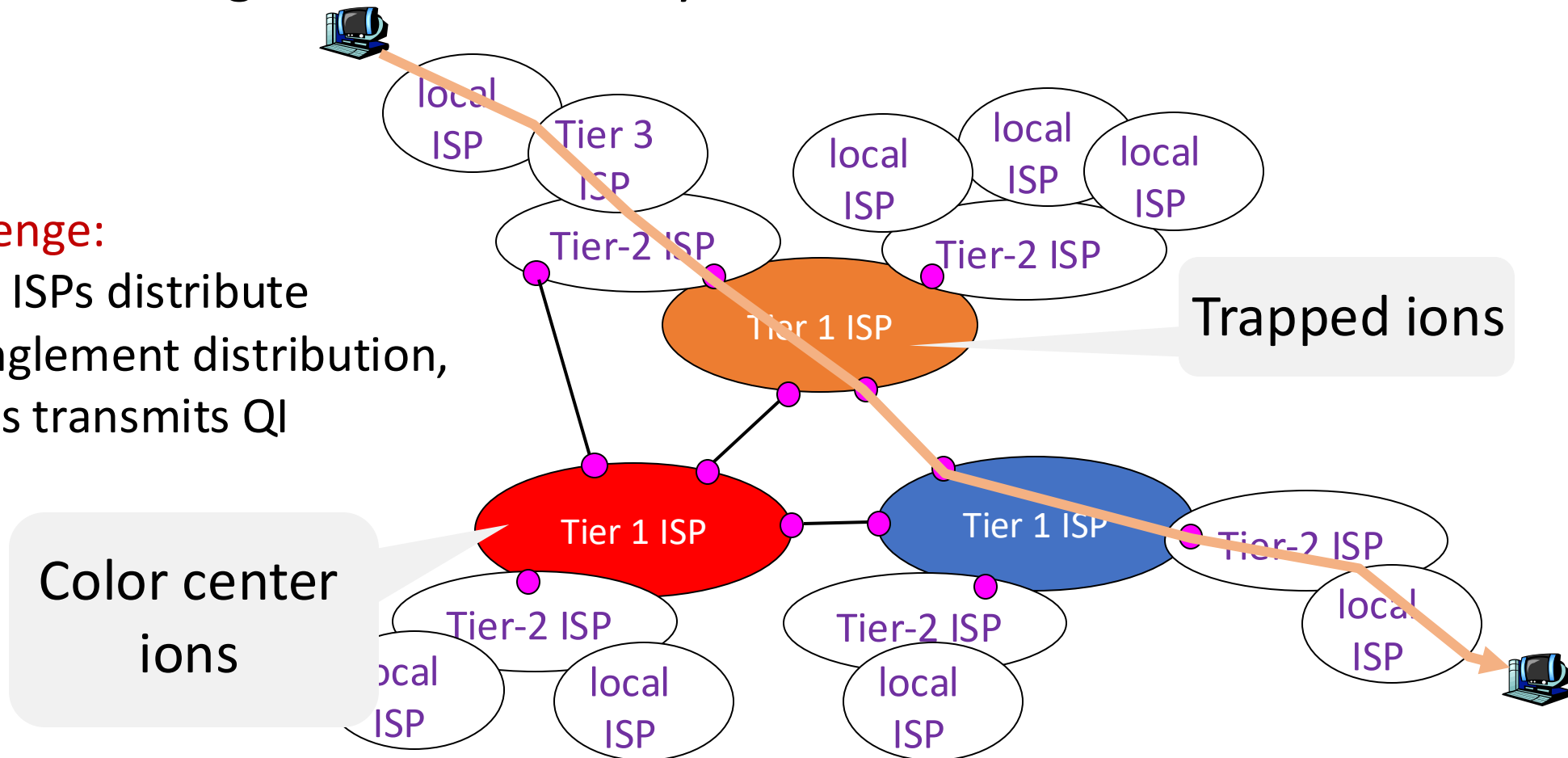
Note: services are interchangeable

Quantum Internet

- Quantum information can pass through many networks!
- e2e entanglement over many networks

Challenge:

some ISPs distribute entanglement distribution, others transmits QI



One way vs. Two way



Two way

Pros:

- Distillation simpler than QEC
- Bell pairs fungible \Rightarrow
 - high rates
 - pre-shared entanglement
- Tolerates noisy gates

Cons:

- Increased latency due to classical comms
- High memory requirement

One way

Pros:

- No classical comms \Rightarrow low latency
- Low memory requirement
- Allows for pre-sharing entanglement

Cons:

- QEC very challenging, requires high quality gates
 - 100 physical qubits per logical qubit?
- Requires many high quality gates
- May require more q-repeaters

Muralidharan, Sreraman, et al. "Optimal architectures for long distance quantum communication." *Scientific reports* 6.1 (2016): 20463.

Mantri, Prateek, Kenneth Goodenough, and Don Towsley. "Comparing One-and Two-way Quantum Repeater Architectures." *arXiv preprint arXiv:2409.06152* (2024).

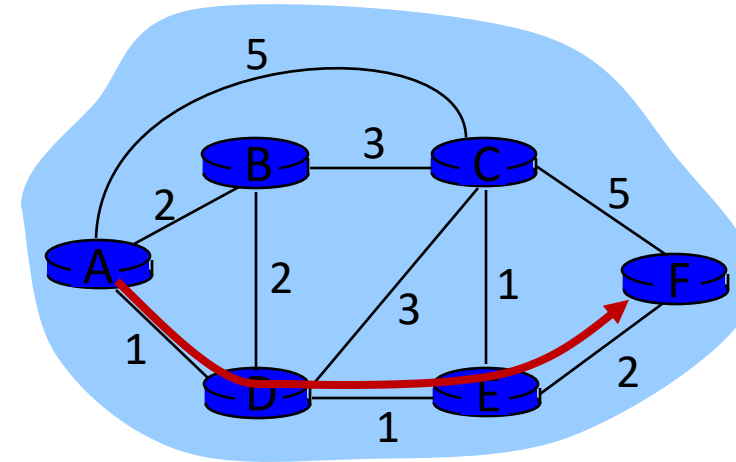
Classical routing

Routing protocol

Goal: determine “good” path
(sequence of routers) thru
network from source to dest.

Graph abstraction for routing
algorithms:

- graph nodes are routers
- graph edges are physical links
 - link cost: delay, \$ cost, or congestion level



“good” path:

- typically means minimum cost path
- other def’s possible
- Dijkstra algorithm

Routing algorithm classification

Q: global or decentralized information?

global:

- central controller has complete topology, link cost info

Decentralized:

- router knows physically-connected neighbors, link costs to neighbors
- iterative process of computation, exchange of info with neighbors
- “distance vector” algorithms

Q: static or dynamic?

static:

- routes change slowly over time

Dynamic:

- routes change more quickly
 - periodic update
 - in response to link cost changes

Current approach



- **(Logical)** central controller with complete topology, link cost info
- Includes policy constraints
 - e.g., party A cannot use link set \mathcal{L}
- Calculation of backup paths
- Diversity for load balancing

Static algorithms:

- shortest paths with link costs:
 - link entanglement rate, $1/R_l$
 - link fidelity, F_l
 - and others

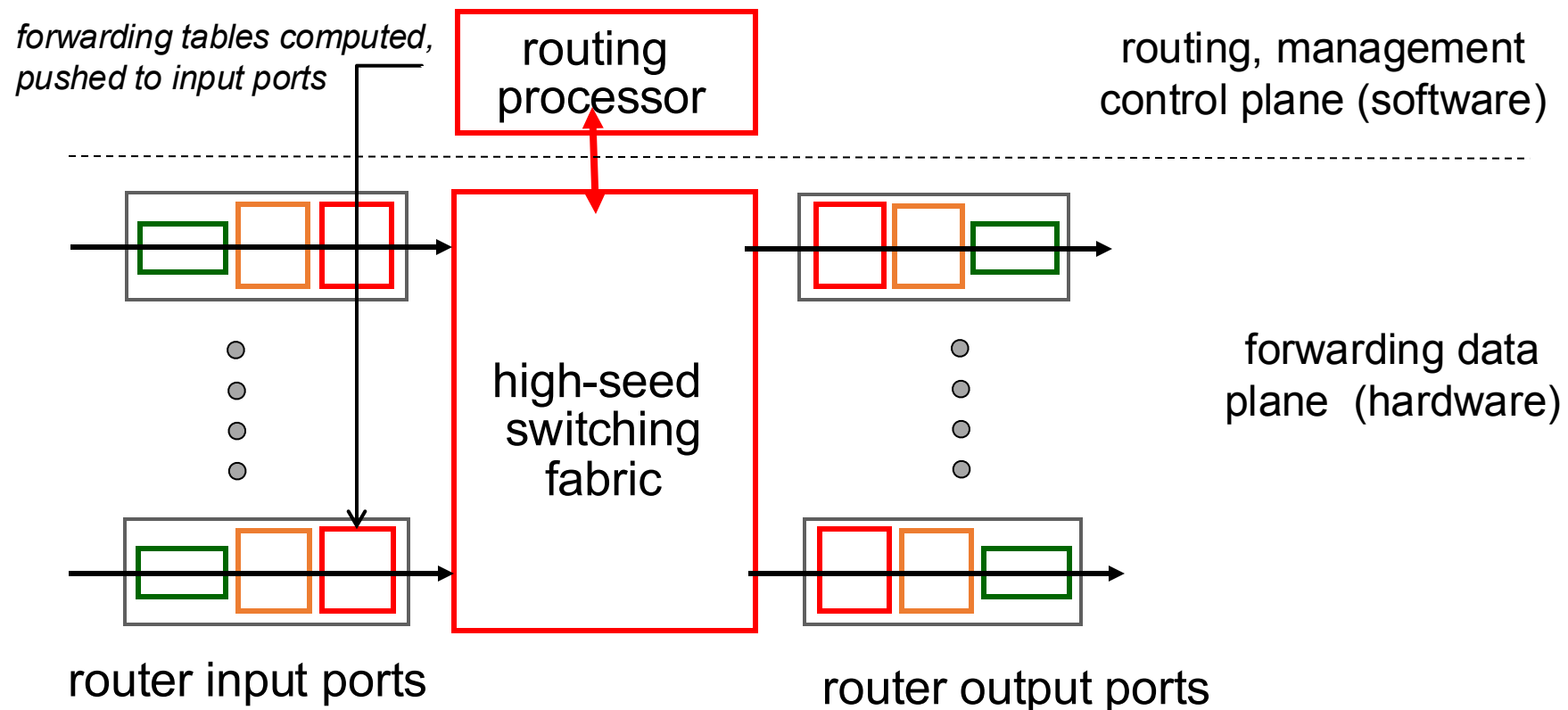
Dynamic algorithms:

- each node chooses neighbors to connect based on local state information

Classical router architecture overview

two key router functions:

- run routing algorithms/protocol
- *forwarding* packets from incoming to outgoing link



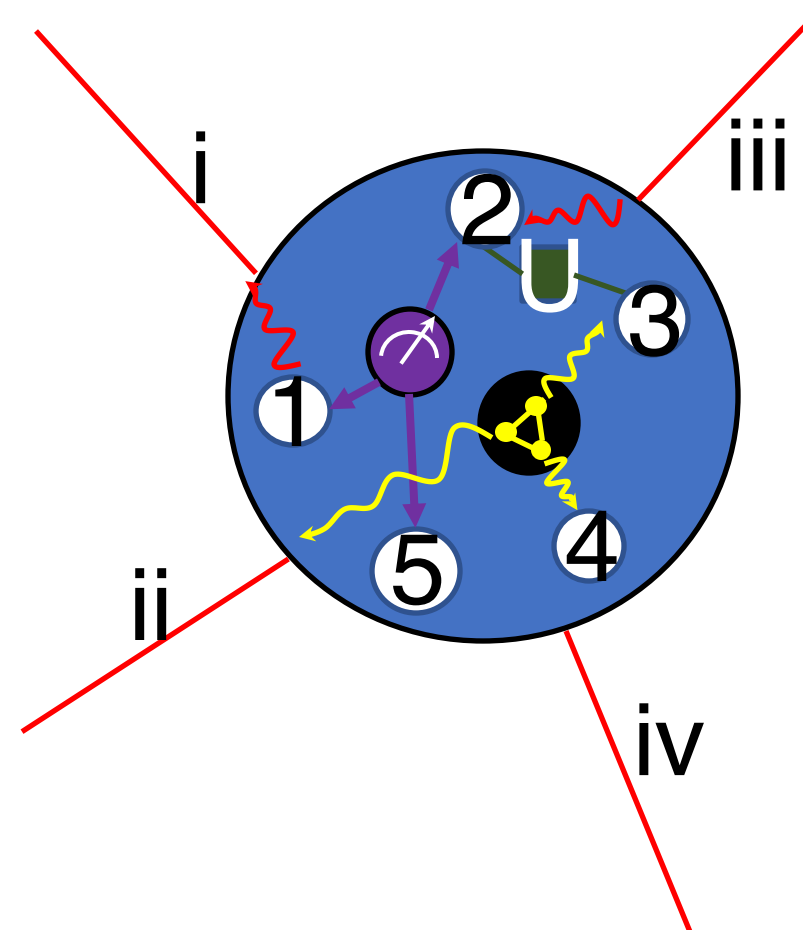
Questions



- capacity of router?
- scheduling policies that achieve capacity? that reduce switching fabric complexity?
 - matching algorithms
 - max weight policies
 - lightweight randomized algorithms

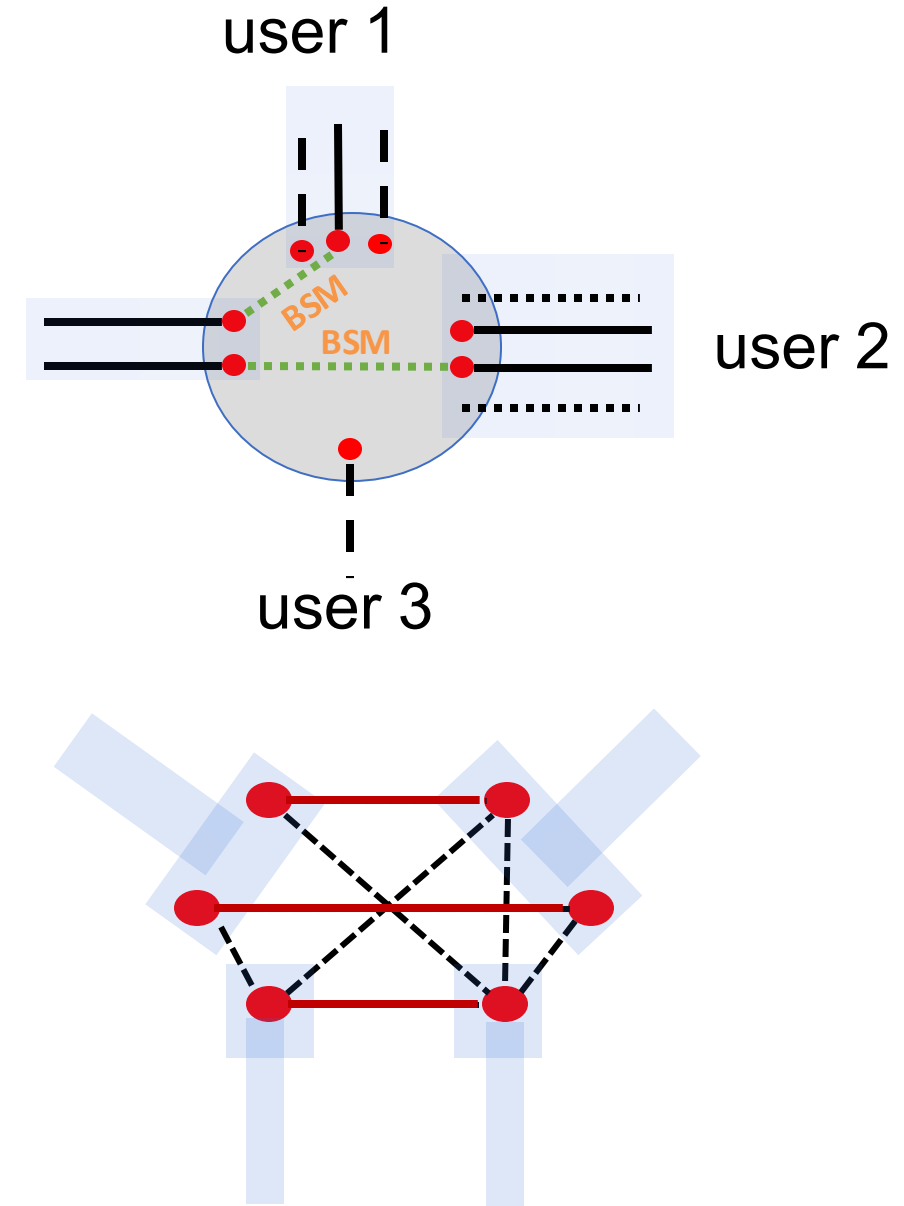
Two-way Quantum switch

- Quantum memories: loading and readout
- Multi-qubit quantum measurements
- Quantum logic across qubits held in QMs
- Multi-photon entanglement sources
- Classical processing and communications



Quantum switch

- User pairs generate requests for Bell pairs
- **Phase 1:** links randomly generate Bell pairs
- **Phase 2:** given outstanding requests, switch selects Bell pairs to measure
 - equivalent to selecting eligible matching in a graph among memories
- Outcomes of BSM matchings form set of end-to-end entanglements between pairs of end nodes



Challenges



- switch design, switching fabric
 - teleportation fabric?
- network capacity, network resource allocation
 - global vs local vs no state information
 - timescale of state information
- memory decoherence, gate errors?
- quality of information – fidelity
 - fidelity degrades over time \Rightarrow youngest qubit first (YQF), deadline scheduling?

\Rightarrow (virtual) circuit switching?

Summary



- entanglement distribution service very different from quantum information transfer service
- quantum networking introduces new problems
 - ... and old problems with new wrinkles
 - resource allocation, path selection, switch & entanglement scheduling
 - delivery of QoS in very noisy environment

Scheduling in repeater chains



Focus on
two-way

Distribution over a chain

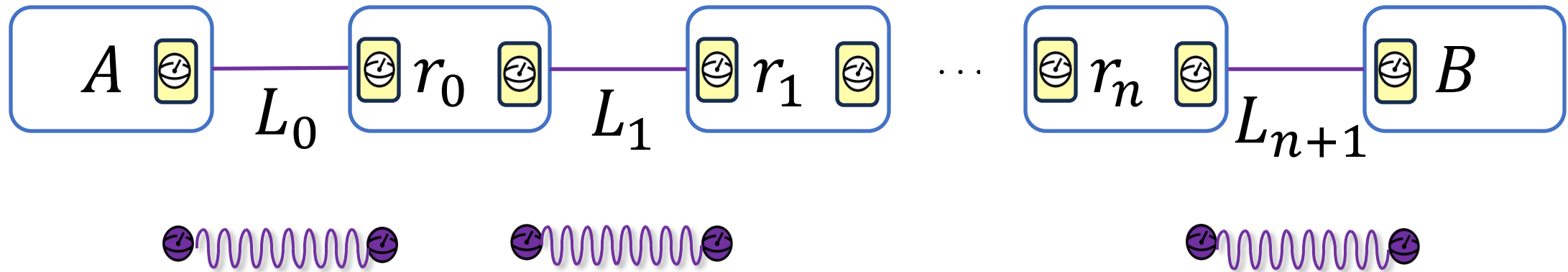


How to generate entanglement between end-nodes?

When to generate link-level entanglement?

When to perform entanglement swapping?

Swaps commute



Assume all LLEG succeeded

BSMs commute with each other

Any swapping ordering distributes e2e
entanglement!

Swap can be performed in parallel!

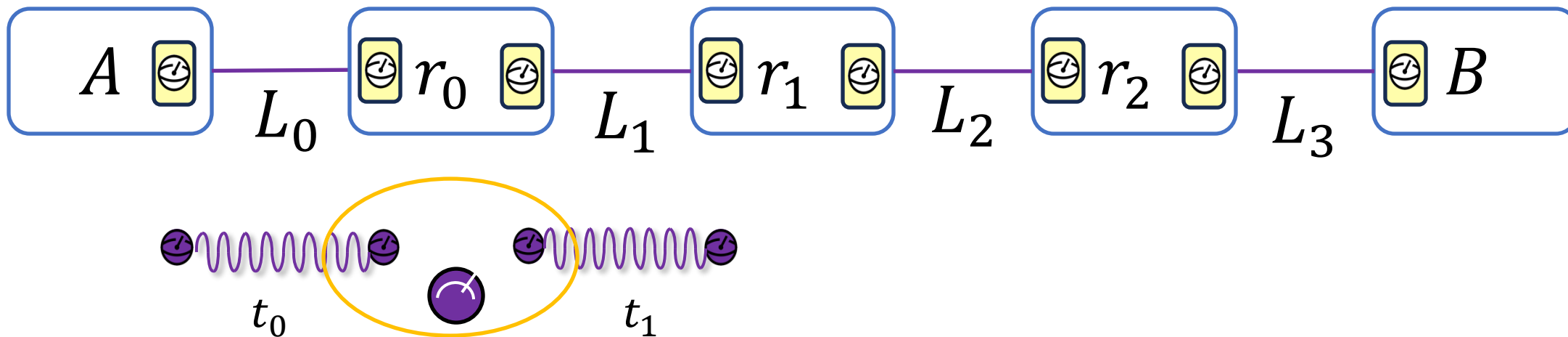
Swap order has profound impacts for
performance, architecture and protocol
design!

SWAP Policies

1. Sequential Swap
2. Swap As Soon As Possible (ASAP)
3. Nested swap

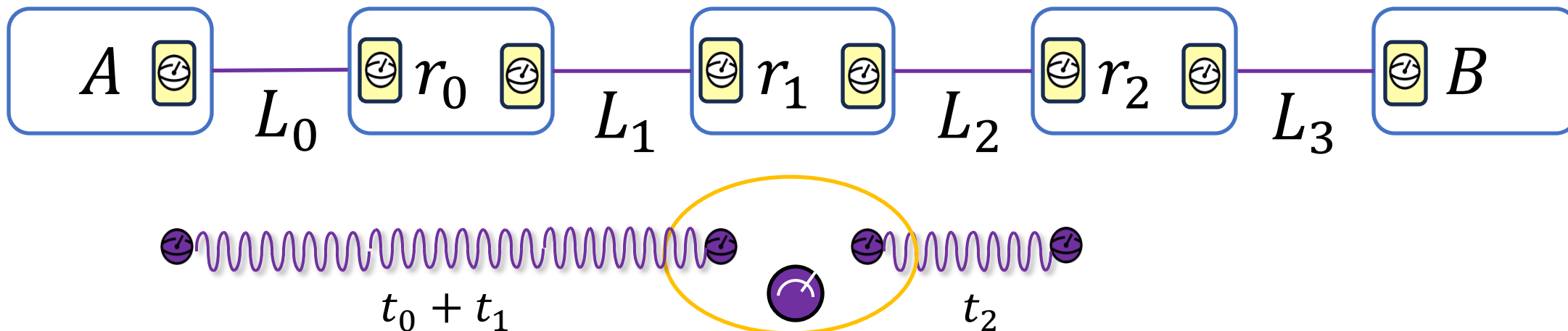
Assume deterministic swaps

Sequential Swap



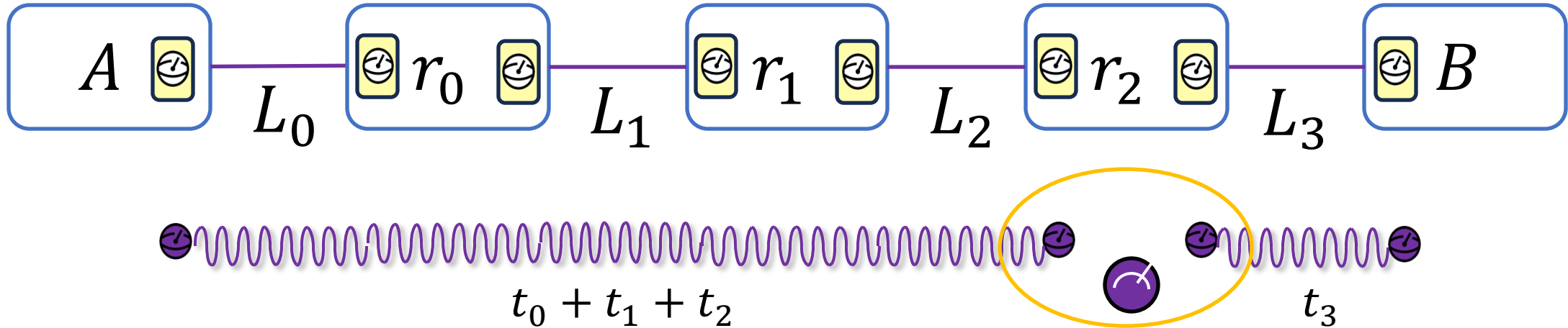
1. A starts LLEG r_0
2. Once succeeded, r_0 starts LLEG with r_1
3. BSM performed in r_0

Sequential Swap



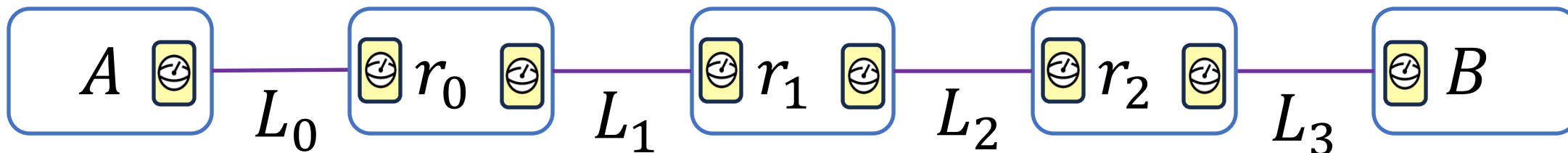
1. A starts LLEG r_0
2. Once succeeded, r_0 starts LLEG with r_1
3. BSM performed in r_0
4. r_i repeats the process until B reached

Sequential Swap



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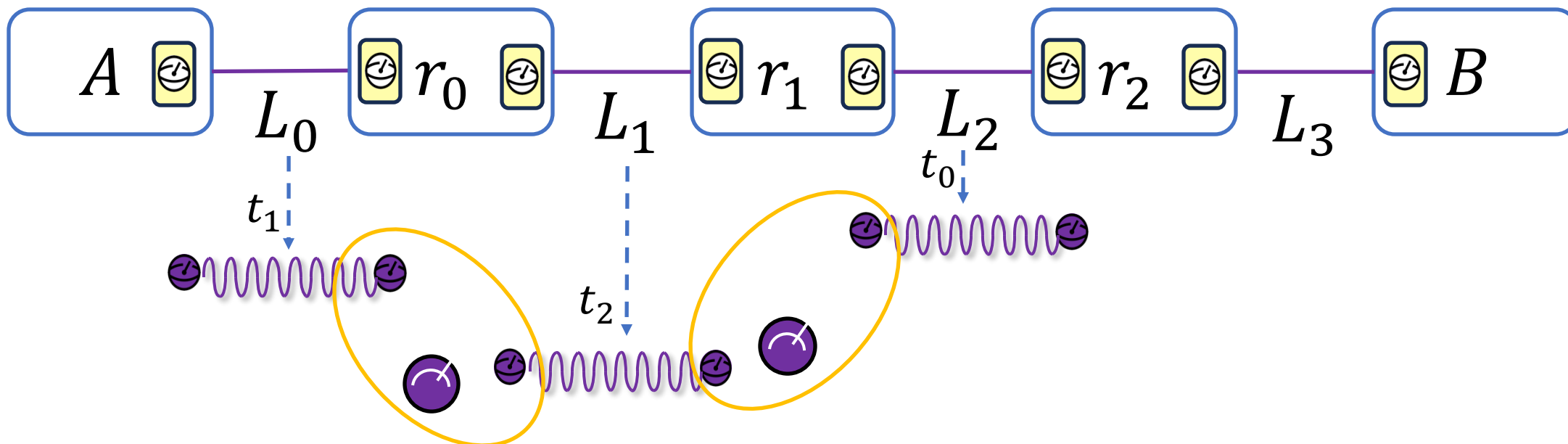
Sequential Swap



$$t_0 + t_1 + t_2 + t_3$$

1. A starts LLEG r_0
2. Once succeeded, r_0 starts LLEG with r_1
3. BSM performed in r_0
4. r_i repeats the process until B reached
5. Latency is $O(n)$ where n is number of repeaters

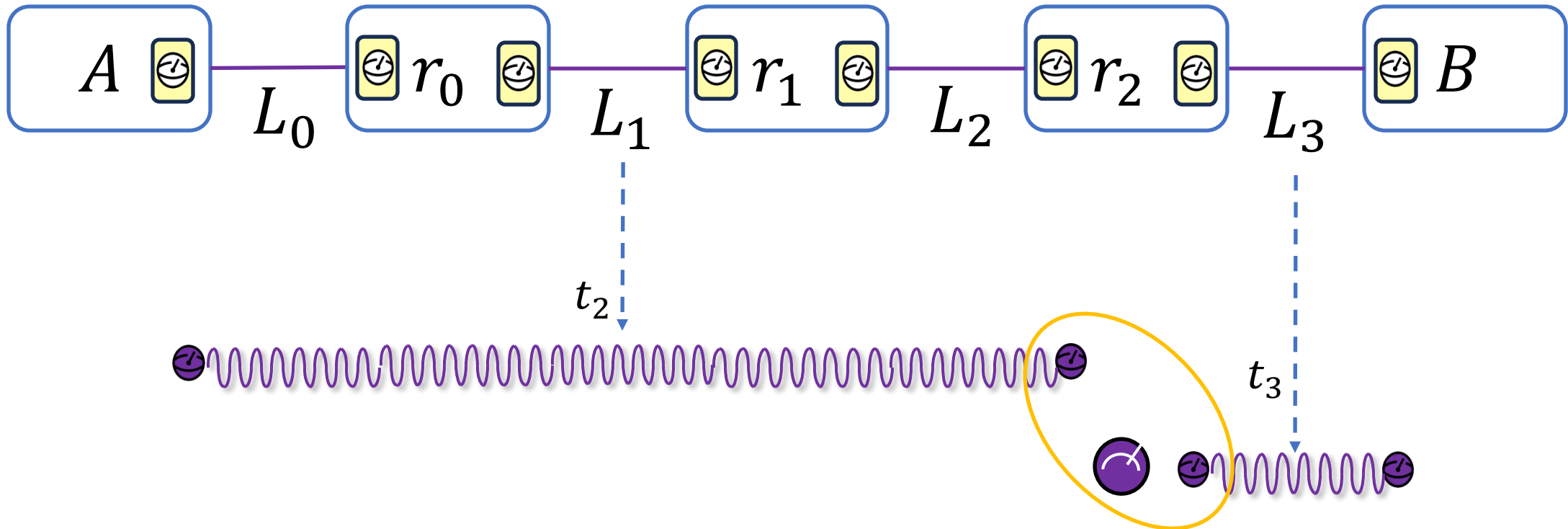
Swap ASAP



$$t_2 \geq t_1 \geq t_0$$

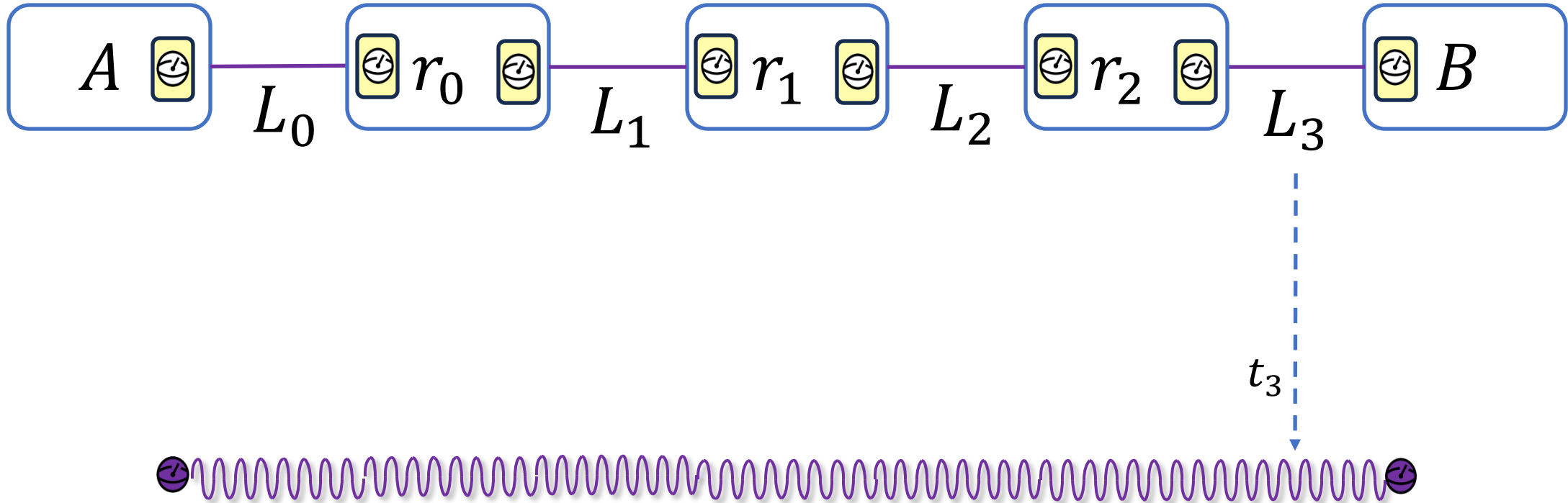
1. LLEG starts simultaneous on every link
2. Perform swap as soon as two links available
3. Stop when all repeaters swap

Swap ASAP



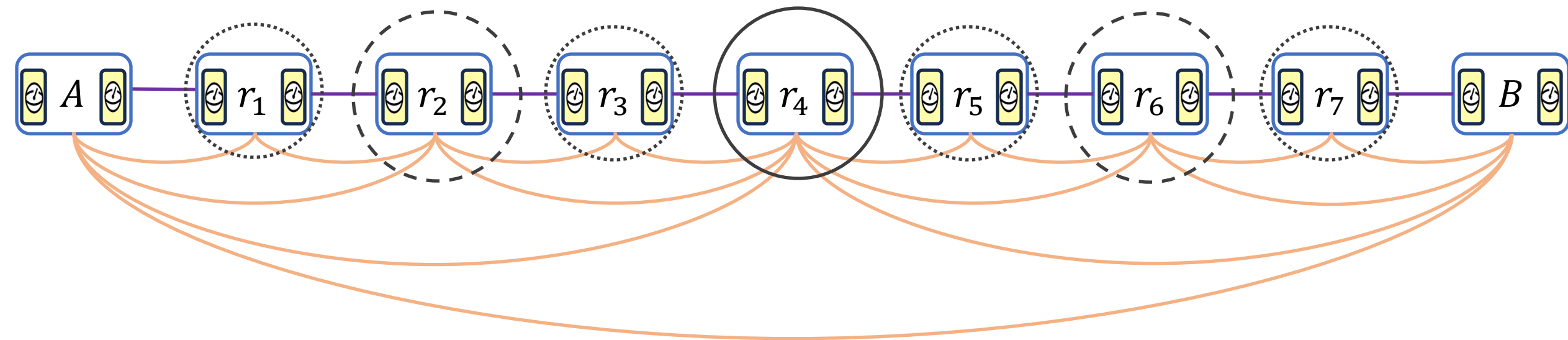
1. LLEG starts simultaneous on every link
2. Perform swap as soon as two links available
3. Stop when all repeaters swap

Swap ASAP



1. LLEG starts simultaneous on every link
2. Perform swap as soon as two links available
3. Latency is $O(\log n)$, n is number of repeaters

Nested Swaps



1. Attempt LLEG in parallel for all links
2. Find repeater at midpoint
3. Swap when both long range links have been created
4. Repeat process until possible



Analysis

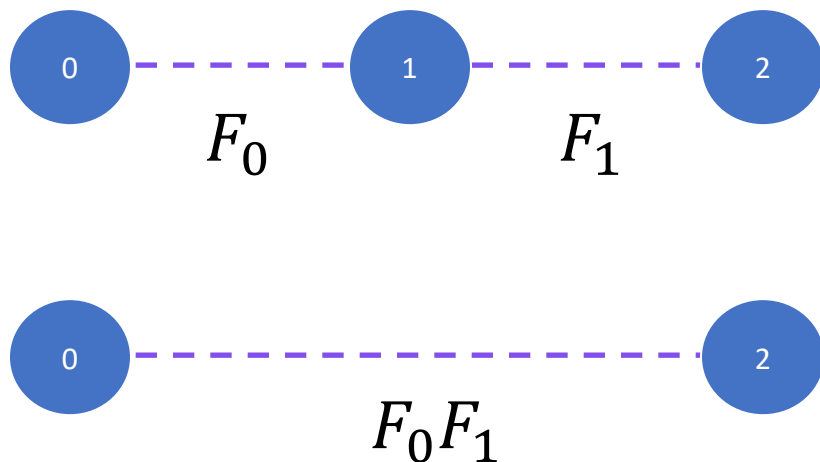
- Swap schedules have different performance
- If swap probability q , decay rate changes drastically
- Sequential simpler to implement, follow classical nets
- ASAP and Nested may require additional control

	Sequential	ASAP	Nested
Probability scaling with n	q^n	q^n	$q^{\log(n)}$
# of Bell pairs decohering	1	multiple	multiple
Latency	$O(n)$	$O(\log(n))$	$O(\log(n))$

Distribution and Noise

What happens with imperfect entanglement?

Assume that noise is depolarizing
(worst case scenario)



Noise accumulates exponentially!

Distillation also requires scheduling!

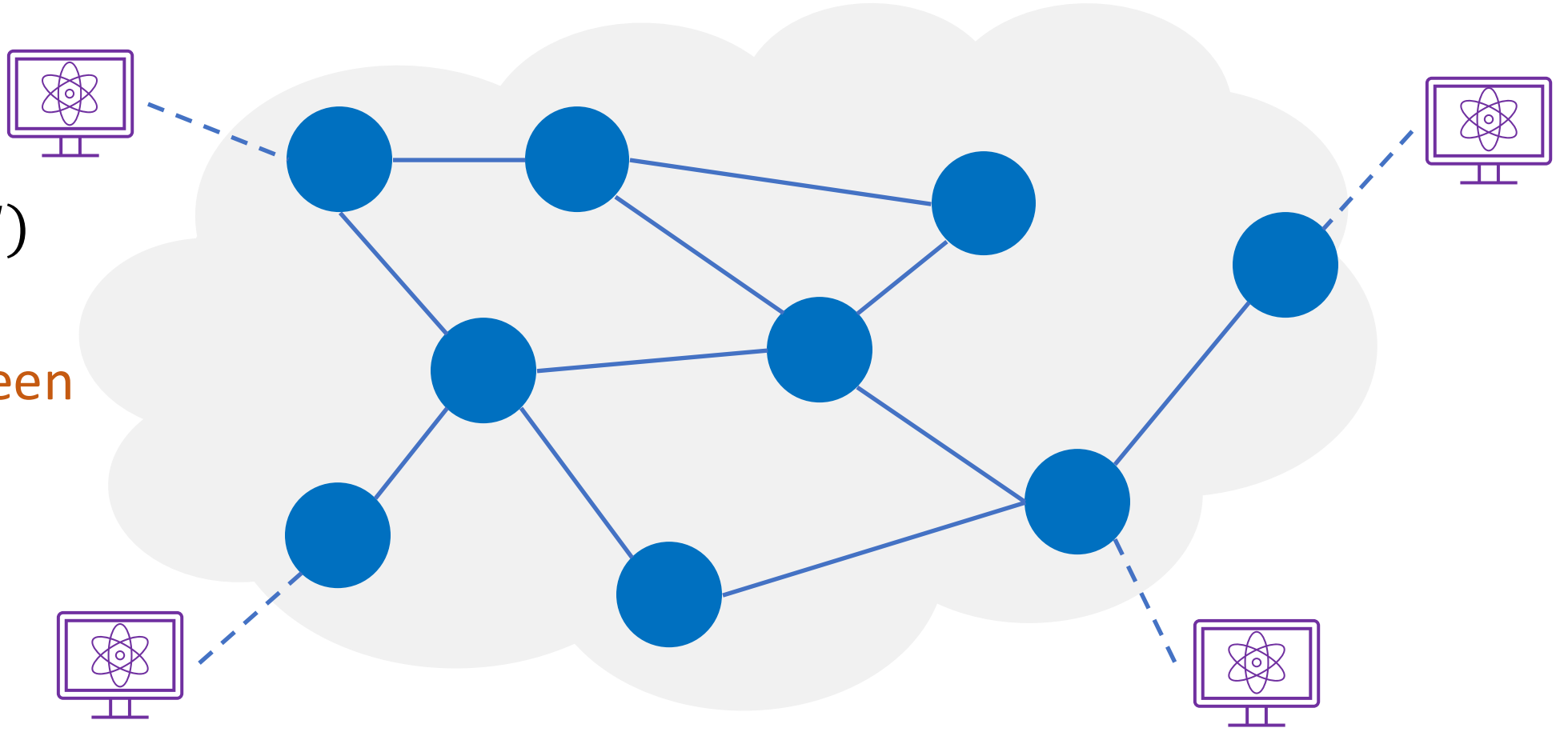
Different distillation schedules
exhibit different performance!

Quantum Network Routing

The Routing Problem

Given $G = (V, E)$

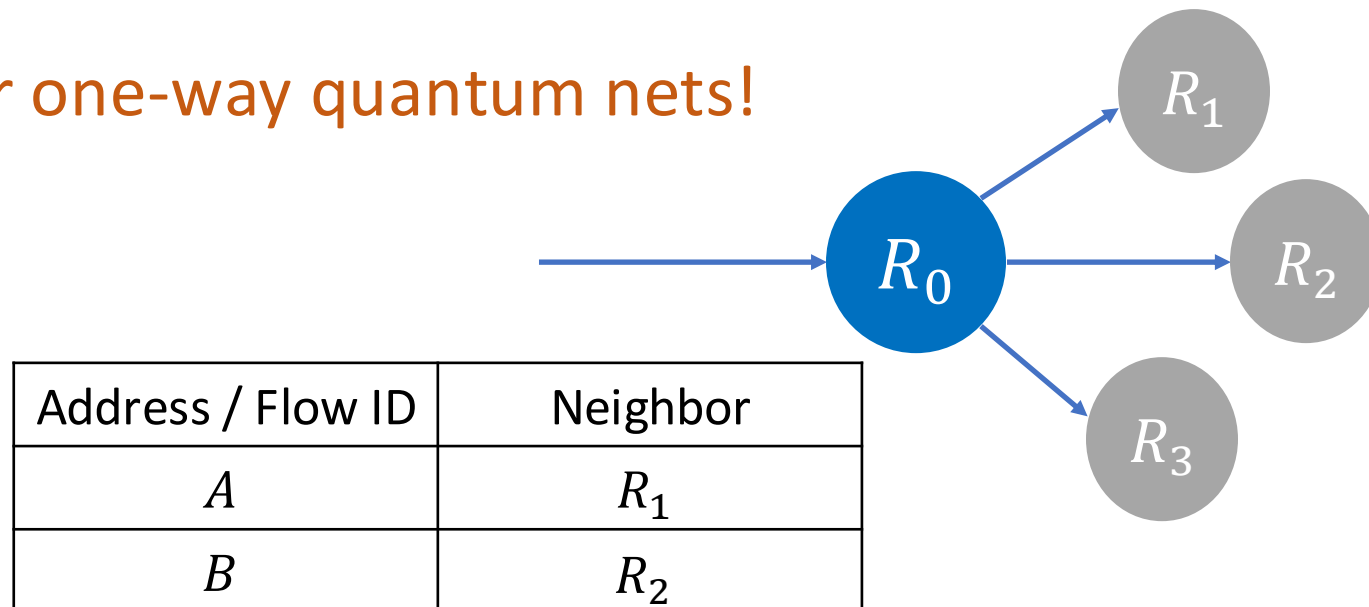
Find path between
any nodes!



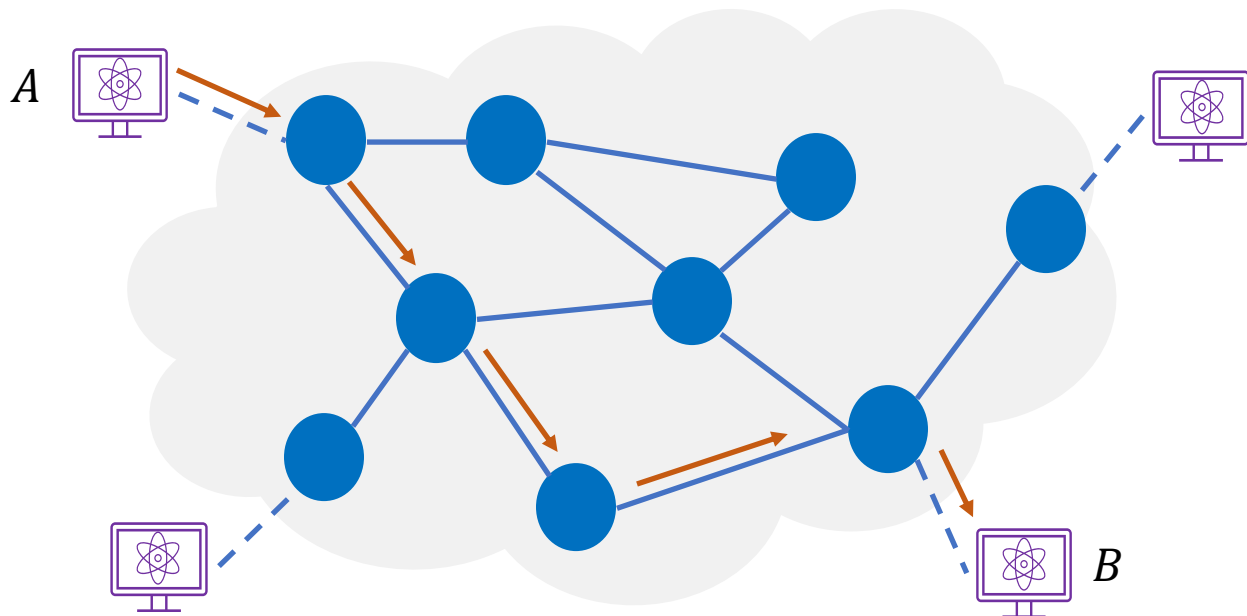
How is it solved in classical?

- Answer depends if circuit/packet switching
- For simplicity, **focus on find paths**
- Assume centralized controller knows topology
- Run shortest path algorithm
- Routers maintain tables with forwarding information

Works for one-way quantum nets!



Do two-way networks differ?



A and B want to gen entanglement

Controller finds path \mathcal{P}

Distribute entanglement with
favorite scheduling policy

Repeat until success

Methods work, although how much
time its required?

Circuit switching / packet switching?

- Caveat: links are channels, not entanglement
- Connecting end-nodes requires LLE
- Process is probabilistic

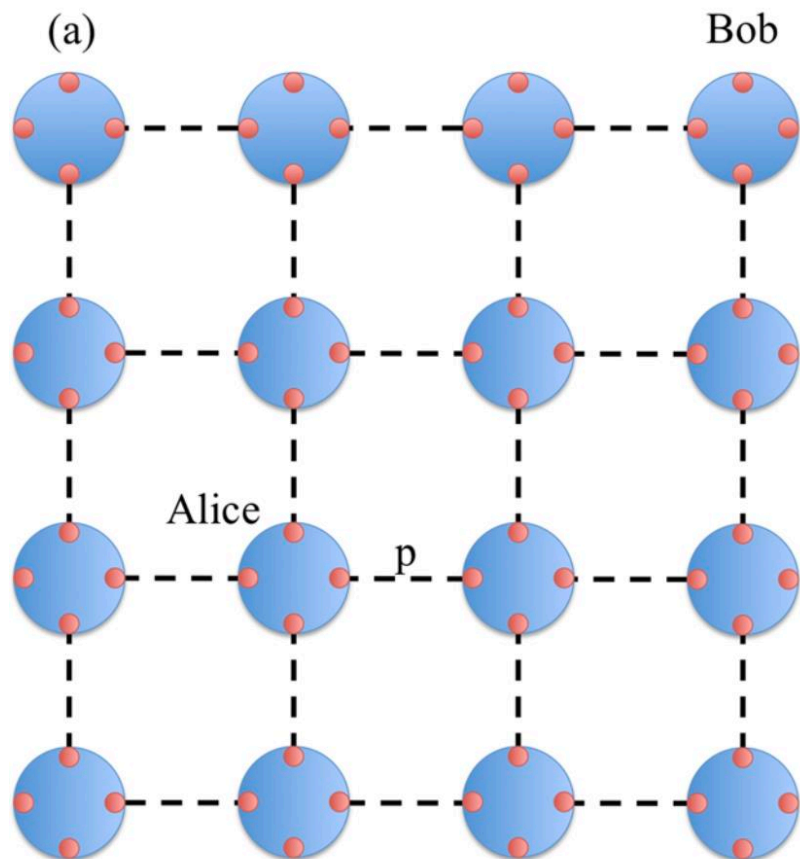
Back-of-the-Envelope Analysis



- Assume 2-phase time-slotted model (LLEG and swap)
- If memories only last for 1 slot -> rate decays exponentially with distance
- If probabilistic swaps -> only nested allows for non-exponential decay
- Precise scaling depends on multiplexing, coherence times, attempt frequency...

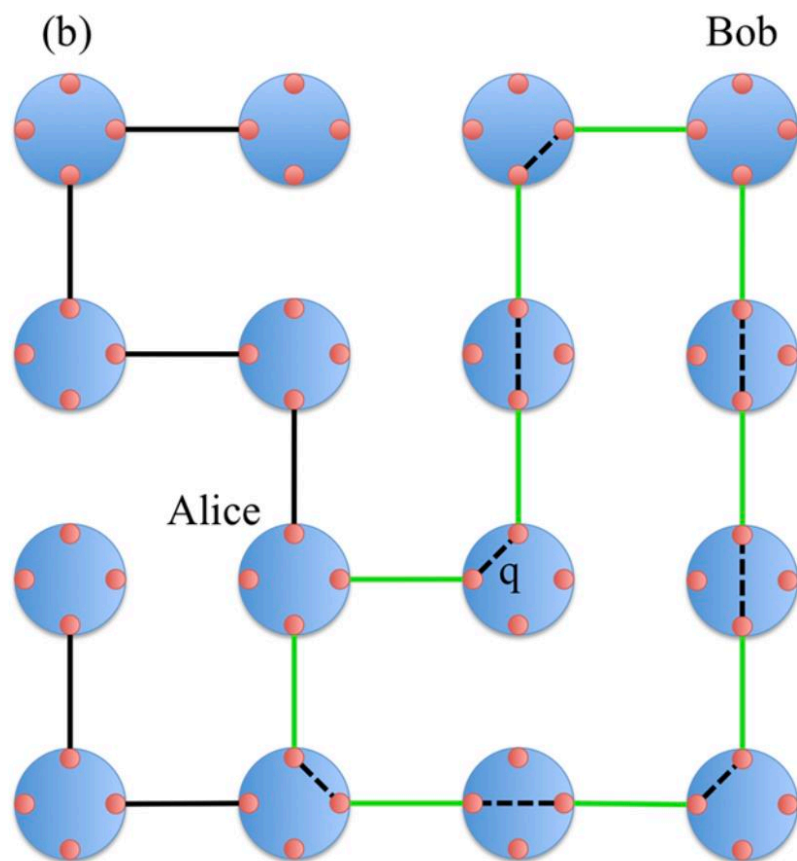
In classical, time scales linearly with distance / rate is constant

A different paradigm



- Slotted time: 1 slot allows LLEG and swap
- Continuous LLEG attempts
- One qubit per link per router
- Swap succeeds with prob. q
- Memory holds states for 1 slot
- Route with successful entanglement through multiple paths!

Routing Decisions



Find set of edge-disjoint paths
between Alice and Bob

Perform swaps according to paths
(1-to-1 mapping between edges and qubits)

Success probability is $\sum_{l \in \mathcal{P}} p_l$ given \mathcal{P}

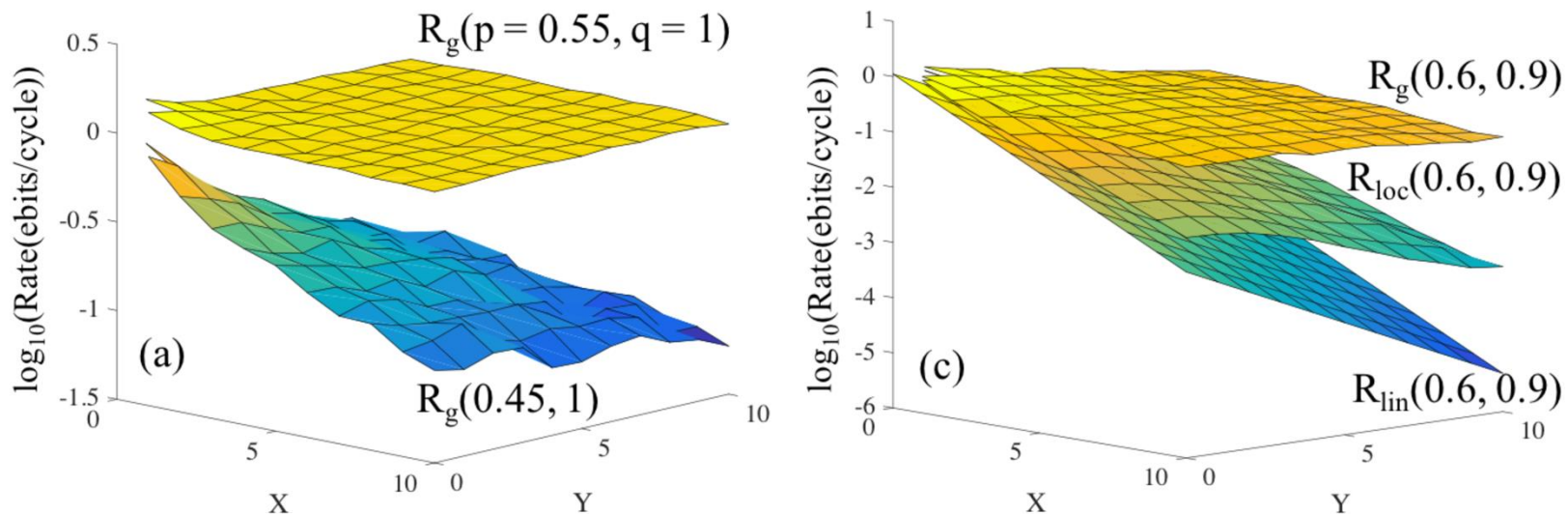
$$p_l \propto q^{|l|}$$

Requires global knowledge on LLEG,
although local information can be used!

Analysis

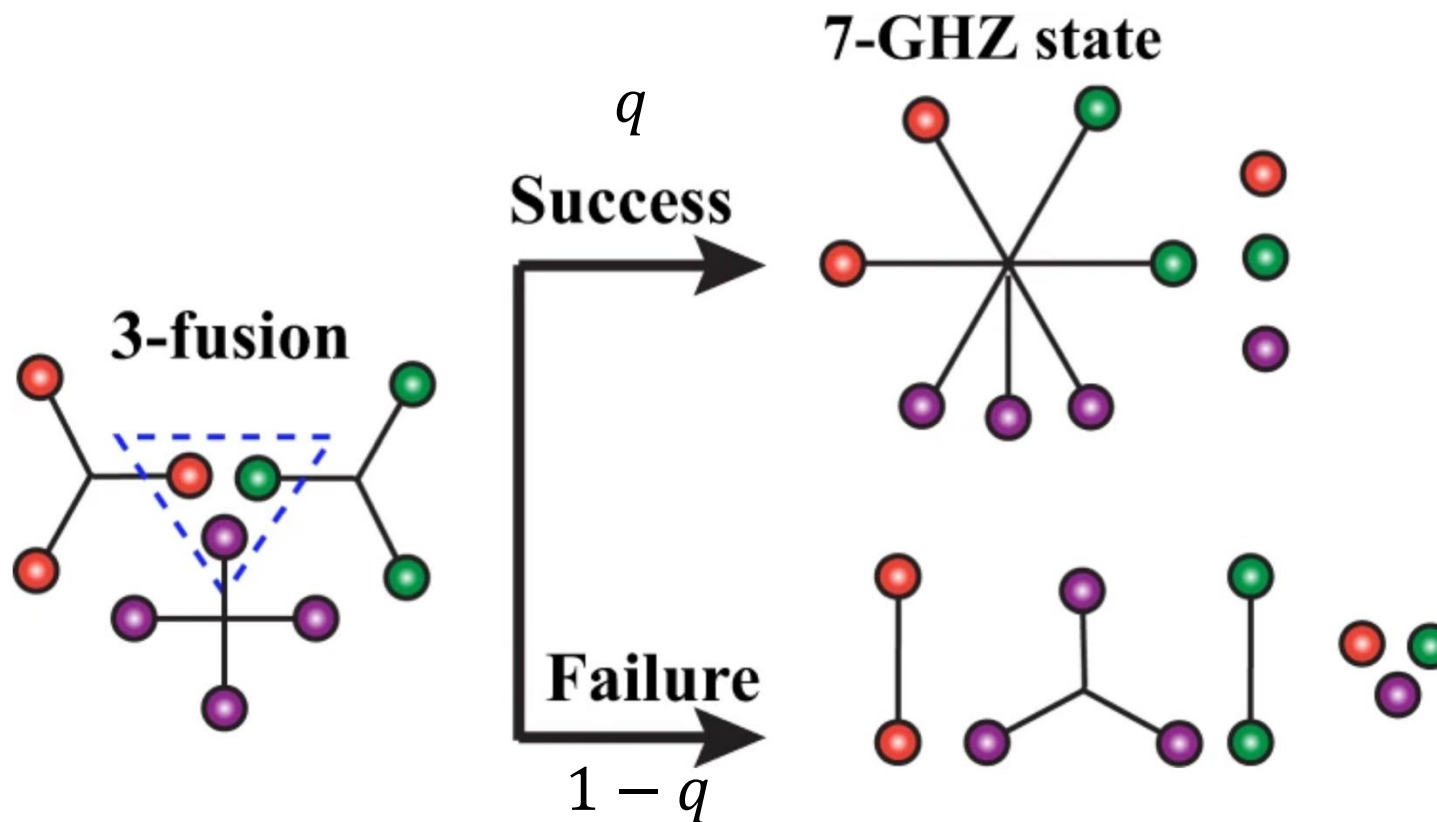
- Performance (rate) analyzed using **percolation theory**
- When $q = 1$, $p > p_{thresh}$ yields distance independent rates
- When $q < 1$, rate decays exponentially with distance
- Intuition: utilize all network resources available to serve user pair!

Distance-independent rate with $q < 1$?

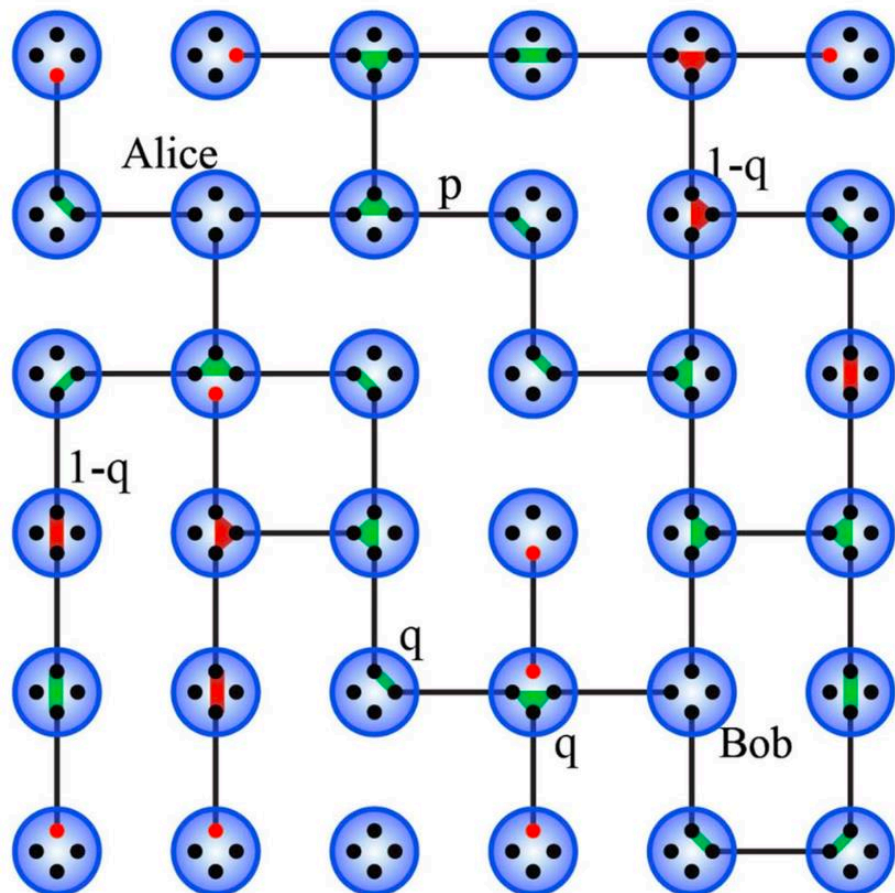


GHZ Fusions

- N -qubit GHZ states “generalizes” Bell pairs
- Entanglement swapping becomes GHZ fusion



Routing with GHZ



All nodes besides Alice and Bob
perform GHZ fusions

Approach resembles broadcast

If p and q sufficiently large, percolation!
Distance independent rate even when $q < 1$

Problem deals with probabilistic generation,
although not resilient to noise!

Different protocols obtained depending on
how fusions are performed!

Connectionless Architecture for 2-way Quantum Networks

A connectionless Qnet architecture



- little or no flow state at switches/repeaters
 - each entanglement request potentially handled separately

 - swap ASAP
 - nested swapping
- } requires path-level synchronization (connection state)
- sequential swapping
 - link resources given to flow **only** when request needs them
 - only requires link-level synchronization

Analogous to **datagram service** in classical Internet

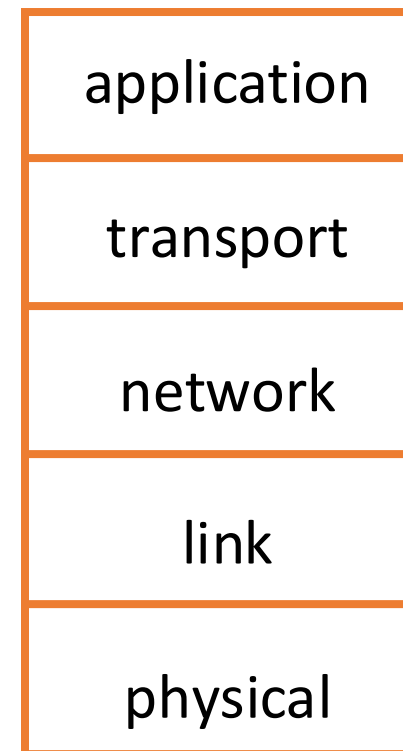
Pros



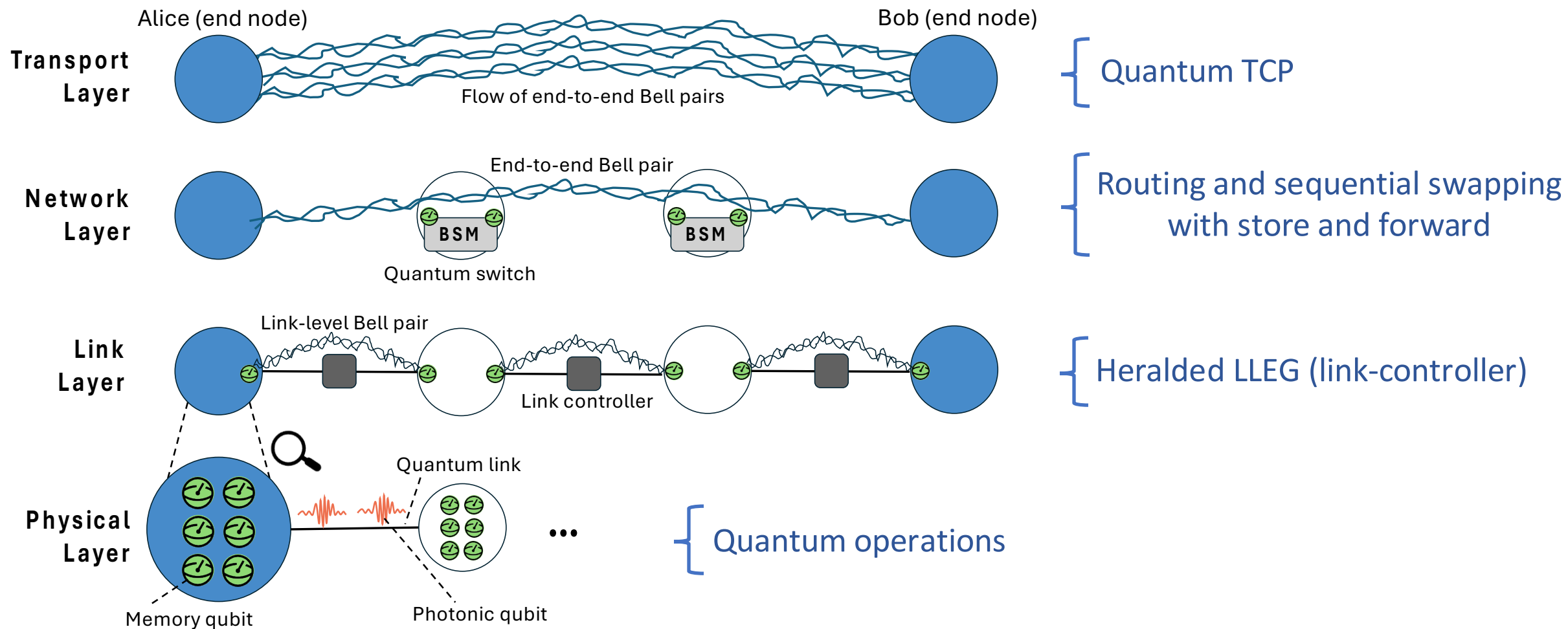
- simplifies synchronization; link by link classical comms
- amenable to analytics (understandability, management)
- allows deployment of classical Internet protocols
 - Q-TCP
 - destination-based routing
 - multicast \leftrightarrow multipartite state distribution

Connectionless Architecture

- Packet switching – Store and Forward
- TCP/IP network stack
- Statistical multiplexing enables multiple users
- Fundamental for Internet performance

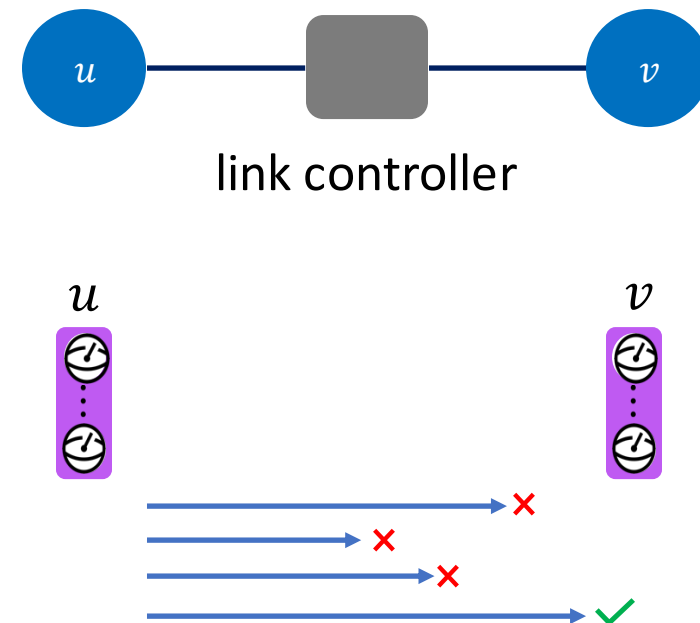


Architecture



Link Layer

- Heralded LLEG (multiplexed)
- Nodes requests entanglement from controller
- Controller orchestrates generation
- Entanglement indexed by request ID and link-level label
- **Can accommodate purification!**

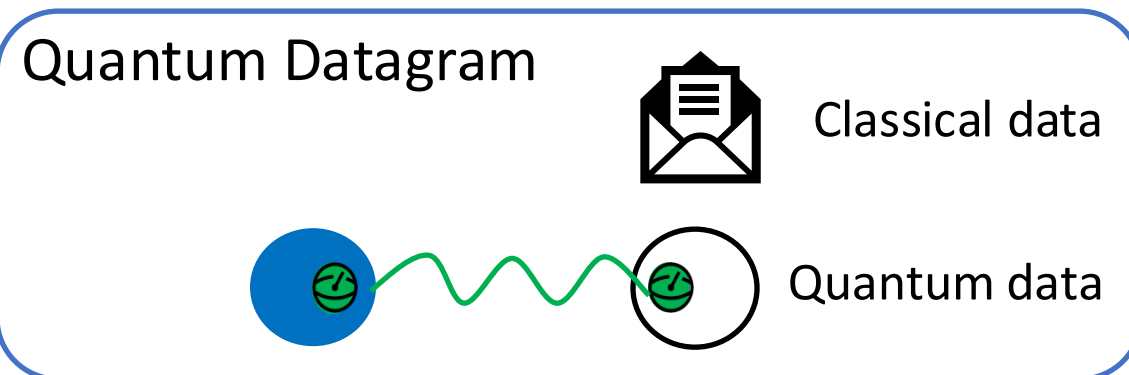
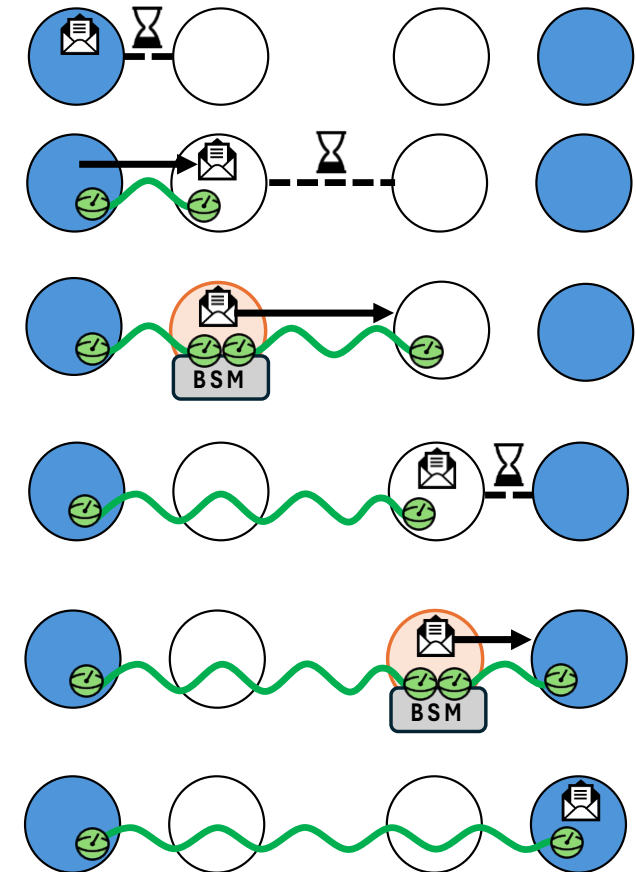


Entangled state ($ID, label$) generated for u

Link-layer generated entangled links for
network layer to consume!!!

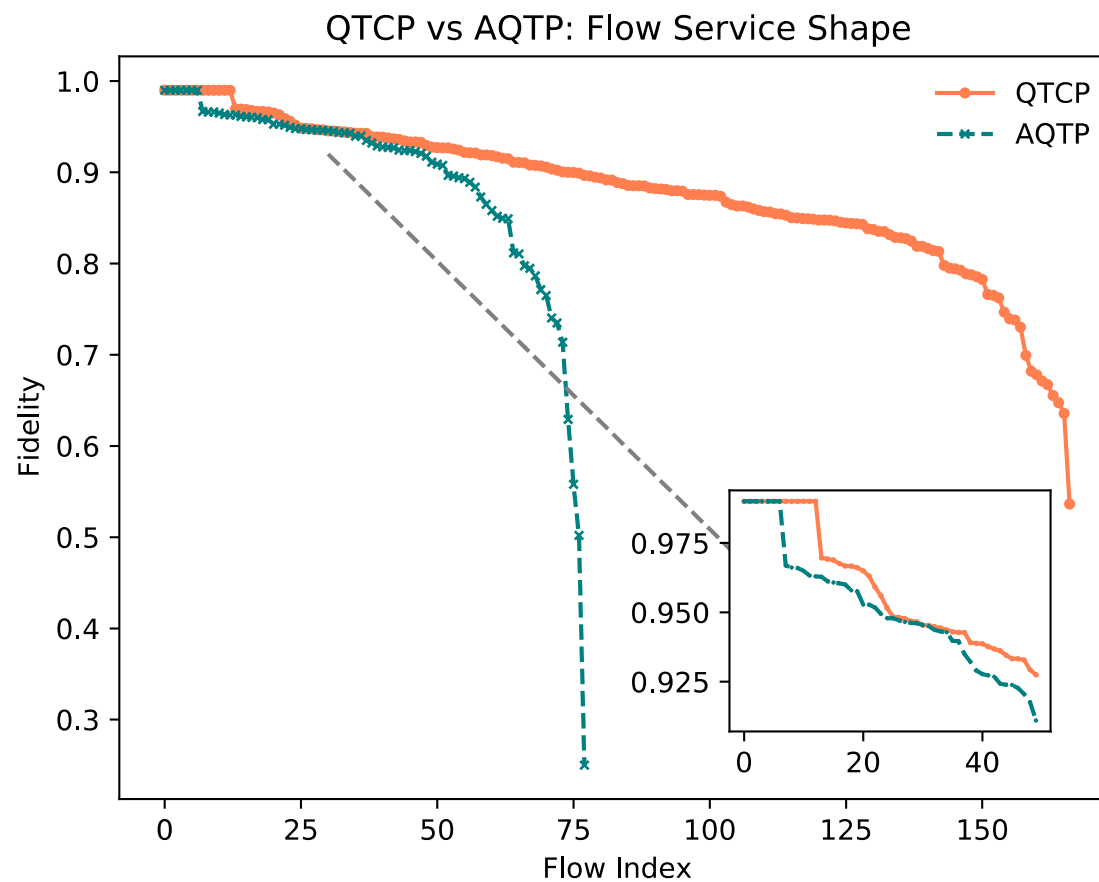
Network Layer

- Forward Bell states like packets via sequential entanglement swapping
- Quantum datagram
- No resource reservation
- Best-effort service
- Supports different schedules for link-level entanglement consumption



Transport Layer: QTCP

- TCP-like: opening handshake, data flow, closing handshake
- Ack's for delivered q-datagrams
- Control E2E q-datagram rate (congestion control, AIMD)
- Active queue management reduces time states decohere



Poll Question

What are the benefits of a connectionless quantum network architecture? Select all that apply.

- A. Requires minimal synchronization
- B. Reduces total amount of required classical communication.
- C. Permits borrowing ideas from classical Internet for application to quantum Internet
- D. Provides distance independent entanglement rates
- E. I don't know

Answer

What are the benefits of a connectionless quantum network architecture? Select all that apply.

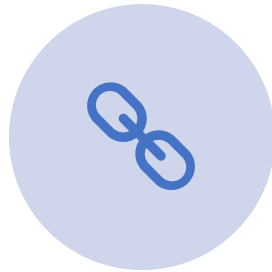
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Quantum Network Tomography

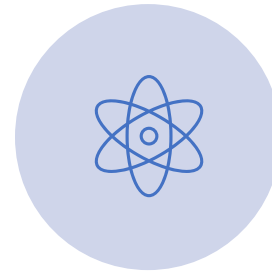
Outline



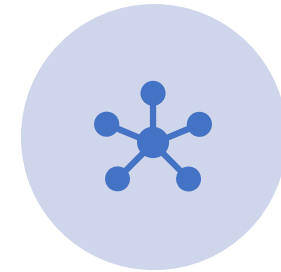
NETWORK MANAGEMENT
AND TOMOGRAPHY
OVERVIEW



CLASSICAL NETWORK
TOMOGRAPHY



QUANTUM NETWORK
TOMOGRAPHY (QNT)

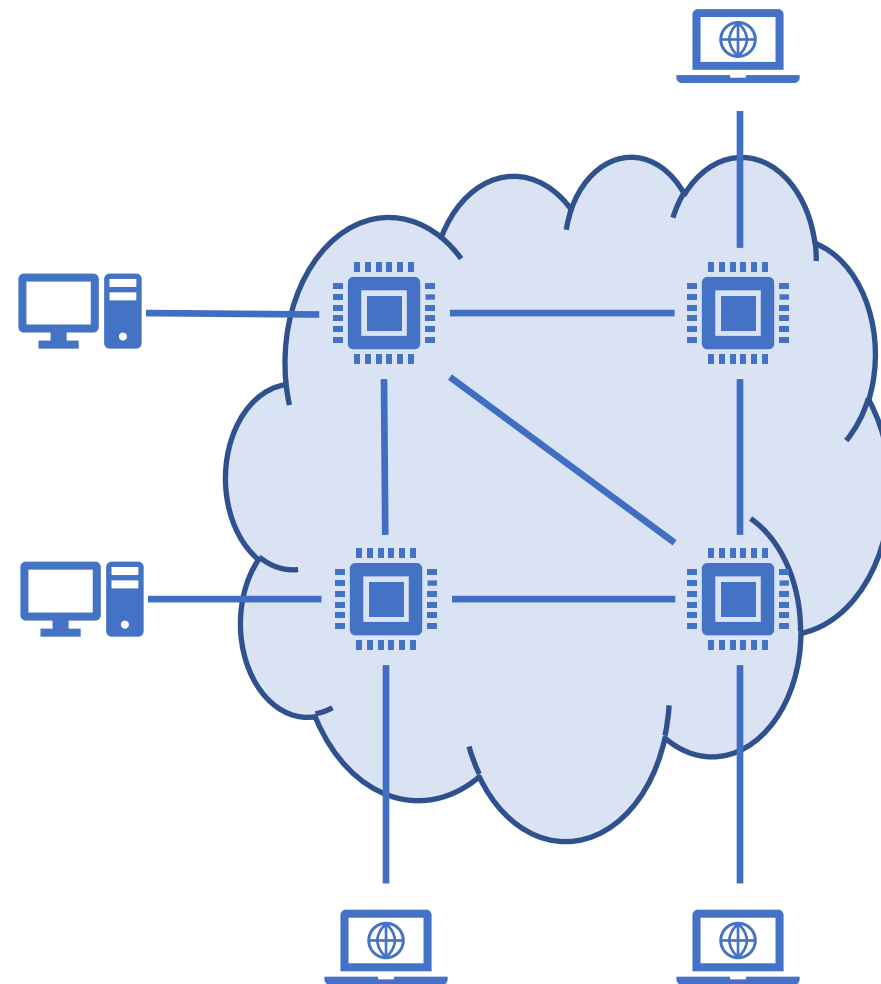


CHARACTERIZING STAR
NETWORKS

Network management



- Network component data collection
- Information to aid decision making
- Fault-detection for hardware / software
- Determine traffic patterns



Network tomography

Goal

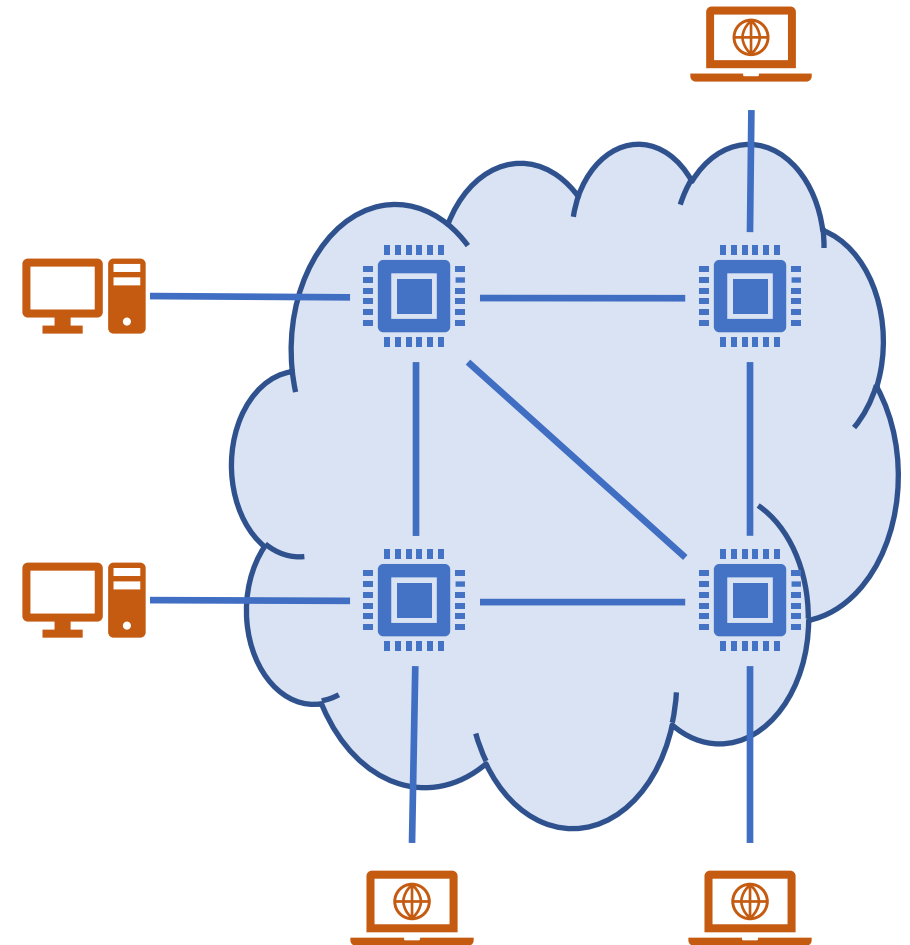
Infer internal behavior in network
from external nodes

In practice

Estimate error parameters for internal
components from end-to-end
measures

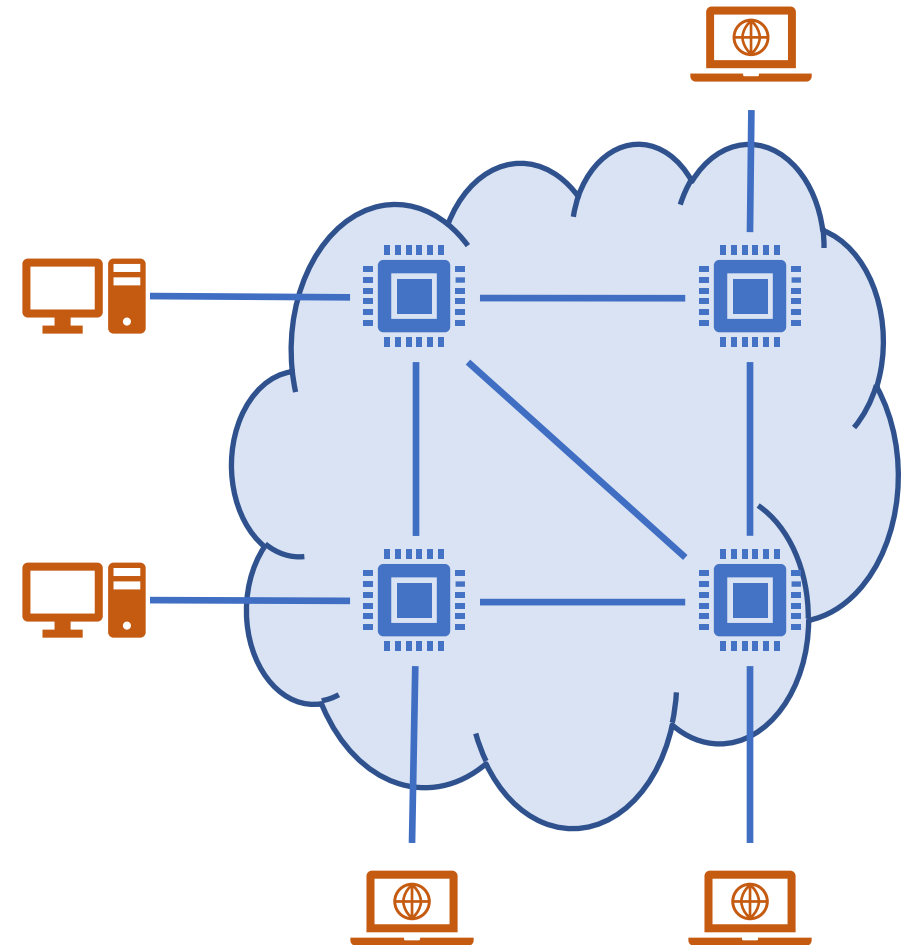
Identifiability

Obtain one value for parameters given
a set of observations



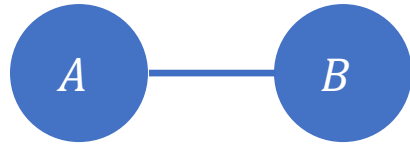
Why end-to-end?

- No participation by network needed
 - Measurement probes regular packets
- No administrative access needed
- Inference across multiple domains
 - No cooperation required
 - Monitor service level agreements
- Reconfigurable applications
 - Video, audio, reliable multicast



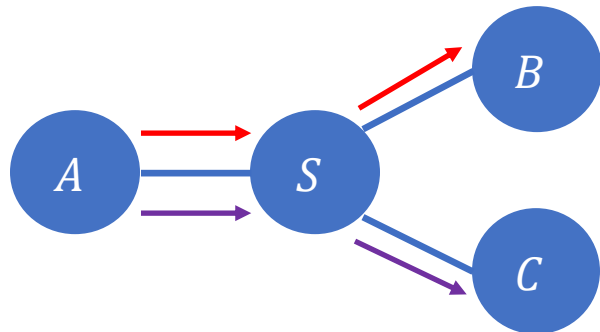
Definitions

Link-level metrics



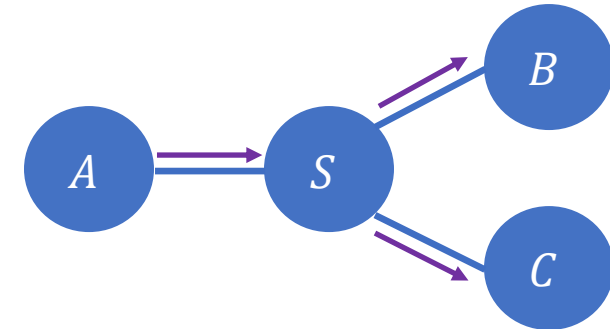
E.g: delay, loss, bit-flip rate

Unicast communication



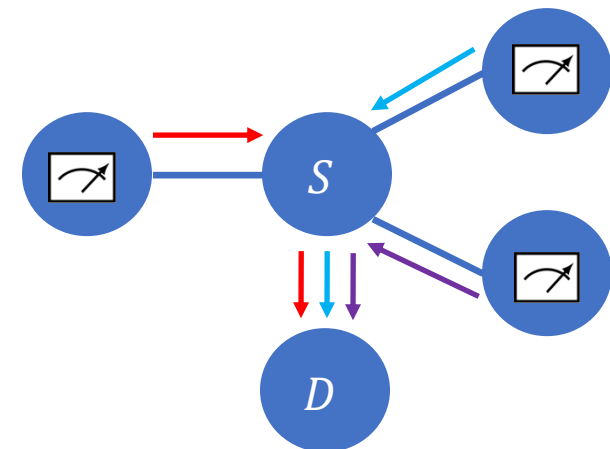
One-to-one

Multicast communication



One-to-many

Estimation



Data sent to fusion center

Unicast Tomography

Assumptions

- Links are **symmetric**
- Additive metrics

Results

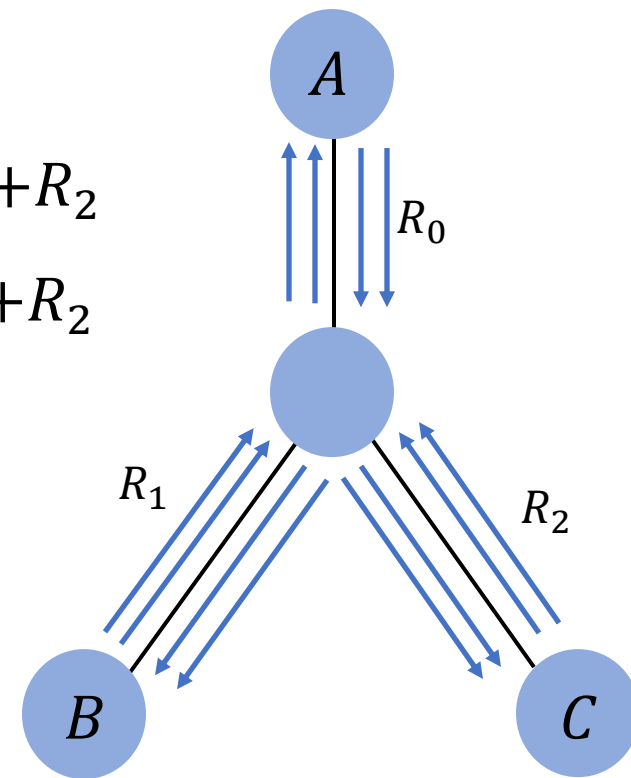
- Linear independence!
(**identifiable**)
- True for general trees
- Can infer some link delays within
general graph
- Measurements over cycles

Routing Matrix

$$R_{AB} = R_0 + R_1$$

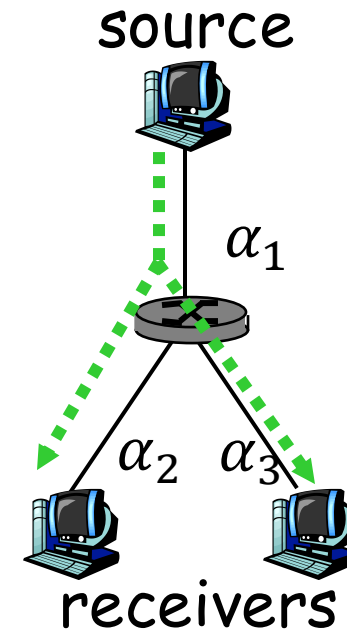
$$R_{AC} = R_0 + \quad + R_2$$

$$R_{BC} = \quad R_1 + R_2$$



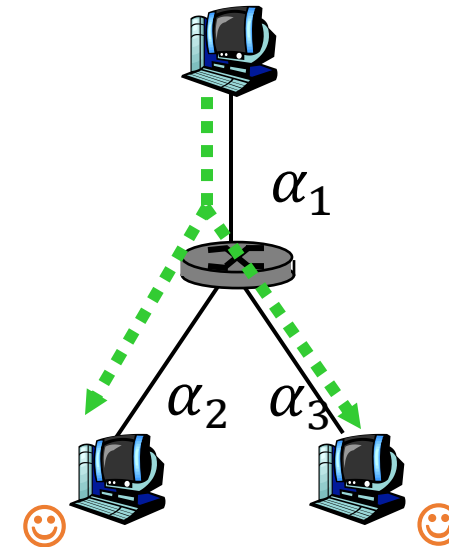
Multicast Tomography

- Multicast probes
 - copies made *as needed* within network
 - Communication through trees
- Receivers observe correlated performance
- *Exploit* correlation to get link behavior
 - Loss rates
 - Delays



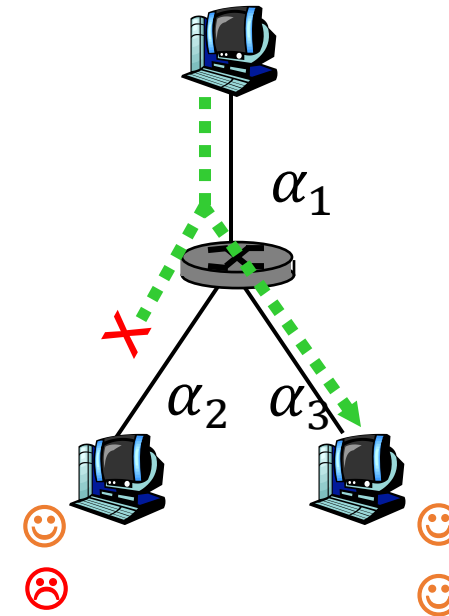
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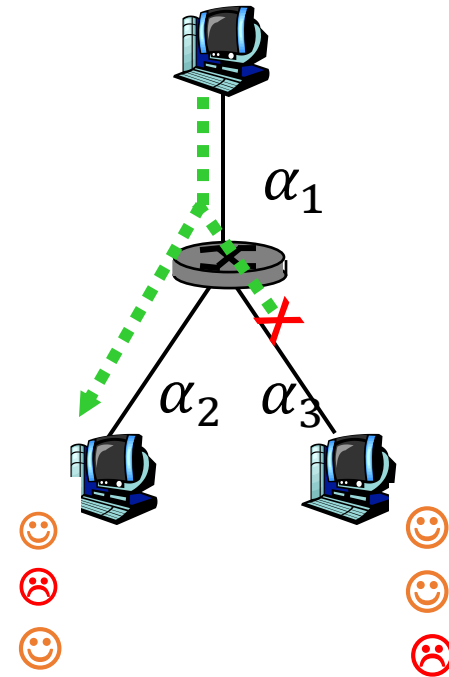
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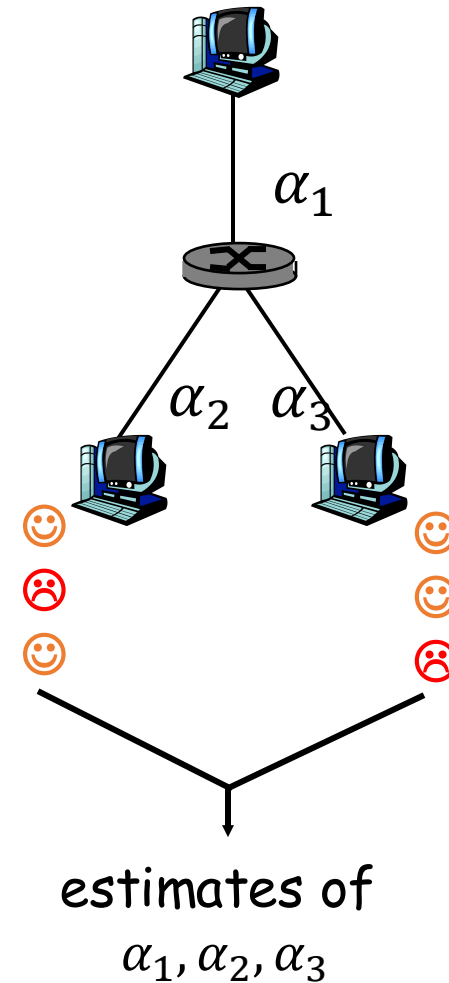
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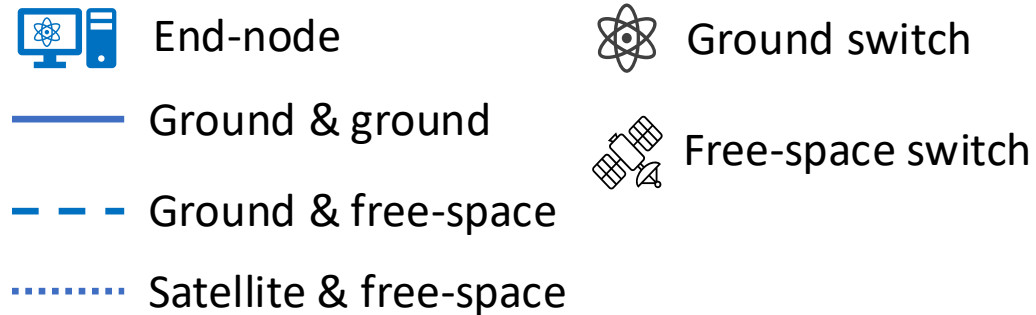
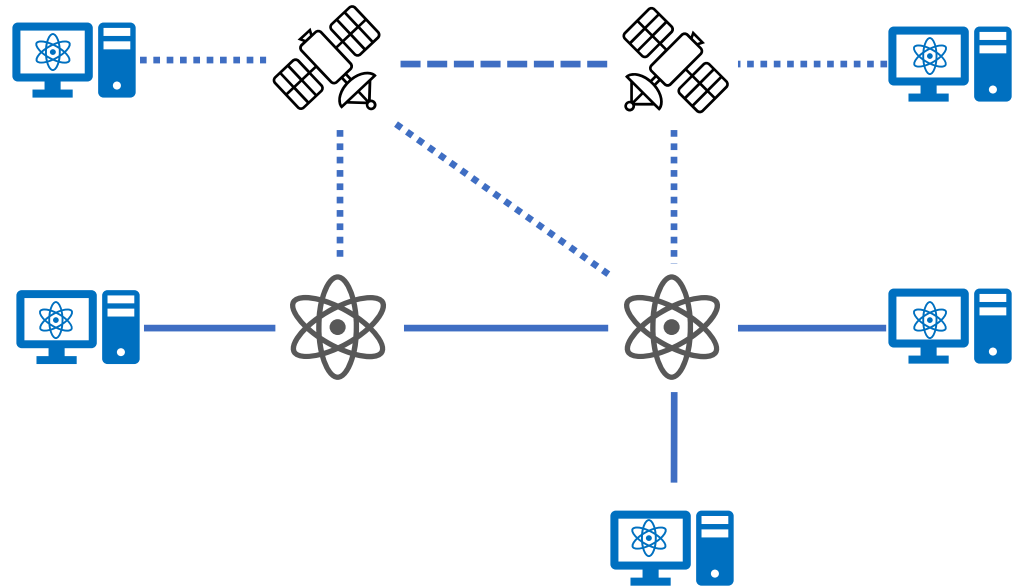


Multicast Tomography

- Multicast probes
 - copies made *as needed* within network
 - Communication through trees
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- *Exploit* correlation to get link behavior
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 - Delays



Quantum Network Tomography



How to characterize links with end-to-end measurements?

End-nodes

- Perform quantum circuits
- Request network state distribution
- Specify circuits for intermediate nodes

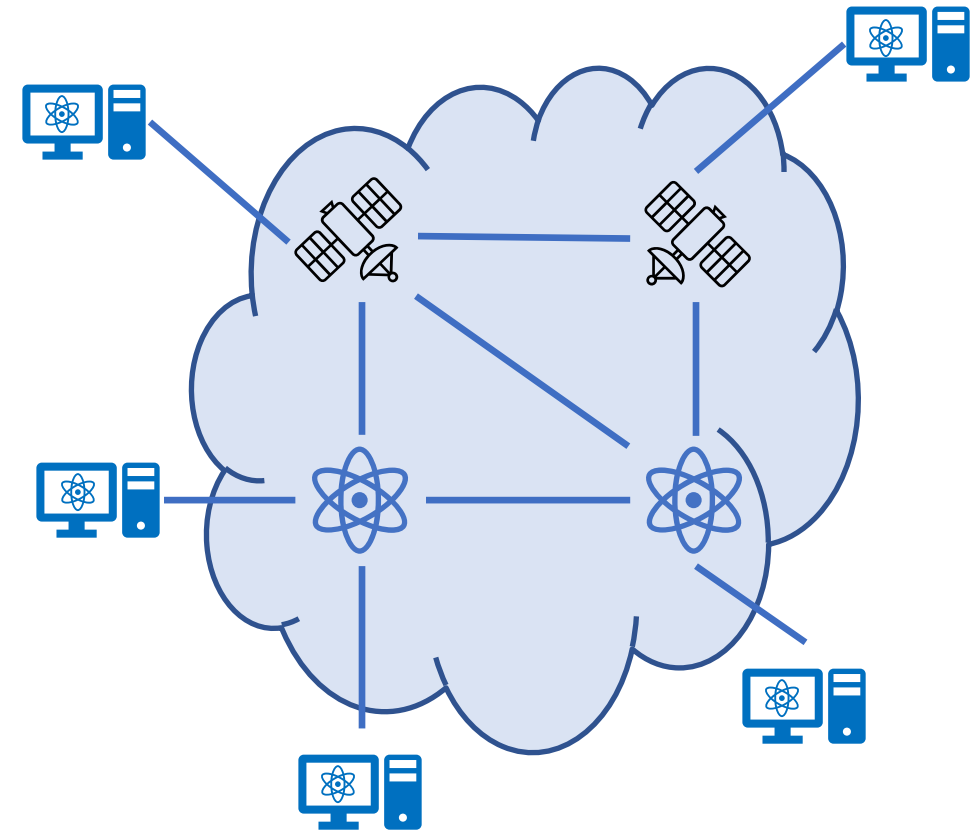
Intermediate nodes

- Receive requests for circuits
- Ancilla qubits
- No measurements for estimation

Motivation



- Inhomogeneous quantum hardware
- Hybrid communication media
- Network management
 - Faulty network hardware identification
 - Improved decision-making in resource utilization
 - Noise-informed quantum error correction
- Quality assurance
- Reconfigurable applications



From Classical to Quantum



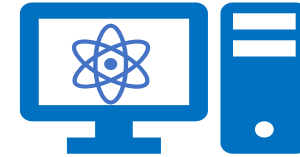
Classical	Quantum
Link-level metrics	Quantum channel parameters
Probes	State Distribution
Unicast	Bipartite state distribution
Multicast	Multipartite state distribution
End-to-end measurements	Measurements in end-nodes

Operational Assumptions



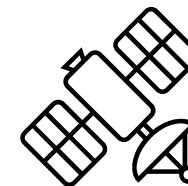
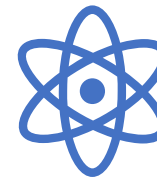
End-nodes V_E

- Perform quantum circuits
- Request network state distribution
- Specify circuits for intermediate nodes



Intermediate nodes V_I

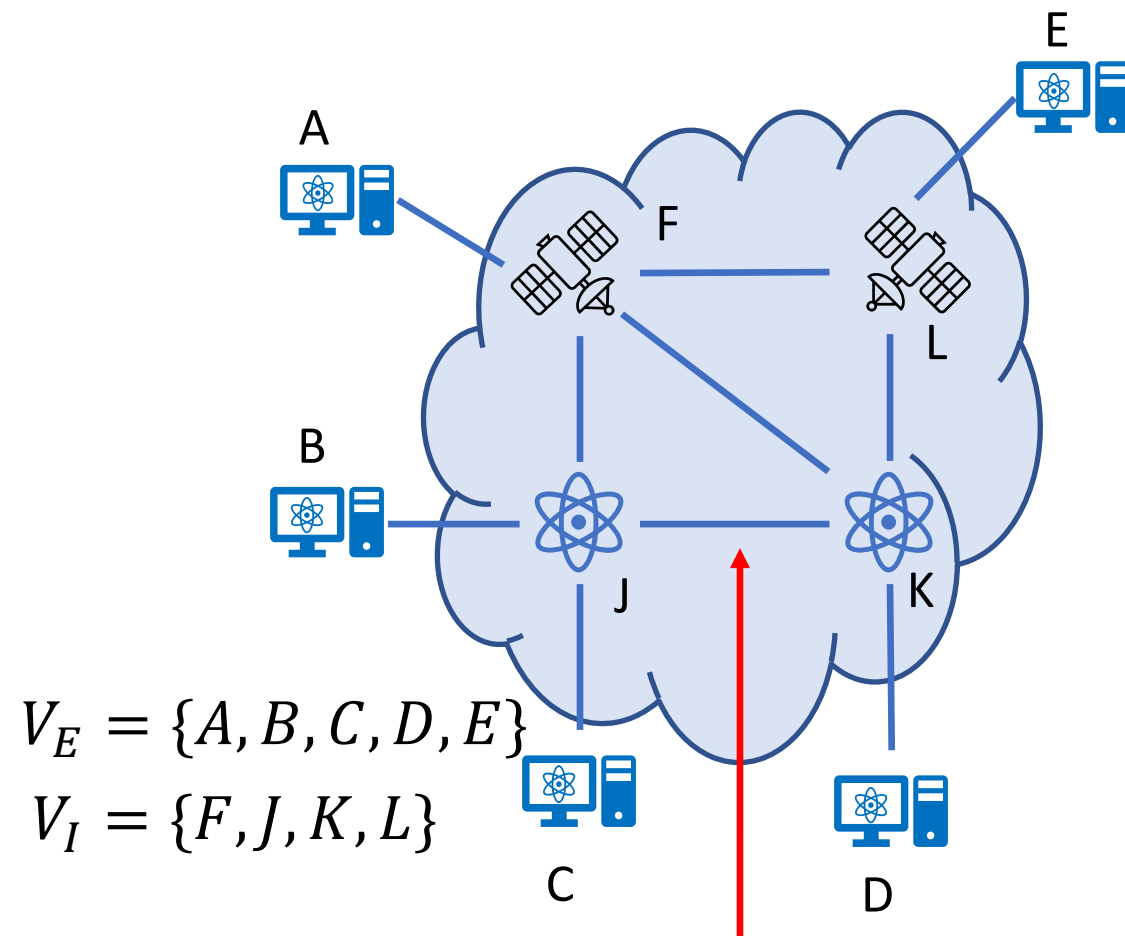
- Receive requests for circuits
- Ancilla qubits
- No measurements for estimation



Quantum Network Model

- Network is graph $G = (V, E)$
 - V : quantum processors
 - E : fiber optics, free space channels
- End and intermediate nodes
- Links: single-qubit quantum channels
- One-way quantum transmission

Goal: estimate θ from E2E measurements!

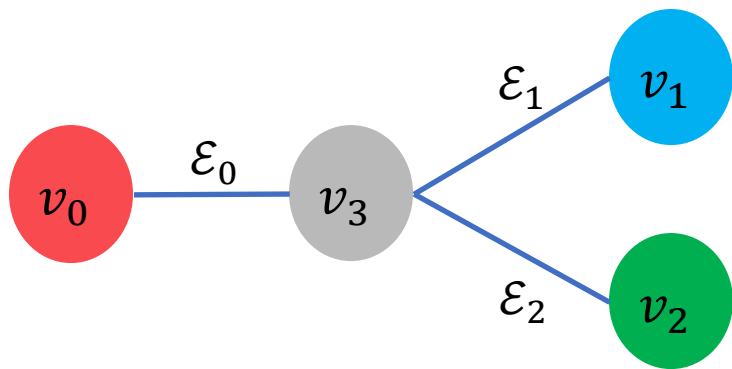


$$V_E = \{A, B, C, D, E\}$$

$$V_I = \{F, J, K, L\}$$

$$\mathcal{E}_e(\rho) = \sum_k \theta_{ek} \sigma_{ek} \rho \sigma_{ek}$$

Characterizing bit-flip stars



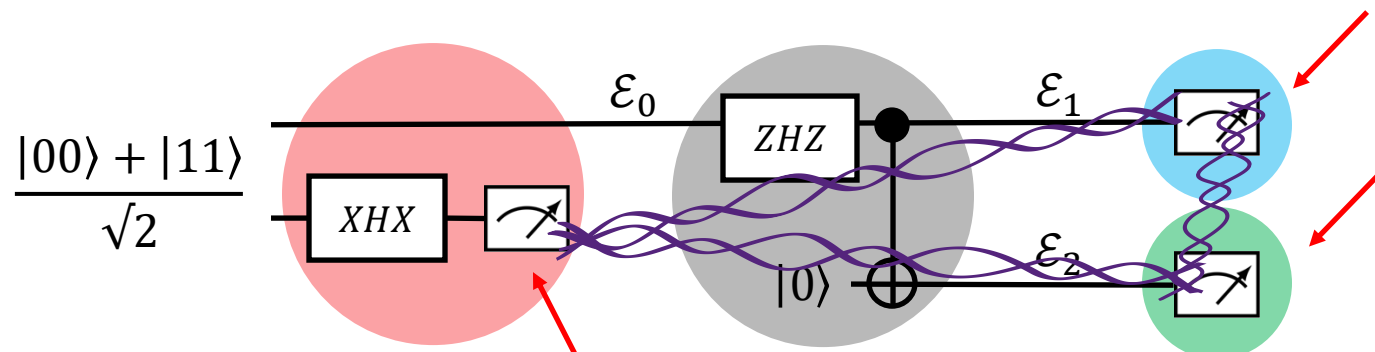
$$\mathcal{E}_e = \theta_e \rho + (1 - \theta_e) X \rho X$$

$|0\rangle \rightarrow |1\rangle$ with prob. θ_e

$|1\rangle \rightarrow |0\rangle$ with prob. θ_e

GHZ measurement implement with pre-shared entanglement

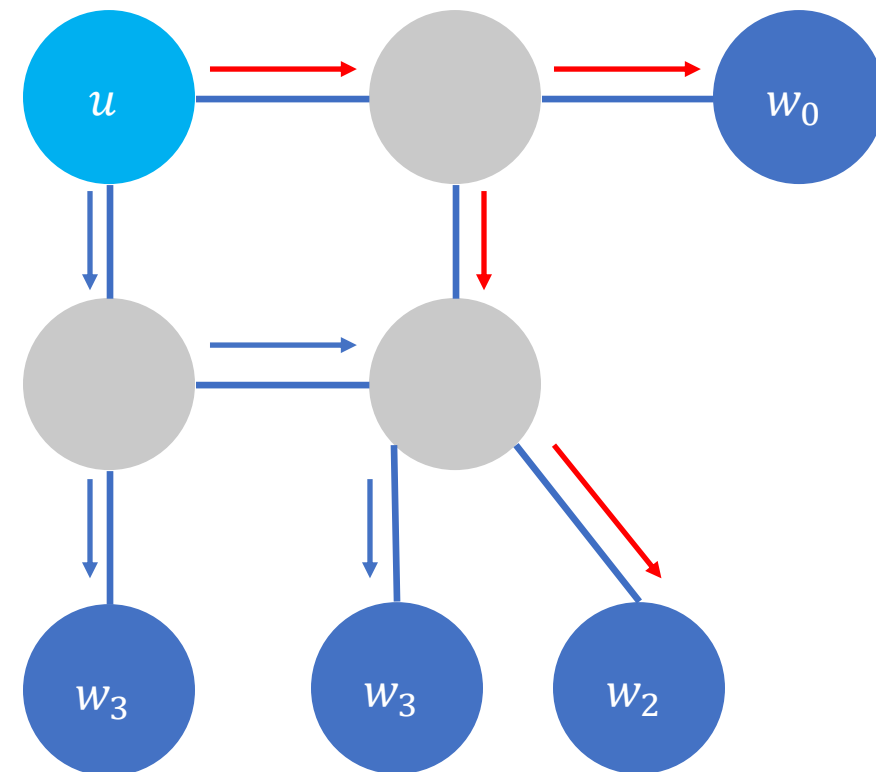
Probe: Independent Encoding



Probability of measuring $|1\rangle$ in v_i is θ_i !

- State preparation for stars generalizes to rooted trees of G
- Characterize links in tree
- Characterize graphs through tree covering
- Compatible with one-way, two-way architectures

Active area of research!



Challenges



- What are the optimal estimation strategies for stars?
- How to generalize estimators for arbitrary trees?
- How to partition network in trees for estimation?
- How do bipartite and multipartite compare?
- Under which conditions entanglement provides advantage?
- Under which conditions are trees identifiable?
- How to generalize efficient estimators for Pauli channels?

Recap



- Fundamentals of quantum communications
- Classical and Quantum Networks
- Scheduling in repeater chains
- Routing in quantum networks
- Connectionless quantum networks
- Quantum network tomography

Thank you!